

## Lecture 30 Synopsis

(Fri, 6 May 2016)

In this lecture — the last lecture — I thought that it would be fun to return to Euler and the PNT.

At the very end, to keep things more current, I <sup>(relented and)</sup> interjected a quick note about  $S(T)$  [Lec 15, p. 26] which plays a role in ongoing computer tests of the Riemann Hypothesis.

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### EULER FIRST!

Recall

$$\mu(n) = \begin{cases} (-1)^r, & n = p_1 \cdots p_r, \text{ distinct } p_i \\ 0, & n \text{ not squarefree} \end{cases}$$

[ $\mu(1) = 1$ ]. See Lec 19 p. 14. We have:

$$\frac{1}{\zeta(s)} = \sum_{n=1}^{\infty} \frac{\mu(n)}{n^s}, \quad \operatorname{Re}(s) > 1$$

$$M(x) = \sum_{n \leq x} \mu(n), \quad x \geq 1, x \in \mathbb{R}.$$

Taking  $s > 1$  and letting  $s \rightarrow 1^+$ , one would suspect (with Euler, 1748) that

$$\sum_{n=1}^{\infty} \frac{\mu(n)}{n} = 0.$$

Of course, the "rub" here is that the convergence of the LHS is far from obvious!

NOTE. IF the LHS converges, its value must be 0. This follows immediately from Lec 21, p. 11 FACT 2(b) [with  $\delta = \frac{\pi}{2}$ ]; see also Lec 25 p. 3 on "super Stolz".

In Lec 20 p. 28 [D], I pointed out that it was known very early — using only ELEMENTARY techniques — that the following statements are equivalent:

recall  
Lec 2 p. 2

- (i)  $\psi(x) \sim x$  (i.e., the PNT) ;
- (ii)  $M(x) = o(x)$  ;
- (iii)  $\sum_{n=1}^{\infty} \frac{\mu(n)}{n} = 0$  .

See Lec 20 pp. 33 - 40 for (i)  $\Rightarrow$  (ii) ; then Lec 21

pp. ①-⑧ for (ii)  $\Rightarrow$  (i).

③

We now "finish" the job.

And bring in Euler.

THM.

$$(i)(ii) \iff (iii) \bullet$$

I.E., Euler  
"already knew"  
the PNT.

Proof

Following Landau's, Handbuch der Primzahlen,  
568 + 569, we set:

$$g(x) = \sum_{n \leq x} \frac{\mu(n)}{n}$$

$$x \geq 1$$

$$f(x) = \sum_{n \leq x} \frac{\mu(n) \ln n}{n} \bullet$$

Recall that  $|g(x)| < 1$  for  $x \geq 2$ ; Lec 20 p. ③7.

ASSUME (iii): I.E.,  $g(x) = o(1)$ .

Need to show  $M(x) = o(x)$ . Keep  $x \geq 2$ . Get:

$$\begin{aligned} \sum_{n \leq x} \mu(n) &= 1 + \int_1^x v dg(v) \\ &= 1 + [vg(v)]_1^x \sim \int_1^x g(v) dv \end{aligned}$$

$$= 1 + xg(x) - 1 - \int_1^x g(v) dv \quad (4)$$

$$= xg(x) - \int_1^x g(v) dv \cdot$$

But  $|g(t)| \leq \varepsilon$  for  $t \geq T_\varepsilon$ . For  $x \geq T_\varepsilon$ , we now get:

$$|M(x)| \leq \varepsilon x + \int_1^{T_\varepsilon} |g(v)| dv + \int_{T_\varepsilon}^x \varepsilon dv$$

$$\leq \varepsilon x + O(T_\varepsilon) + \varepsilon x$$

$$= 2\varepsilon x + O(T_\varepsilon)$$

$\Downarrow$

$$\limsup_{x \rightarrow \infty} \frac{|M(x)|}{x} \leq 2\varepsilon$$

$\Downarrow$

$$M(x) = o(x) \cdot$$

(OK)

NEXT, ASSUME (i)(ii). Must prove  $g(x) = o(1)$ .

To achieve this, we require 2 lemmas.

Lemma 1

$$g(x) - \frac{f(x)}{\ln x} = O\left(\frac{1}{\ln x}\right) \cdot$$

$$x \geq 2$$

(So:  $g(x) = o(1) \iff f(x) = o(\ln x)$ .)

Pf

Recall that  $\sum_{d|N} \mu(d) = \begin{cases} 1, & N=1 \\ 0, & N>1 \end{cases}$  (Lec 20 p. 29).

Then that

$$1 = \sum_{n \leq x} \mu(n) \left\lfloor \frac{x}{n} \right\rfloor$$

via  $\sum_{nk \leq x} \mu(n)$  and the hyperbola method à la

Lec 20 p. 37. This idea can be generalized

using  $k \equiv \frac{N}{n}$

$$1 = \sum_{nk \leq x} \frac{\mu(n)}{nk} \leftarrow \sum_{N \leq x} \frac{1}{N} \left\{ \sum_{n|N} \mu(n) \right\}$$

$$\Downarrow$$
$$1 = \sum_{n \leq x} \frac{\mu(n)}{n} \left\{ \sum_{k \leq \frac{x}{n}} \frac{1}{k} \right\}$$

$$= \sum_{n \leq x} \frac{\mu(n)}{n} \left[ \ln\left(\frac{x}{n}\right) + \gamma + O\left(\frac{n}{x}\right) \right]$$

↖ Lec 18 p. 40 bottom

(6)

$$= \sum_{n \leq x} \frac{\mu(n)}{n} (\ln x - \ln n)$$

$$+ \gamma \sum_{n \leq x} \frac{\mu(n)}{n}$$

$$+ O(1) \sum_{n \leq x} \frac{1}{n} \frac{n}{x}$$

$$\Downarrow$$

$$1 = (\ln x) g(x) - \sum_{n \leq x} \frac{\mu(n) \ln n}{n} + \gamma g(x) + O(1)$$

$$|g(x)| \leq 1$$

$$\Downarrow$$

$$1 = (\ln x) g(x) - f(x) + O(1) \quad \text{by (3)}$$

$$\Downarrow$$

$$(\ln x) g(x) - f(x) = O(1)$$

$$\Downarrow$$

$$g(x) - \frac{f(x)}{\ln x} = O\left(\frac{1}{\ln x}\right) \quad \square$$

Must now seek to prove  $f(x) = o(\ln x)$  .

Lemma 2

For  $x \geq 2$ , we have:

$$\sum_{n \leq x} \frac{1}{n} g\left(\frac{x}{n}\right) = 1.$$

Pf

On (5) middle, we saw that

$$1 = \sum_{nk \leq x} \frac{\mu(n)}{nk}$$

think  
hyperbolas

{sum the other direction!}

$$= \sum_{k \leq x} \frac{1}{k} \left\{ \sum_{n \leq \frac{x}{k}} \frac{\mu(n)}{n} \right\}$$

$$= \sum_{k \leq x} \frac{1}{k} g\left(\frac{x}{k}\right) \quad (\text{by def of } g).$$



Recall that

$$-\mu(N) \ln N = \sum_{kl=N} \lambda(k) \mu(l)$$

by Lec 20 p. 33. Accordingly

$$f(x) = \sum_{N \leq x} \frac{\mu(N) \ln N}{N} \quad \text{ala } \textcircled{3}$$

$$= - \sum_{kl \leq x} \frac{\lambda(k) \mu(l)}{kl}$$

{ think hyperbolas }

$$= - \sum_{k \leq x} \frac{\lambda(k)}{k} \left\{ \sum_{l \leq \frac{x}{k}} \frac{\mu(l)}{l} \right\}$$

$$f(x) \approx - \sum_{k \leq x} \frac{\lambda(k)}{k} g\left(\frac{x}{k}\right)$$

recall  
 $\lambda(1) = 0$   
 $\lambda(p^m) = \ln p, m \geq 1$

Integration by parts with  $\psi(x)$  is cumbersome since  $g(*)$  is not continuous. Avoid it!!

To prove  $f(x) = o(\ln x)$  as  $x \rightarrow \infty$ , there is no loss of generality in taking  $x = \text{integer}$ . See (3).

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We have

$$\psi(v) = \sum_{n \leq v} 1(n)$$

and  $\psi(v) = 0, v < 2$ . Since we have assumed (i)(ii), it makes sense to write

$$\psi(v) = v + v\varepsilon(v), v > 0$$

with

$$\lim_{v \rightarrow \infty} \varepsilon(v) = 0$$

and

$$\varepsilon(v) \equiv -1 \text{ for } 0 < v < 2.$$

For convenience, declare:

$$\varepsilon(0) = 0.$$

We still have  $\psi(0) = 0 + 0\varepsilon(0)$ .

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CHOOSE  $\eta$  so that

$$|\varepsilon(v)| \leq \eta, \text{ all } v \geq 0.$$

In (8) box, notice that  $(x = \text{integer})$

$$\sum_{k \leq x} \frac{1(k)}{k} g\left(\frac{x}{k}\right)$$

$$= \sum_{k \leq x} \frac{\psi(k) - \psi(k-1)}{k} g\left(\frac{x}{k}\right)$$

$$\left\{ \psi(k) \approx k + k \varepsilon(k) \right\}$$

$$= \sum_{1 \leq k \leq x} \frac{1 + k \varepsilon(k) - (k-1) \varepsilon(k-1)}{k} g\left(\frac{x}{k}\right)$$

$$\approx \sum_{k=1}^x \frac{1}{k} g\left(\frac{x}{k}\right)$$

$$+ \sum_{k=1}^x [\varepsilon(k) - \varepsilon(k-1)] g\left(\frac{x}{k}\right)$$

$$+ \sum_{k=1}^x \frac{\varepsilon(k-1)}{k} g\left(\frac{x}{k}\right)$$

$\approx 1$  ← by Lemma 2 on (7)

$$+ \sum_{k=1}^x \varepsilon(k) g\left(\frac{x}{k}\right) - \sum_{k=2}^x \varepsilon(k-1) g\left(\frac{x}{k}\right)$$

$$+ \sum_{k=2}^x \frac{\varepsilon(k-1)}{k} g\left(\frac{x}{k}\right)$$

$\varepsilon(0) = 0$

$$= 1 + \sum_{k=1}^x \varepsilon(k) g\left(\frac{x}{k}\right) - \sum_{\ell=1}^{x-1} \varepsilon(\ell) g\left(\frac{x}{\ell+1}\right) \\ + \sum_{\ell=1}^{x-1} \frac{\varepsilon(\ell)}{\ell+1} g\left(\frac{x}{\ell+1}\right)$$

$$= 1 + \sum_{k=1}^x \varepsilon(k) g\left(\frac{x}{k}\right) - \sum_{\ell=1}^{\overset{x}{\circlearrowleft}} \varepsilon(\ell) g\left(\frac{x}{\ell+1}\right) \\ + \sum_{\ell=1}^{\overset{x}{\circlearrowleft}} \frac{\varepsilon(\ell)}{\ell+1} g\left(\frac{x}{\ell+1}\right)$$

IF WE DECLARE  
 $g(v) \equiv 0$  for  $v < 1$

$$= 1 + \sum_{k=1}^x \varepsilon(k) \left[ g\left(\frac{x}{k}\right) - g\left(\frac{x}{k+1}\right) \right] \\ + \sum_{\ell=1}^x \frac{\varepsilon(\ell)}{\ell+1} g\left(\frac{x}{\ell+1}\right)$$

$$(-f(x)) = 1 + \sum_1 + \sum_2, \text{ say.}$$

Fix any tiny  $\varepsilon > 0$ . Let  $|\varepsilon(k)| \leq \varepsilon$  for  
 all  $k \geq K_\varepsilon + 1$ . Keep  $x \geq K_\varepsilon + 100$ .

Notice that:

$$|\Sigma_2| \leq \sum_{l=1}^{K_\epsilon} \frac{g_l}{l+1} \cdot 1$$

$$|g(v)| \leq 1$$

$$+ \sum_{K_\epsilon+1 \leq l \leq x} \frac{\epsilon}{l+1} \cdot 1$$

$$\leq C_\epsilon + \epsilon \ln x + O(1)$$

$$\leq C'_\epsilon + \epsilon \ln x \quad \text{//}$$

To estimate  $\Sigma_1$ , notice that

$$g\left(\frac{x}{k}\right) - g\left(\frac{x}{k+1}\right) = \sum_{\frac{x}{k+1} < n \leq \frac{x}{k}} \frac{u(n)}{n}$$

by p. ③ and ⑪ lines 5+6. Accordingly,

$$\left| g\left(\frac{x}{k}\right) - g\left(\frac{x}{k+1}\right) \right| \leq \left\{ \sum_{\left(\frac{x}{k+1}, \frac{x}{k}\right]} \frac{1}{n} \right\}^2$$

In the foregoing,  $\mathbb{Z} \cap (\frac{x}{k+1}, \frac{x}{k}]$  may well be empty. Empty sums are zero.

Continue to keep  $x \geq K_\epsilon + 100$ . By def,

$$\Sigma_1 = \sum_{k=1}^x \epsilon(k) \left[ g\left(\frac{x}{k}\right) - g\left(\frac{x}{k+1}\right) \right]. \quad (1)$$

Hence:

$$|\Sigma_1| \leq \sum_{1 \leq k \leq K_\epsilon} M \cdot 2 + \sum_{K_\epsilon+1 \leq k \leq x} \epsilon \left( \sum_{(\frac{x}{k+1}, \frac{x}{k}]} \frac{1}{n} \right)$$

||| THE UNION OVER k WITH  $(\frac{x}{k+1}, \frac{x}{k}]$  IS DISJOINT |||

$$\leq 2M K_\epsilon + \epsilon \sum_{n \leq \frac{x}{K_\epsilon+1}} \frac{1}{n}$$

$$\leq 2M K_\varepsilon + \varepsilon (\ln x + O(1)) \quad (14)$$

↓

$$|\Sigma_1| \leq C_\varepsilon'' + \varepsilon \ln x \quad \blacksquare$$

Upon reviewing (8) box, (10), (11), (12) (middle), and line 2 above, we see that

$$|F(x)| \leq 1 + C_\varepsilon''' + 2\varepsilon \ln x$$

for  $x \geq K_\varepsilon + 100$ . Since  $\varepsilon > 0$  was arbitrary on (11), we get  $f(x) = o(\ln x)$ , as required on (6) bottom. By Lemma 1, on (5), we therefore have  $g(x) = o(1)$ , as needed.

This completes the proof of THM on (3).  $\blacksquare$

(I couldn't resist)

Finally, a very brief comment about  $\zeta(T)$ !

We recall that:

vis à vis verifications of RH

$$\zeta(T) = \frac{1}{\pi} \operatorname{Arg} \zeta\left(\frac{1}{2} + iT\right)$$

'ala Lec 15 (26) (middle). (cf. also here Lec 15 pp. (22) - (25) and Lec 27 (3) (FACT) - (6) (top).

When  $T \neq$  all  $y$ , we have

$$\zeta(T) = -\frac{1}{\pi} \operatorname{Im} \int_{\frac{1}{2}}^{\infty} \frac{\zeta'(s+iT)}{\zeta(s+iT)} ds$$

If  $T =$  some  $y$ , one defines  $\zeta(T)$  by right continuity to preserve Lec 15 (26) (lines 6+7).

One knows that:

$$\zeta(T) = O(\ln T)$$

THM

Define  $\text{Log } \zeta(s)$  in the usual up and across way beginning at point  $A \in \mathbb{R}$ ,  $A \gg 1$ .  
 Keep  $T \geq 2$ ,  $T \neq$  all  $\gamma$ , and  $-1 \leq \sigma \leq 2$ .

We then have:

$$\text{Log } \zeta(s) = O(\ln T) + \sum_{\substack{\rho \\ |\gamma - T| \leq 1}} \text{Log}(s - \rho) \cdot$$

Here  $s = \sigma + iT$  and  $\text{Log} =$  the standard principal value.

Pf

Review Lec 27 (3) (bot) - (5) and Lec 15 (22) (bot).

$$\begin{aligned} \text{Log } \zeta(s) &= 0 + \int_{\infty}^{\sigma} \frac{\zeta'(u+iT)}{\zeta(u+iT)} du \\ &= \int_{\infty}^2 O(2^{-u}) du + \int_2^{\sigma} \frac{\zeta'(u+iT)}{\zeta(u+iT)} du \\ &= O(1) + \int_2^{\sigma} \left[ O(\ln T) + \sum_{\substack{\rho \\ |\gamma - T| \leq 1}} \frac{1}{u+iT-\rho} \right] du \end{aligned}$$

↖  
 Lec 15 (13) (8)

$$= O(1) + O(\ln T)$$

$$+ \sum_{\substack{\rho \\ |y-T| \leq 1}} [\text{Log}(s-\rho) - \text{Log}(2+iT-\rho)]$$

$$= O(\ln T) + \sum_{\substack{\rho \\ |y-T| \leq 1}} \text{Log}(s-\rho)$$

{ since  $0 \leq \text{Re}(\rho) \leq 1$  } • 

For  $T \neq$  all  $\gamma$ , it is customary to define:

$$\zeta_1(T) = -\frac{1}{\pi} \text{Re} \int_{\frac{1}{2}}^{\infty} (s-\frac{1}{2}) \frac{\zeta'(s+iT)}{\zeta(s+iT)} ds$$

For  $T =$  some  $\gamma$ , use right continuity (cf. page (18) box).

It is easy to see that  $\zeta_1(T)$  is well-defined and nicely continuous insofar as  $T$  remains away from all  $\gamma$ . Lec 15, pp. (22) (bot) - (23) (top).

Notice here that:

$$\begin{aligned} S_1 &= -\frac{1}{\pi} \operatorname{Re} \int_{\frac{1}{2}}^{\infty} (\sigma - \frac{1}{2}) d \operatorname{Log} I(\sigma + iT) \\ &= -\frac{1}{\pi} \operatorname{Re} \left[ 0 - 0 - \int_{\frac{1}{2}}^{\infty} \operatorname{Log} I(\sigma + iT) d\sigma \right] \end{aligned}$$

$\left. \begin{array}{l} \text{by integ by parts and} \\ \operatorname{Log} I(\sigma) = O(2^{-\sigma}), \sigma \geq 2 \end{array} \right\}$

$$\Downarrow$$

$$S_1(T) = \frac{1}{\pi} \int_{\frac{1}{2}}^{\infty} \ln |I(\sigma + iT)| d\sigma$$

This is for  $T \neq$  all  $\gamma$ . By baby analysis and

$$\int_0^1 |\ln u| du < \infty,$$

the box remains true for all  $T \geq 2$  (i.e., there is good continuous behavior).

Notice further (for  $T \neq \text{all } \gamma$ ) that:

$$S_1(T) = \frac{1}{\pi} \operatorname{Re} \int_{\frac{1}{2}}^{\infty} \operatorname{Log} \zeta(\sigma + iT) d\sigma \quad (18) \text{ line 3}$$

$\swarrow$   
 $O(2^{-\sigma}) \quad \sigma \geq 2$

$$\Rightarrow \underline{S_1'(T)} = \frac{1}{\pi} \operatorname{Re} \int_{\frac{1}{2}}^{\infty} \frac{\zeta'}{\zeta}(\sigma + iT) i d\sigma \quad \left\{ \begin{array}{l} \text{Leibnitz's} \\ \text{rule} \end{array} \right\}$$

$$= -\frac{1}{\pi} \operatorname{Im} \int_{\frac{1}{2}}^{\infty} \frac{\zeta'}{\zeta}(\sigma + iT) d\sigma$$

$$= S(T) \text{ by (15) middle}$$

{ with good uniform convergence }  
locally wrt  $T$



$$S_1'(T) = S(T), \quad T \neq \text{all } \gamma$$

$$S_1(T) = c_1 + \int_2^T S(u) du, \quad \text{all } T \geq 2$$

continuous wrt  $T$

# THM (very fundamental)

In the zero-counting formula

$$N(T) = \underbrace{\frac{T}{2\pi} \ln\left(\frac{T}{2\pi e}\right) + \frac{7}{8} + O\left(\frac{1}{T}\right)}_{C^{\infty}} + S(T)$$

ala Lec 15 (22) + (26), we have

$$S(T) = O(\ln T), \quad \underline{\underline{\int_2^T S(u) du = O(\ln T)}}.$$

Pf.

Only the last assertion remains to be proved.  
WLOG  $T \neq$  all  $\gamma$ . Apply (18) box. Get:

$$\begin{aligned} J_1(T) &= \frac{1}{\pi} \int_2^{\infty} \ln |I(\sigma + iT)| d\sigma + \frac{1}{\pi} \int_{\frac{1}{2}}^2 \ln |I| d\sigma \\ &\quad \uparrow O(2^{-\sigma}) \\ &= O(1) + \frac{1}{\pi} \int_{\frac{1}{2}}^2 \ln |I(u + iT)| du \end{aligned}$$

(21)

$$= O(1) + \frac{1}{\pi} \int_{\frac{1}{2}}^2 [O(\ln T) + \sum_{\substack{\rho \\ |1-\gamma| \leq 1}} \ln |u+iT-\rho|] du$$

by p. (16) THM  
 observe too that we have  
 $|u-\rho| \leq |(u-\rho)+i(T-\gamma)| \leq 3$

$$= O(1) + O(\ln T) + O(\ln T)$$

$$= O(\ln T) \quad \blacksquare$$

It is hardly surprising that the implied constants for  $\ln T$  in p. (20) THM can be made explicit.

The relation  $\int_2^T \zeta(u) du = O(\ln T)$  qualitatively states that the average value of  $\zeta(u)$  is 0.

These last 2 points are important. In the early 1950s, Alan Turing used these facts to develop a numerical criterion (now known as Turing's Law) by which the Riemann Hypothesis on interval  $[T_1, T_2]$  can be checked simply by locating enough (that is to say, a requisite number of) sign-changes for the REAL-VALUED function

$$\zeta\left(\frac{1}{2} + it\right) \quad \text{or, better,} \quad \frac{\zeta\left(\frac{1}{2} + it\right)}{|\zeta\left(\frac{1}{2} + it\right)|}$$

in an interval slightly bigger than  $[T_1, T_2]$ . \*

recall Lec 23, p. ③ lines 1-3 and (e); Lec 15, p. ②⑥ (line -3)

OF COURSE, there is nothing to guarantee in advance that the requisite number will be found. One simply has to try!!!

The point here is 4-fold:

- (A) if the requisite number is reached (by the machine), one is assured by Turing's theorem that all zeros are accounted for, and that there are none having  $\text{Re}(s) \neq \frac{1}{2}$ ;
- (B) there is no need to check RH for  $T < T_1$  first;

\* with special attention paid to the pattern near  $T_1, T_2$

(C) there is no need to compute any  $J$ -values with  $\text{Re}(s) \neq \frac{1}{2}$ ;

(D) there is no need to "zero in" on the crossings in  $[T_1, T_2]$  attached to the sign-changes detected by the machine.

To understand why Turing's Law is at least believable, simply pretend that one somehow knew that  $S(t)$  was exactly zero in some short intervals centered at  $T_1$  and  $T_2$ . See p. (20) THEOREM [the formula for  $N(T)$ ] and ponder the logical consequences which ensue!

For further details about Turing and  $J(s)$ , see the paper of Hejhal and Odlyzko in the Turing Centenary volume "Alan Turing: His Work and Impact" published by Elsevier. The story is a VERY interesting one. ← with links to lec 29 p. (29) and S. Skewes

Turing's Law is used in all modern computational work aimed at verifying the RH. When the approach is successful, the zeros of  $J$  in the range  $[T_1, T_2]$  are also known to be simple. Cf. (A) on (22).