

(February 6, 2024)

# Discriminants and Resultants: multiple and simultaneous zeros

Paul Garrett garrett@umn.edu <https://www-users.cse.umn.edu/~garrett/>

(All this goes back to mid-1800's, if not earlier!)

1. Discriminants and multiple zeros
2. Euclidean algorithm for  $\gcd(f, f')$
3.  $\gcd(f, f')$  is discriminant of  $f$
4. Resultants and common zeros

---

## 1. Discriminants and multiple zeros

For  $k$  a field, and polynomial  $f \in k[x]$ , the *discriminant* of  $f$  is (with no universal notation),

$$\prod_{i \neq j} (\theta_i - \theta_j) \quad (\theta_i \text{ the zeros of } f \text{ in } \bar{k})$$

where  $\bar{k}$  is an algebraic closure of  $k$ . The obvious intentional point is that this is 0 if and only if there is a repeated/multiple zero of  $f$ .

Because that expression is invariant under permutations of those zeros, it is expressible in terms of the elementary symmetric polynomials in the  $\theta$ 's, which are (up to signs) the coefficients of  $f$ . Beyond the quadratic case, it is tedious to execute the algorithm to obtain that expression of the discriminant. It is barely palatable in the cubic case.

---

## 2. Euclidean algorithm for $\gcd(f, f')$

At the same time,  $f$  has a multiple root/factor if and only if  $\gcd(f, f')$  is non-trivial, because any repeated factor of  $f$  will persist to  $f'$ . And conversely. In some extreme cases, it is feasible to formulaically describe the outcome of the Euclidean algorithm applied to  $f$  and  $f'$ . For example, for  $f(x) = x^n + ax + b$ :

$$f(x) - \frac{x}{n} \cdot f'(x) = (x^n + ax + b) - \frac{x}{n}(nx^{n-1} + a) = ax + b - \frac{x}{n}a = a\left(1 - \frac{1}{n}\right)x + b$$

Of course, we don't care about  $n = 1$ , and we decide now to not care about  $a = 0$  (which would be easy to appraise separately). Thus, this linear factor is essentially the same as

$$x + \frac{b}{a\left(1 - \frac{1}{n}\right)} = x - \frac{-b}{a\left(1 - \frac{1}{n}\right)}$$

From the Euclidean algorithm for polynomials over a field, we know that the remainder, upon dividing  $g(x)$  by  $x - \alpha$ , is  $g(\alpha)$ . Thus, the next step in this slightly larger Euclidean algorithm is

$$f'(x) - [?] \cdot \left(x - \frac{-b}{a\left(1 - \frac{1}{n}\right)}\right) = f'\left(\frac{-b}{a\left(1 - \frac{1}{n}\right)}\right)$$

where we do not care about the *dividend*. This is

$$n \cdot \left(\frac{-b}{a\left(1 - \frac{1}{n}\right)}\right)^{n-1} + a = (-1)^{n-1} \cdot \left(a\left(1 - \frac{1}{n}\right)\right)^{1-n} \cdot \left(nb^{n-1} + a \cdot \left(a\left(\frac{1}{n} - 1\right)\right)^{n-1}\right)$$

We can adjust by non-zero constants, to obtain

$$n^n b^{n-1} + (1-n)^{n-1} a^n$$

That is, the latter expression vanishes if and only if  $f$  has a repeated factor.

### 3. $\gcd(f, f')$ is discriminant of $f$

[3.1] **Claim:** For  $f(x) = x^n + ax + b$ , the expression  $n^n b^{n-1} + (1-n)^{n-1} a^n$  obtained above, by applying Euclidean algorithm to  $f$  and  $f'$ , is the discriminant of  $f$ .

*Proof:* The heuristic is about degree considerations, in terms of the zeros of  $f$  in an algebraic closure of  $k$ . Namely, on one hand,  $\prod_{i \neq j} (\theta_i - \theta_j)$  is apparently of degree  $n(n-1)$  in the zeros  $\theta_i$ . On the other hand,  $a = \pm s_{n-1}$  and  $b = s_n$ , the elementary symmetric polynomials in the zeros, which are of degrees  $n-1$  and  $n$ . Thus, the expression obtained via the Euclidean algorithm is apparently of degree  $(n-1)n$ , as well.

However, for one thing, if the  $\theta_i$  are merely *numbers* of some kind, or abstract field elements, this notion of *degree* does not have obvious content. This problem can be overcome by treating the *universal* version of the situation, namely, where  $k$  is the fraction field  $K(t_1, \dots, t_n)$  of a polynomial ring  $K[t_1, \dots, t_n]$ , and  $f \in k[x]$  has zeros  $t_i$ . The notion of (total) degree does make sense in  $K[t_1, \dots, t_n]$ , so we might want to consider the alleged identity in  $K[x, t_1, \dots, t_n]$ , even though we did the computation in a larger ring.

That is, in  $K[t_1, \dots, t_n]$ , indeed  $s_\ell$  is of (total) degree  $\ell$ . So  $a$  is indeed of degree  $n-1$ , and  $b$  of degree  $n$ , so  $a^n$  and  $b^{n-1}$  are both of degree  $n(n-1)$ , as the heuristic gives. And the product defining the discriminant, likewise, is of (total) degree  $n(n-1)$  in  $K[t_1, \dots, t_n]$ .

By unique factorization in polynomial rings over fields, since both expressions vanish (as polynomials in  $K[t_1, \dots, t_n]$ ) whenever any  $t_i$  and  $t_j$  are mapped to the same element of any target ring, both are divisible by all  $t_i - t_j$ . In both cases, by degree arguments, this does not leave any room for further factors of either.

### 4. Resultants and common zeros

For field  $k$  and  $f, g \in k[x]$ , the *resultant*  $R(f, g)$  of  $f$  and  $g$  is intended to be a polynomial (with coefficients in  $k$ ) in the coefficients of  $f$  and  $g$  whose vanishing is equivalent to  $f$  and  $g$  having simultaneous zeros. Thus, by the derivation criterion for repeated factors/roots, it should be that, the *discriminant* of a single polynomial  $f$  is the *resultant* of  $f$  and  $f'$ .

Letting  $\alpha_i$  and  $\beta_j$  be the zeros (with multiplicities) of  $f, g$  in an algebraic closure of  $k$ , up to constants, the resultant should be

$$R(f, g) = \prod_{i,j} (\alpha_i - \beta_j)$$

Since this  $R(f, g)$  is invariant under permutations of the  $\alpha_i$ , and under permutations of the  $\beta_j$ , by the theory of symmetric functions, it is a polynomial in the elementary symmetric polynomials in the  $\alpha_i$  and the  $\beta_j$ . Up to signs, these are the coefficients of  $f$  and  $g$ . This is *one* proof of the *existence* of the resultant.

However, the basic algorithm to express symmetric polynomials in terms of the elementary ones is qualitatively opaque, and, being completely general, ignores structural features of a given situation.

Another, more structured/intelligible approach: let  $f$  be of degree  $d$  and  $g$  of degree  $e$ . Let  $P_{<n}$  be the  $k$ -vectorspace of polynomials of degrees  $< n$ . The linear map

$$P_{<e} \oplus P_{<d} \longrightarrow P_{<e+d} \quad \text{by} \quad A \oplus B \longrightarrow Af + Bg$$

is a  $k$ -linear map from one  $(e + d)$ -dimensional space to another. It has non-zero kernel exactly when the determinant of the matrix giving the map, in whatever coordinates, is 0.

When the determinant is 0, there are non-zero polynomials  $A, B$ , of degrees less than those of  $g, f$  (in that order), such that  $Af + Bg = 0$ . That is,  $Af = -Bg$ . By unique factorization in  $k[x]$ , since the degree of  $B$  is strictly less than that of  $f$ , some factor of  $f$  must divide  $g$ . So, again, we have *existence* of a resultant, namely, that determinant.

That determinant can be *expressed* formulaically in terms of the natural basis for polynomials consisting of monomials  $x^i$ . Letting  $T : P_{<e} \oplus P_{<d} \rightarrow P_{<e+d}$  be that map, and  $f(x) = a_0 + a_1x + \dots + a_dx^d$ , and  $g(x) = b_0 + b_1x + \dots + a_dx^d$ ,

$$\begin{array}{rclcl}
 T(1 \oplus 0) & = & 1 \cdot f & = & a_0 + a_1x + \dots + a_dx^d \\
 T(x \oplus 0) & = & x \cdot f & = & a_0x + a_1x^2 + \dots + a_dx^{d+1} \\
 \dots & & & & \\
 T(x^{e-1} \oplus 0) & = & x^{e-1} \cdot f & = & a_0x^{e-1} + a_1x^e + \dots + a_dx^{e+d-1} \\
 T(0 \oplus 1) & = & 1 \cdot g & = & b_0 + b_1x + \dots + b_ex^e \\
 T(0 \oplus x) & = & x \cdot g & = & b_0x + b_1x^2 + \dots + b_ex^{e+1} \\
 \dots & & & & \\
 T(0 \oplus x^{d-1}) & = & x^{d-1} \cdot g & = & b_0x^{d-1} + b_1x^d + \dots + b_dx^{e+d-1}
 \end{array}$$

More later! :)

This does lead to a classic algebraic-curve fact, namely, Bézout's theorem, that two plane algebraic curves over  $\mathbb{C}$ , defined by polynomials  $f, g$ , intersect in  $(\deg f) \cdot (\deg g)$  points, counting multiplicities and points at infinity.

---