

CS 2011: Practice Quiz 3 SOLUTION

Summer 2018

University of Minnesota

Quiz period: 15 minutes

Points available: 40

Use the following table of letter/number correspondences for encryption problems.

-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----													
	0		1		2		3		4		5		6
	A		B		C		D		E		F		G
	-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----												
	13		14		15		16		17		18		19
	N		O		P		Q		R		S		T
	-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----												
	20		21		22		23		24		25		
	U		V		W		X		Y		Z		
	-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----												

Problem 1 (10 pts): Encrypt the following message using the Caesar cipher with the key provided.

- Message: ZEBRA
- Key: 8
- **Show some work or full credit.**
- Give the letters of the final encrypted message.

SOLUTION:

- $Z \equiv 25, 25 + 8 \pmod{26} = 33 \pmod{26} = 7 \equiv H$
- $E \equiv 4, 4 + 8 \pmod{26} = 12 \pmod{26} = 12 \equiv M$
- $B \equiv 1, 1 + 8 \pmod{26} = 9 \pmod{26} = 9 \equiv J$
- $R \equiv 17, 17 + 8 \pmod{26} = 25 \pmod{26} = 25 \equiv Z$
- $A \equiv 0, 0 + 8 \pmod{26} = 8 \pmod{26} = 8 \equiv I$

Encrypted message: HMJZI

Problem 2 (10 pts): Arrange the following functions in order of big-O growth rate right with slowest growing functions listed first.

SOLUTION:

- | | |
|--------------------------|----------------------------|
| (A) $2n^3 \log n$ | • (D) $7\sqrt{n} + \log n$ |
| (B) $5n(\log n)^2$ | • (B) $5n(\log n)^2$ |
| (C) $4n + n!$ | • (A) $2n^3 \log n$ |
| (D) $7\sqrt{n} + \log n$ | • (E) $3^n + n \log n$ |
| (E) $3^n + n \log n$ | • (C) $4n + n!$ |

Problem 3 (10 pts): Use a proof by induction to show that for integer $n > 4$, $2^n > n^2$. Clearly illustrate Base and Inductive Cases and identify the Inductive Hypothesis.

SOLUTION:

BASE CASE: For $n = 5$, $2^5 = 32$ and $5^2 = 25$. Since $32 > 25$ the property holds.

INDUCTION CASE:

1. Assume that for k , the inequality $2^k > k^2$ holds.
2. Show for $k + 1$, the inequality $2^{k+1} > (k + 1)^2$ holds.
3. Starting with the left side of the inequality transform as follows.

$$\begin{aligned}
 2^{k+1} &= 2 \cdot 2^k && \text{Def of Exponentiation} \\
 &> 2 \cdot k^2 && \text{By Inductive hypothesis (1)} \\
 &> k^2 + 4k && \text{Because } k \geq 5 \\
 &> k^2 + 2k + 1 && \text{Because } k \geq 5 \\
 &= (k + 1)^2 && \text{Factor polynomial}
 \end{aligned}$$

4. Have shown that left hand side of inequality is indeed larger than right hand side so induction case holds.

Problem 4 (10 pts): The algorithm `complete(A)` determines if array **A** of length N contains exactly the numbers $0, 1, 2, \dots, N - 1$ in some order in which case **true** is returned. Otherwise, if some number is missing, **false** is returned.

Analyze the code below and give a big-O estimate of its worst-case runtime.

```

1: complete(A[] : int array){
2:   for(n=0; n<length(A); n++){
3:     found = false
4:     i = 0
5:     while(i < length(A)
6:       and found == false)
7:     {
8:       if(A[i] == n){
9:         found = true
10:      }
11:      i++;
12:    }
13:    if(found == false){
14:      return false;
15:    }
16:  }
17:  return true;
18: }
```

SOLUTION: The outer loop from line 2 to 16 iterates N times where N is the length of the array. The operations at lines 3 and 4 are constant time. The inner loop from lines 5 to 12 also runs for N iterations and each operation within the loop is constant time. These leads to the combined complexity of the loops being $O(N^2)$. In the worst case the algorithm will check for all values $0, 1, \dots, (N-1)$. This means the overall runtime is $O(N^2)$.