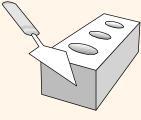


### The Entity-Relationship Model

Chapter 2



## Overview of Database Design

Requirement analysis

Iterative process

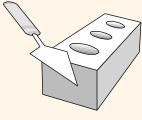
Conceptual DB design

Logical DB design

Schema refinement

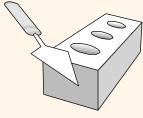
Physical DB design

App & DB security



# Overview of Database Design

- Suilding a conceptual database design should be done after informal discussions with the customers.
- Your design should reflect all the details and operations that your customer need
- Examples of customers are: University databases, Company database, Banking, Airline reservation,...



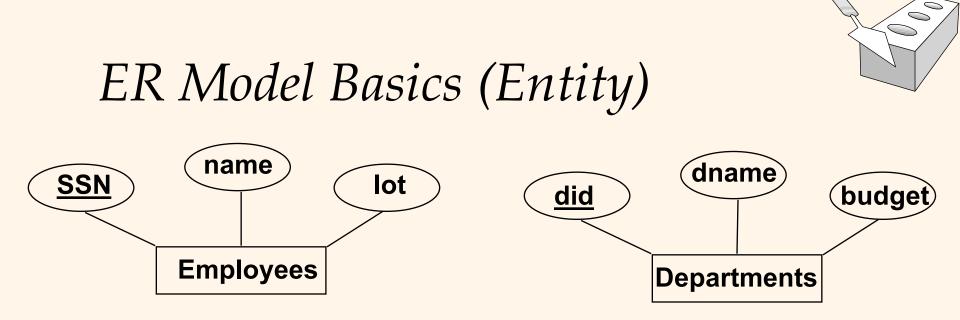
# Overview of Database Design

- \* <u>Conceptual design</u>: (ER Model is used at this stage.)
  - What are the *entities* and *relationships* in the enterprise?
  - What information about these entities and relationships should we store in the database?
  - What are the *integrity constraints* or *business rules* that hold?
  - A database `schema' in the ER Model can be represented pictorially (*ER diagrams*).
  - Can map an ER diagram into a relational schema.

### 

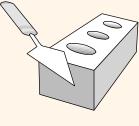
### ER Model Basics

- *Entity*: Real-world object distinguishable from other objects. An entity is described (in DB) using a set of *attributes*.
- *Entity Set*: A collection of similar entities.
  E.g., all employees.
  - All entities in an entity set have the same set of attributes. (Until we consider ISA hierarchies, anyway!)
  - Each entity set has a *key*.
  - Each attribute has a *domain*.

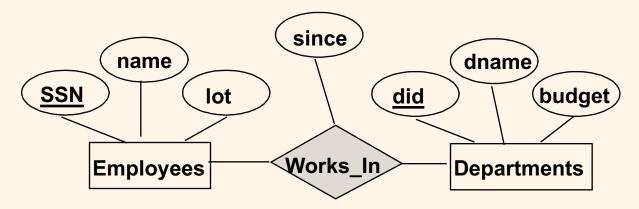


\* Based on the discussion with our customers, we find that:

- Each employee has there attributes: SSN, name, and parking lot
- Each department has three attributes: number, name, and budget
- Employees can be uniquely identified by their SSN (<u>SSN</u>), departments can be identified by their numbers (<u>did</u>).
- ✤ A *key* is a minimal set of attributes whose values uniquely identify an entity in the set
  - If there are more than one *candidate* key, we designate one of them as the *primary* key



# ER Model Basics (Relationship)

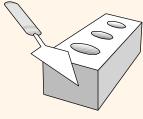


\* <u>*Relationship*</u>: Association among two or more entities.

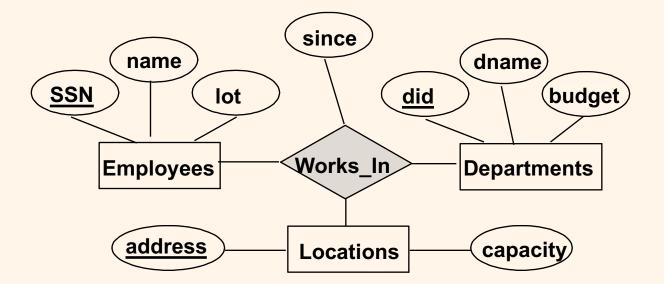
• E.g., Attishoo works in Pharmacy department.

✤ <u>Relationship Set</u>: Collection of similar relationships.

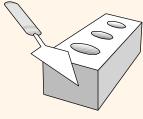
- \* A relationship set may have *descriptive* attributes:
  - E.g., John starts working in the CS department since August 2005



N-ary Relationship



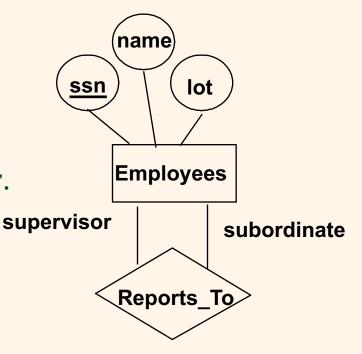
An n-ary relationship set R relates n entity sets E1 ... En; each relationship in R involves entities e1 E1, ..., En

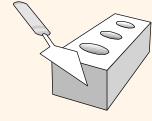


## More on Relationship Set

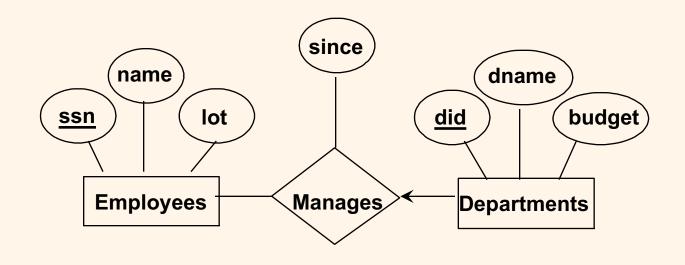
#### The entity set in a relation may not be distinct

- An Employees can supervise another employee
- In this case, a role indicator should be labeled (supervisor. subordinate)
- Same entity set could participate in different relationship sets, or in different "roles" in same set.



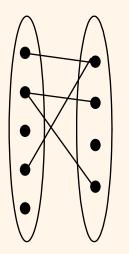


# Key Constraints



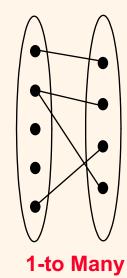
 Each dept has at most one manager, according to the <u>key constraint</u> on Manages.

### Key Constraints (Cont.)

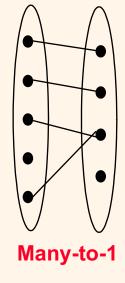


Many-to-Many

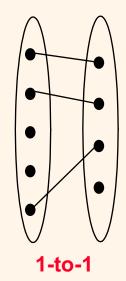
Works\_In: An employee can work in many departments; a department can have many employees.



Manages: A department is managed only by one employee. One employee can manage many departments



Works\_In: If we add a constraint that each employee can work in only one department

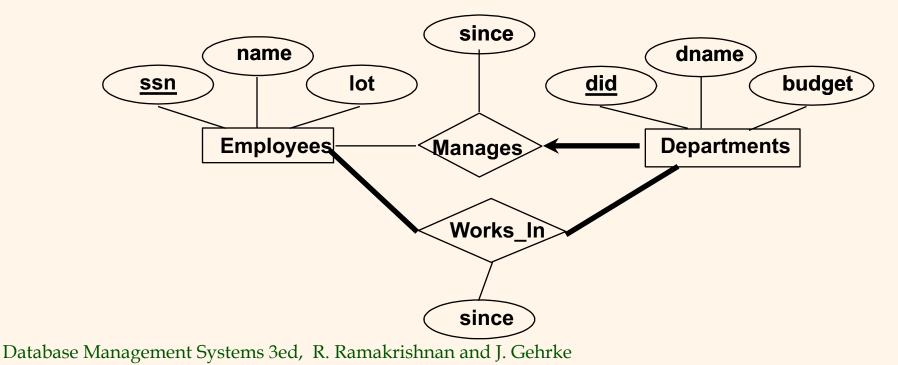


Manages: If we add a constraint that one employee can manage only one department

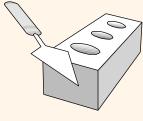
# Participation Constraints

Does every department have a manager?

- If so, this is a *participation constraint*: the participation of Departments in Manages is said to be *total* (vs. *partial*).
  - Every *did* value in Departments table must appear in a row of the Manages table (with a non-null *ssn* value!)

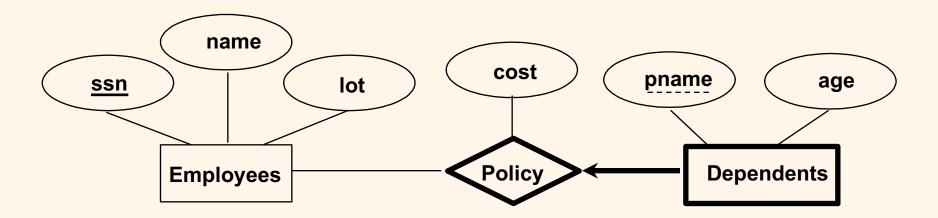


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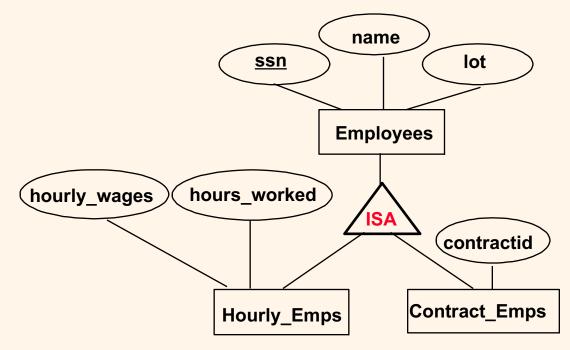


### Weak Entities

- \* A *weak entity* can be identified uniquely only by considering the primary key of another (*owner*) entity.
  - Owner entity set and weak entity set must participate in a one-tomany relationship set (one owner, many weak entities).
  - Weak entity set must have total participation in this *identifying* relationship set.



### ISA (`is a') Hierarchies

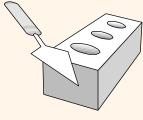


\* As in C++, or other PLs, attributes are inherited.

If we declare A ISA B, every A entity is also considered to be a B entity.

# **\*ISA** Relationship can be viewed as either specialization or generalization

Database Management Systems 3ed, R. Ramakrishnan and J. Gehrke



# ISA (`is a') Hierarchies

- Overlap constraints: Can Joe be an Hourly\_Emps as well as a Contract\_Emps entity? (Allowed/disallowed)
  - Hourly\_Emps OVERLAPS Senior\_Emps
- Covering constraints: Does every Employees entity also have to be an Hourly\_Emps or a Contract\_Emps entity? (Yes/no)
  - Hourly\_Emps AND Contract\_Emps COVERS Employees
- Reasons for using ISA:
  - To add descriptive attributes specific to a subclass.
  - To identify entitities that participate in a relationship.