

FM 5012
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the Generalized Minimal Residual (GMRES) algorithm for solving non-symmetric linear systems.*

this method was introduced by Saad and Schultz in 1986 to solve iteratively matrix equations

$$A x = b,$$

where the matrix A is not symmetric.

① Motivation: the method of conjugate gradients as a Galerkin method.

We have seen that if the k -th iterate of the CG method is x^k , the error, $e^k = x - x^k$ satisfies the equality

$$w^T A e^k = 0 \quad \forall w \in K_k,$$

where

$$K_k = \text{span} \left\{ A^i r_0 \right\}_{i=0}^{k-1}, \quad r_0 = Ax_0 - b.$$

* From: Iterative methods for sparse linear systems, Y. Saad.

We immediately see that the above equation states that the k -th residual,

$$r_k = Ax_k - b,$$

is orthogonal to K_k since

$$\begin{aligned} w^T r_k &= w^T (Ax_k - b) \\ &= w^T (Ax_k - Ax) \\ &= -w^T A e_k \\ &= 0 \quad \forall w \in K_k. \end{aligned}$$

Now, let $\{v_i\}_{i=1}^k$ be an l_2 -orthonormal basis of K_k and set

$$V_k = (v_1, v_2, \dots, v_k).$$

V_k is an $N \times k$ matrix, where $N = \text{order of } A$. With this notation, we can write that

$$x_k = x_0 + V_k y_k,$$

where $y_k \in \mathbb{R}^k$ is the solution of

$$(V_k^T A V_k) y_k = -V_k^T r_0$$

Note that the above equation defines a Galerkin approximation, $\mathbf{V}_k \mathbf{Y}_k$, in K_k to the solution of

$$\mathbf{A} \mathbf{z} = -\mathbf{r}_0.$$

Of course, we have that $\mathbf{x} = \mathbf{y}_0 + \mathbf{z}$!

In other words, the iterates of the CG method are of the form $\mathbf{x}_k = \mathbf{x}_0 + \mathbf{z}_k$, where \mathbf{z}_k is the \mathbf{A} -projection of \mathbf{z} into the space K_k .

Note that if we set

$$\mathbf{v}_k = \mathbf{r}_0 / \| \mathbf{r}_0 \|,$$

we have

$$\begin{aligned}\mathbf{x}_k &= \mathbf{x}_0 + \mathbf{V}_k \mathbf{Y}_k, \\ \mathbf{Y}_k &= -\| \mathbf{r}_0 \| \mathbf{H}_k^{-1} \mathbf{e}_1, \\ \mathbf{H}_k &= \mathbf{V}_k^T \mathbf{A} \mathbf{V}_k.\end{aligned}$$

This is a rewriting of the CG method that is possible to use for matrices that are not symmetric and positive definite. Note that if k is small, the matrix \mathbf{H}_k is easy to invert!

Arnoldi's orthogonalization. (of K_k).

We can obtain the basis $\{v_i\}_{i=1}^k$ of the Krylov subspace K_k as follows:

$$(1) \quad \hat{v}_1 = r_0, \quad v_1 = \hat{v}_1 / \|\hat{v}_1\|.$$

(2) For $j = 1, \dots, k-1$,

$$(3) \quad h_{xj} = v_x^T A v_j \quad x = 1, \dots, i \\ \hat{v}_{j+1} = A v_j - \sum_{i=1}^j h_{xj} v_i$$

$$(4) \quad h_{j+1,j} = \|\hat{v}_{j+1}\|, \quad v_{j+1} = \hat{v}_{j+1} / \|\hat{v}_{j+1}\|.$$

endFor

Notice that if $i \leq j$,

$$v_l^T \hat{v}_{j+1} = v_l^T A v_j - \sum_{i=1}^j h_{xj} v_l^T v_i$$

So, if $v_l^T v_i = \delta_{li}$ for $i, l \leq j$ then we have that $v_l^T \hat{v}_{j+1} = 0$. Hence, if $\hat{v}_{j+1} \neq 0$, we have that $v_l^T v_i = 0$ for $i, l \leq j+1$.

Next, let us show that if $\dim K_k = k$, then the algorithm (1), (2) does not break down.

Let us proceed by induction on k . For $k=1$, the result is obvious. Assume the result holds for $k=m$ and let us show it holds for $k=m+1$. Since it holds for $k=m$, we have that

$$\begin{aligned} K_m &= \text{span} \left\{ A^i r_0 \right\}_{i=0}^{m-1} \\ &= \text{span} \left\{ A^{i-1} v_i \right\}_{i=1}^m \\ &= \text{span} \{ v_1, \dots, v_m \}. \end{aligned}$$

Now, this implies that

$$K_{m+1} = \text{span} \{ v_1, \dots, v_m, A^m v_1 \}.$$

By (3) and (4), for $j \leq m$,

$$\begin{aligned} \hat{v}_{j+1} &= A v_j - \sum_{i=1}^j h_{ij} v_i \\ \Rightarrow A v_j &= \sum_{i=1}^{j+1} h_{ij} v_i \\ \Rightarrow A v_j &\in \text{span} \{ v_1, \dots, v_{j+1} \} \\ \Rightarrow A^m v_1 &\in \text{span} \{ v_1, \dots, v_{m+1} \}. \end{aligned}$$

and so, $v_{m+1} \neq 0$. This implies that the algorithm does not break down.

Finally, notice that if $m > j+1$,

$$U_m^\top A V_j = \sum_{i=1}^{j+1} h_{ij} U_m^\top V_i = 0,$$

by the inductive hypothesis. This shows that H_k is an upper Hessenberg matrix which is tridiagonal if A is symmetric.

③ the Full Orthogonalization Method (FOM)

The above discussion motivates the introduction of the FOM:

- Pick the initial guess x_0 , compute $r_0 = Ax_0 - b$ and set $v_1 = r_0 / \|r_0\|$

- For $j = 1, 2, \dots, k$ do

$$h_{ij} = v_i^\top A v_j \quad i = 1, 2, \dots, j$$

$$\tilde{v}_{j+1} = A v_j - \sum_{i=1}^j h_{ij} v_i$$

$$h_{j+1,j} = \|\tilde{v}_{j+1}\|$$

$$v_{j+1} = \tilde{v}_{j+1} / h_{j+1,j}$$

end for

- Set $x_k = x_0 + V_k Y_k$ where

$$Y_k = -H_k^{-1} \|r_0\| e_1$$

Note that the algorithm breaks down only if for some $j \leq k$, $h_{j+1,j} = 0$. Next we show that this happens only if $r^j = 0$!

The relations (3) can be rewritten in matrix form as

$$AV_k = V_{k+1} \bar{H}_k.$$

Note that

$$H_k = V_k^T A V_k = V_k^T V_{k+1} \bar{H}_k,$$

hence

$$\bar{H}_k = \begin{bmatrix} H_k \\ 0 \dots 0 & h_{k+1,k} \end{bmatrix}.$$

Since

$$x_k = x_0 + V_k y_k,$$

we get

$$\begin{aligned} r_k &= r_0 + AV_k y_k \\ &= r_0 + V_{k+1} \bar{H}_k y_k \\ &= r_0 + V_k H_k y_k + h_{k+1,k} (e_k^T y_k) V_{k+1} \end{aligned}$$

$$\Rightarrow r_k = h_{k+1,k} (e_k^T y_k) r_{k+1}.$$

This means that

$$\|r_k\| = |h_{k+1,k}| |e_k^T y_k|,$$

and $r_k = 0$ if $h_{k+1,k} = 0$, as claimed.

④ The GMRES method.

The difference, and only difference, between the FOM and the GMRES methods is the way in which y_k is computed. In FOM, y_k is computed by requesting that

$$\begin{aligned} r^k &= A(x_0 + V_k y_k) - b \\ &= r_0 + A V_k y_k \end{aligned}$$

be orthogonal to K_k . In GMRES, instead, it is asked that

$$\|r^k\|^2 = \|r^k(y_k)\|^2 =: J(y_k).$$

is a minimum, hence the name of the method!

Next, notice that

$$\begin{aligned} J(y) &= \| r_0 + AY_k y \|^2 \\ &= \| r_0 + V_{k+1} \bar{H}_k y \|^2 \\ &= \| \|r_0\| e_1 + \bar{H}_k y \|^2 \end{aligned}$$

This implies that

$$y_k = -(\bar{H}_k^T \bar{H}_k)^{-1} \bar{H}_k e_1 \quad (\text{roll}),$$

and shows that this method is different than the FOM.

To actually compute y_k , we can still exploit the structure of J . Suppose that we obtain a $(k+1) \times (k+1)$ matrix Q_k , the accumulated product of rotation matrices, and an upper triangular matrix R_k such that

$$Q_k \bar{H}_k = R_k$$

(Note that the last row of R_k is zero!) - then

$$\begin{aligned} J(y) &= \| Q_k (\|r_0\| e_1 + \bar{H}_k y) \|^2 \\ &= \| \underbrace{\|r_0\| Q_k e_1}_{-g_k} + R_k y \|^2 \end{aligned}$$

Now if we write

$$\hat{g}_k = \begin{bmatrix} \hat{g}_k \\ g_{k,k+1} \end{bmatrix}, \quad \hat{R}_k = \begin{bmatrix} \hat{R}_k \\ 0 \dots 0 \end{bmatrix},$$

then $y_k = \hat{R}_k^+ \hat{g}_k$ and $J(y_k) = \|g_{k,k+1}\|$.

⑤ Convergence analysis

To analyze the GMRES method, we begin by noting that if the Arnoldi process does not break down, then \hat{R}_k is invertible. To see this, consider the beginning of the algorithm leading to the equation $Q_k \bar{F}_k = R_k$:

$$\left[\begin{array}{cc|ccc} c & -s & 0 & 0 & 0 \\ s & c & 0 & 0 & 0 \\ \hline 0 & 0 & & & \\ 0 & 0 & & \text{Id} & \\ 0 & 0 & & & \end{array} \right] \left[\begin{array}{cccc} r & x & x & x \\ h & x & x & x \\ 0 & x & x & x \\ 0 & 0 & x & x \\ 0 & 0 & 0 & x \end{array} \right] = \left[\begin{array}{cccc} x & y & y & y \\ 0 & y & y & y \\ 0 & x & x & x \\ 0 & 0 & x & x \\ 0 & 0 & 0 & x \end{array} \right]$$

$$\text{where } c = \frac{r}{\sqrt{r^2+h^2}}, \quad s = -\frac{h}{\sqrt{r^2+h^2}}, \quad t = \sqrt{r^2+h^2}$$

We see that if $h \neq 0$, then $t > 0$. Now notice that in every multiplication by a rotation, $h = h_{j+1,j}$, which is non-zero because the Arnoldi process did not break down for $j \leq k$! This establishes that \hat{R}_k is invertible.

this means that the GMRES method can only break down if for some $j < N$, $b_{j+1,j} = 0$. Next, we prove that this happens if and only if $r_j = 0$, that is, if and only if $Ax_j = b$.

Assume that $b_{j+1,j} = 0$. This implies that

$$AV_j = V_j H_j.$$

In this case, we have

$$\begin{aligned} J(\gamma) &= \|r_0 + AV_j \gamma\|^2 \\ &= \|r_0 + V_j H_j \gamma\|^2 \\ &= \|V_j (\|r_0\| e_1 + H_j \gamma)\|^2 \\ &= \|\|r_0\| e_1 + H_j \gamma\|^2 \\ &= 0 \end{aligned}$$

if

$$\gamma = -\frac{\|r_0\|}{\|H_j\|} e_1.$$

This means that $r_j = 0$, as claimed.

Assume now that $r_j = 0$ and that $r_i \neq 0$ $i < j$.

then

$$\|r_j\| = |g_{j,j+1}| = |s_j| |g_{j+1,j}| = |s_j| \|r_{j-1}\|$$

This implies $0 = s_j \Rightarrow b_{j+1,j} = 0$, as claimed.

We have thus proven the following result.

Theorem. The GMRES method converges in at most N iterations, where $N = \text{order of } A$.

We can also obtain an error estimate of the residual in some cases.

Theorem. Let A be a diagonalizable matrix.
Then

$$\|r_m\|_2 \leq K \epsilon_m \|r_0\|_2,$$

where

$$K = \|X\|_2 \|\tilde{X}\|_2, \\ \epsilon_m = \min_{P \in P^m} \max_{\substack{1 \leq i \leq N \\ P(i) = 1}} |\rho(\lambda_i)|$$

$$\text{and } \text{diag}\{\lambda_i, i=1, \dots, N\} = \Lambda = \tilde{X}^T A X.$$

Proof. Since

$$x_m = x_0 + V_m y_m,$$

we have

$$r_m = r_0 + A V_m y_m$$

and since $v_m = k_m$,

$$r_m = p(A) r_0 , \text{ degree } p = m, p(0) = 1.$$

Since x_m minimizes the residual over $x_0 + k_m$, we set

$$\|r_m\|_2 = \min_{\substack{p \in P^m \\ p(0)=1}} \|p(A)r_0\|.$$

But

$$p(A) = X p(\Lambda) X^{-1}$$

and so

$$\|p(A)\|_2 \leq k \max_{1 \leq i \leq n} |p(\lambda_i)|.$$

This completes the proof. \square

If all the eigenvalues of A are inside the ellipse of center $(c, 0)$ focal distance d and major semi-axis a

