CSCI 2021: Introduction

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Last Updated:
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CSCI 2021 - Logistics

Reading
- Bryant/O’Hallaron: Ch 1
- C references: basic syntax, types, compilation

Goals
- Basic Model of Computation
- Begin discussion of C
- Course Mechanics

Ongoing
Due Tue 9/14 11:59pm
- Lab01: Setup, submit to Gradescope
- HW01: Basics, online Gradescope Quiz
“Von Kauffman” Model: CPU, Memory, Screen, Program

Most computers have 4 basic, physical components:

1. A CPU which can execute instructions
2. CPU knows WHICH instruction to execute at all times
3. MEMORY where data is stored and can change
4. Some sort of Input/Output device like a SCREEN

The CPU is executes a set of instructions, usually called a program, which change MEMORY and the SCREEN

Example of a Running Computer Program

<table>
<thead>
<tr>
<th>CPU: at instruction 10:</th>
<th>MEMORY:</th>
<th>SCREEN:</th>
</tr>
</thead>
<tbody>
<tr>
<td>&gt; 10: set #1024 to 1801</td>
<td>Addr</td>
<td>Value</td>
</tr>
<tr>
<td>11: set #1028 to 220</td>
<td>------+-------</td>
<td></td>
</tr>
<tr>
<td>12: sum #1024,#1028 into #1032</td>
<td>#1032</td>
<td>-137</td>
</tr>
<tr>
<td>13: print #1024, &quot;plus&quot;, #1028</td>
<td>#1028</td>
<td>12</td>
</tr>
<tr>
<td>14: print &quot;is&quot;, #1032</td>
<td>#1024</td>
<td>19</td>
</tr>
</tbody>
</table>

1Of course it’s a little more complex than this but the addage, “All models are wrong but some are useful.” applies here. This class is about asking “what is really happening?” and going deep down the resulting rabbit hole.
Sample Run Part 1

CPU: at instruction 10:
> 10: set #1024 to 1801
  11: set #1028 to 220
  12: sum #1024,#1028 into #1032
  13: print #1024, "plus", #1028
  14: print "is", #1032

CPU: at instruction 11:
  10: set #1024 to 1801
  > 11: set #1028 to 220
  12: sum #1024,#1028 into #1032
  13: print #1024, "plus", #1028
  14: print "is", #1032

CPU: at instruction 12:
  10: set #1024 to 1801
  11: set #1028 to 220
  > 12: sum #1024,#1028 into #1032
  13: print #1024, "plus", #1028
  14: print "is", #1032
Sample Run Part 2

CPU: at instruction 13:

10: set #1024 to 1801 | Addr | Value |
11: set #1028 to 220 |-------+-------|
12: sum #1024,#1028 into #1032 | #1032 | 2021 |
> 13: print #1024, "plus", #1028 | #1028 | 220 |
14: print "is", #1032 | #1024 | 1801 |

CPU: at instruction 14:

10: set #1024 to 1801 | Addr | Value |
11: set #1028 to 220 |-------+-------|
12: sum #1024,#1028 into #1032 | #1032 | 2021 |
13: print #1024, "plus", #1028 | #1028 | 220 |
> 14: print "is", #1032 | #1024 | 1801 |

CPU: at instruction 15:

10: set #1024 to 1801 | Addr | Value |
11: set #1028 to 220 |-------+-------|
12: sum #1024,#1028 into #1032 | #1032 | 2021 |
13: print #1024, "plus", #1028 | #1028 | 220 |
14: print "is", #1032 | #1024 | 1801 |
> 15: ....
Observations: CPU and Program Instructions

- Program instructions are usually small, simple operations:
  - Put something in a specific memory cell using its **address**
  - Copy the contents of one cell to another
  - Do arithmetic (add, subtract, multiply, divide) with numbers in cells and specified constants like 5
  - Print stuff to the screen

- The CPU keeps track of which instruction to execute next

- In many cases after executing it moves ahead by one instruction but you all know **jumping** around is also possible

- This program is in **pseudocode**: not C or Java or Assembly…

- Pseudocode can have almost anything in it so long as a human reader understands the meaning

- Real machines require more precise languages to execute as they are (still) much dumber than humans
Observations: Screen and Memory

Screen versus Memory

- Nothing is on the screen until it is explicitly printed by the program
- Normally you don’t get to see memory while the program runs
- Good programmers can quickly form a mental picture of what memory looks like and draw it when needed
- You will draw memory diagrams in this class to develop such mental models

Memory Cells

- Memory cells have Fixed ADDRESS Changeable CONTENTS
- Random Access Memory (RAM): the value in any memory cell can be retrieved FAST using its address
- My laptop has 16GB of memory = 4,294,967,296 (4 billion) integer boxes (!)
  - Cell Address #’s never change: always cell #1024
  - Cell Contents frequently change: set #1024 to 42
Dealing with raw memory addresses is tedious

Any programming language worth its salt will have **variables**: symbolic names associated with memory cells

**You pick variable names**; compiler/interpreter automatically translates to memory cell/address

---

### PROGRAM ADDRESSES ONLY

**CPU**: at instruction 50:

- > 50: copy #1024 to #1032
- 51: copy #1028 to #1024
- 52: copy #1032 to #1028
- 53: print "first", #1024
- 54: print "second", #1028

| Addr | Value |
|-------+-------|
| #1024 | 42 |
| #1028 | 31 |
| #1032 | ? |

**MEMORY**:

| Addr | Value |
|-------+-------|
| #1024 | 42 |
| #1028 | 31 |
| #1032 | ? |

### PROGRAM WITH NAMED CELLS

**CPU**: at instruction 51:

- > 50: copy x to temp
- 51: copy y to x
- 52: copy temp to y
- 53: print "first", x
- 54: print "second", y

| Addr | Name  | Value |
|-------+-------+-------|
| #1032 | temp  | ?     |
| #1028 | y     | 31    |
| #1024 | x     | 42    |

**MEMORY**:

| Addr | Name  | Value |
|-------+-------+-------|
| #1032 | temp  | ?     |
| #1028 | y     | 31    |
| #1024 | x     | 42    |

---
Correspondence of C Programs to Memory

- C programs require memory cell names to be declared with the **type of data** they will hold (*a novel idea when C was invented*).

- The equal sign (=) means
  
  “store the result on the right in the cell named on the left”

- Creating a cell and giving it a value can be combined
  
  ```
  int x; // need a cell named x, holds an integer
  x = 42; // put 42 in cell x
  int y = 31; // need a cell named y and put 31 in it
  int tmp = x + y; // cell named tmp, fill with sum of x and y
  ```

Other Rules

- C/Java compilers read whole programs to figure out how many memory cells are needed based on declarations like `int a;` and `int c=20;`

- Lines that only declare a variable do nothing except indicate a cell is needed to the compiler

- In C, uninitialized variables may have arbitrary crud in them making them dangerous to use: *we’ll find out why in this course*
Exercise: First C Snippet

- Lines starting with // are comments, ignored
- `printf("%d %d\n",x,y)` shows variable values on the screen

```
50: int x;
51: x = 42;
52: int y = 31;
53: // swap x and y (?)
54: x = y;
55: y = x;
56: printf("%d %d\n",x,y);
```

CPU: at line 50
MEMORY: | Addr | Name | Value |
--- | --- | --- | --- |
> 50: int x; | | | |
51: x = 42; | | | |
52: int y = 31; | #1032 | y | ? |
53: // swap x and y (?) | #1028 | x | ? |
54: x = y; | #1024 | | |
55: y = x; |
56: printf("%d %d\n",x,y);

SCREEN:

With your nearby Room colleagues:

1. Show what memory / screen look like after running the program
2. Correct the program if needed: make swapping work

I will chat with a couple folks about their answers which will earn them an Engagement Point.
Answer: First C Snippet

CPU: at line 54

```c
50: int x;
51: x = 42;
52: int y = 31;
53: // swap x and y (?)
54: x = y;  // swap x and y (?)
55: y = x;
56: printf("%d %d\n",x,y);
```

CPU: at line 55

```c
50: int x;
51: x = 42;
52: int y = 31;
53: // swap x and y (?)
54: x = y;  // swap x and y (?)
55: y = x;
56: printf("%d %d\n",x,y);
```

CPU: at line 57

```c
50: int x;
51: x = 42;
52: int y = 31;
53: // swap x and y (?)
54: x = y;  // swap x and y (?)
55: y = x;
56: printf("%d %d\n",x,y);
```

Clearly **incorrect**: how does one swap values properly? (fix `swap_main_bad.c`)
First Full C Program: swap_main.c

1 /* First C program which only has a main(). Demonstrates proper
2 swapping of two int variables declared in main() using a third
3 temporary variable. Uses printf() to print results to the screen
4 (standard out). Compile run with:
5
6 > gcc swap_main.c
7 > ./a.out
8 */

9 #include <stdio.h> // headers declare existence of functions
10 // printf in this case
11 int main(int argc, char *argv[]){ // ENTRY POINT: always start in main()
12     int x; // declare a variable to hold an integer
13     x = 42; // set its value to 42
14     int y = 31; // declare and set a variable
15     int tmp = x; // declare and set to same value as x
16     x = y; // put y's value in x's cell
17     y = tmp; // put tmp's value in y's cell
18     printf("%d %d\n",x,y); // print the values of x and y
19     return 0; // return from main(): 0 indicates success
20 }

▶ Swaps variables using tmp space (exotic alternatives exist)
▶ Executables always have a main() function: starting point
▶ Note inclusion of stdio.h header to declare printf() exists, allusions to C's (limited and clunky) library system
# Exercise: Functions in C, swap_func.c

```c
1 // C program which attempts to swap using a function.
2 //
3 // > gcc swap_func.c
4 // > ./a.out
5
6 #include <stdio.h> // declare existence printf()
7 void swap(int a, int b); // function exists, defined below main
8
9 int main(int argc, char *argv[]){ // ENTRY POINT: start executing in main()
10   int x = 42;
11   int y = 31;
12   swap(x, y); // invoke function to swap x/y (?)
13   printf("%d %d\n",x,y); // print the values of x and y
14   return 0;
15 }
16
17 // Function to swap (?) contents of two memory cells
18 void swap(int a, int b){ // arguments to swap
19   int tmp = a; // use a temporary to save a
20   a = b; // a <- b
21   b = tmp; // b <- tmp=a
22   return;
23 }
```

Does `swap()` “work”? Discuss with neighbors and justify why the code works or why not
Answers: The Function Call Stack and \texttt{swap()}

9: int main(...){

STACK: Caller \texttt{main()}, prior to \texttt{swap()}

<table>
<thead>
<tr>
<th>FRAME</th>
<th>ADDR</th>
<th>SYM</th>
<th>VALUE</th>
</tr>
</thead>
</table>

10: int x = 42;

11: int y = 31;

+-<12: \texttt{swap(x, y)};

<table>
<thead>
<tr>
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<th>VALUE</th>
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</thead>
</table>

V 15: }

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</table>

+-<18: void \texttt{swap(int a, int b)}{

19: int tmp = a;

20: a = b;

21: b = tmp;

22: return;

23: }

STACK: Callee \texttt{swap()} takes control

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\begin{itemize}
  \item \textbf{Caller function} \texttt{main()} and \textbf{Callee function} \texttt{swap()}
  \item Caller \textbf{pushes} a stack frame onto the \textbf{function call stack}
  \item Frame has space for all Callee parameters/locals
  \item Caller tracks where it left off to resume later
  \item Caller copies values to Callee frame for parameters
  \item Callee begins executing at its first instruction
\end{itemize}
**Answers: Function Call Stack: Returning from swap()**

9: int main(...){
10: int x = 42;
11: int y = 31;
12: swap(x, y);
  
+->13: printf("%d %d\n",x,y);
| 14: return 0;
| 15: }

18: void swap(int a, int b){
 | 19: int tmp = a;
 | 20: a = b;
 | 21: b = tmp;

+<22: return;

23: }

STACK: Callee swap() returning

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<th>ADDR</th>
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</tr>
</thead>
<tbody>
<tr>
<td>main()</td>
<td>#2048</td>
<td>x</td>
<td>42</td>
</tr>
<tr>
<td>line:12</td>
<td>#2044</td>
<td>y</td>
<td>31</td>
</tr>
</tbody>
</table>

STACK: Caller main() gets control back

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▶ On finishing, Callee stack frame **pops** off, returns control back to Caller which resumes executing next instruction

▶ Callee may pass a return value to Caller but otherwise needs proper setup to alter the Caller stack frame.

▶ **swap()** does NOT swap the variables in **main()**
Motivation for C

Pure Abstraction

Python, JS
Ruby, Shell
C++, D
Assembly
C
VHDL
Binary
Opcodes

If this were Java, Python, many others, discussion would be over:

- Provide many safety and convenience features
- Insulate programmer from hardware for ease of use

C presents many CPU capabilities directly

- Very few safety features
- Little between programmer and hardware

You just have to know C. Why?
Because for all practical purposes, every computer in the world you’ll ever use is a von Neumann machine, and C is a lightweight, expressive syntax for the von Neumann machine’s capabilities.

–Steve Yegge, Tour de Babel
Von Neumann Machine / Architecture (Wikipedia)

Processing

- Wires/gates that accomplish fundamental ops
- +, -, *, AND, OR, move, copy, shift, etc.
- Ops act on contents of memory cells to change them

Control

- Memory address of next instruction to execute
- After executing, move ahead one unless instruction was to jump elsewhere

Memory

- Giant array of bits/bytes so everything is represented as 1’s and 0’s, including instructions
- Memory cells accessible by address number

Input/Output

- Allows humans to interpret what is happening
- Often special memory locations for screen and keyboard

Wait, these items seem kind of familiar...
Exercise: C allows direct use of memory cell addresses
Where/how are these used in the code below?

<table>
<thead>
<tr>
<th>&amp;x</th>
<th>memory address of variable x</th>
</tr>
</thead>
<tbody>
<tr>
<td>int *a</td>
<td>a stores a memory address (pointer to integer(s))</td>
</tr>
<tr>
<td>*a</td>
<td>get/set the memory pointed to by a (dereference)</td>
</tr>
</tbody>
</table>

```c
#include <stdio.h> // declare existence printf()

void swap_ptr(int *a, int *b); // function exists, defined below main

int main(int argc, char *argv[]){ // ENTRY POINT: start executing in main()
    int x = 42;
    int y = 31;
    swap_ptr(&x, &y); // call swap() with addresses of x/y
    printf("%d %d\n",x,y); // print the values of x and y
    return 0;
}

// Function to swap contents of two memory cells
void swap_ptr(int *a, int *b){ // a/b are addresses of memory cells
    int tmp = *a; // go to address a, copy value int tmp
    *a = *b; // copy val at addr in b to addr in a
    *b = tmp; // copy temp into address in b
    return;
}
```
Swapping with Pointers/Addresses: Call Stack

9: int main(...){
10:   int x = 42;
11:   int y = 31;
+-<12:  swap_ptr(&x, &y);
|  13:  printf("%d %d\n",x,y);
|  14:  return 0;
V 15: }

STACK: Caller main(), prior to swap()
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STACK: Callee swap() takes control
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<td>y</td>
<td>31</td>
</tr>
<tr>
<td>swap_ptr</td>
<td>#2036</td>
<td>a</td>
<td>#2048</td>
</tr>
<tr>
<td>line:19</td>
<td>#2028</td>
<td>b</td>
<td>#2044</td>
</tr>
</tbody>
</table>

- Syntax &x reads "Address of cell associated with x" or just "Address of x". Ampersand & is the address-of operator.
- Swap takes int *a: **pointer** to integer / memory address
- Values associated with a/b are the addresses of other cells
Swapping with Pointers/Addresses: Dereference/Use

9: int main(...){
10:    int x = 42;
11:    int y = 31;
12:    swap_ptr(&x, &y);
13:    printf("%d %d\n",x,y);
14:    return 0;
15: }
18: void swap_ptr(int *a,int *b){
19:    int tmp = *a; // copy val at #2048 to #2032
20:    *a = *b;
21:    *b = tmp;
22:    return;
23: }

LINE 19 executed: tmp gets 42

Syntax *a reads “Dereference a to operate on the cell pointed to by a” or just “Deref a”

Line 19 dereferences via * operator:

- Cell #2040 (a) contains address #2048,
- Copy contents of #2048 (42) into #2032 (tmp)
Aside: Star/Asterisk \* has **3 uses** in C

1. Multiply as in
   \[ w = c*d; \]

2. **Declare** a pointer as in
   ```c
   int *x;  // pointer to integer(s)
   int b=4;
   x = &b;  // point x at b
   int **r; // pointer to int pointer(s)
   ```

3. **Dereference** a pointer as in
   ```c
   int p = *x; // x must be an int pointer
               // retrieve contents at address
   ```

Three different context sensitive meanings for the same symbol makes \* hard on humans to parse, a BAD move by K&R.

```c
int z = *x * *y + *(p+2); // standard, 'unambiguous' C
```

The duck is ready to eat. // English is more ambiguous
Dereference can be used to get values at an address

Can be used on left-hand-side of assignment to set contents at an address

Line 20: dereference a to change contents at #2048
Swapping with Pointers/Addresses: Deref 2

9: int main(...){
10:  int x = 42;
11:  int y = 31;
12:  swap_ptr(&x, &y);
13:  printf("%d %d\n",x,y);
14:  return 0;
15: }
18: void swap_ptr(int *a,int *b){
19:  int tmp = *a;
20:  *a = *b;
21:  *b = tmp;  // copy val at #2032 (42) to #2044 (was 31)
>22:  return;
23: }

LINE 21 executed: alters y using b
| FRAME | ADDR | NAME | VALUE |
|-------+-------+------+-------|
| main() | #2048 | x | 31 |<-+|<--|
| line:12 | #2044 | y | 31->42 |<-|+|

▶ Can be used on left-hand-side of assignment to set contents at an address
▶ Line 21: dereference *b = ... to change contents at #2044
▶ Use of variable name tmp retrieves contents of cell associated with tmp
Swapping with Pointers/Addresses: Returning

9: int main(...){
10: int x = 42;
11: int y = 31;
12: swap_ptr(&x, &y);
13: printf("%d %d\n",x,y);
14: return 0;
15: }
18: void swap_ptr(int *a, int *b){
19: int tmp = *a;
20: *a = *b;
21: *b = tmp;
22: return;
23: }

LINE 22: prior to return
<table>
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<th>VALUE</th>
</tr>
</thead>
</table>
| main() | #2048 | x    | 31    |<--+
| line:12 | #2044 | y    | 42    |<--|+
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
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<td>swap_ptr</td>
<td>#2036</td>
<td>a</td>
<td>#2048</td>
</tr>
</tbody>
</table>
| line:22 | #2028 | b    | #2044 |----+
| | #2024 | tmp  | 42    |

LINE 12 finished/return pops frame
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<td>#2044</td>
<td>y</td>
<td>42</td>
</tr>
<tr>
<td>-----------------------</td>
<td>--------</td>
<td>------</td>
<td>-------</td>
</tr>
</tbody>
</table>

.swap_ptr() finished so frame pops off

Variables x, y in main() have changed due to use of references to them.
Important Principle: Non-local Changes

- Pointers allow functions to change variables associated with other running functions.

- Common beginner example: `scanf()` family which is used to read values from terminal or files.

- Snippet from `scanf_demo.c`:
  ```c
  1 int main(...){
  2   int num = -1;
  3   scanf("%d", &num); // addr
  4   printf("%d\n",num); // val
  4   return 0;
  5 }
  ```

- See `scanf_error.c`: forgetting & yields great badness.

```
<table>
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<th>FRAME</th>
<th>ADDR</th>
<th>NAME</th>
<th>VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>main():3</td>
<td>#2500</td>
<td>num</td>
<td>-1</td>
</tr>
</tbody>
</table>
  |-------|------|------|-------|---+
  | scanf() | #2492 | fmt | #400 | |
  |       | #2484 | arg1 | #2500 |---+

  scanf() called
```

```
<table>
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</tr>
</thead>
<tbody>
<tr>
<td>main():3</td>
<td>#2500</td>
<td>num</td>
<td>5</td>
</tr>
</tbody>
</table>
  |-------|------|------|-------|---+
  | scanf() | #2492 | fmt | #400 | |
  |       | #2484 | arg1 | #2500 |---+

  scanf() changes contents of #2500
```

```
<table>
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<tr>
<th>FRAME</th>
<th>ADDR</th>
<th>NAME</th>
<th>VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>main():4</td>
<td>#2500</td>
<td>num</td>
<td>5</td>
</tr>
</tbody>
</table>
  |-------|------|------|-------|---+
```

```
  scanf() returns
```
Uncle Ben Said it Best...

- Pointers allow any line of C programs to modify any of its data
- A BLESSING: fine control of memory → efficiency, machine’s true capability
- A CURSE: opens up many errors not possible in langs like Java/Python which restrict use of memory

1972 - Dennis Ritchie invents a powerful gun that shoots both forward and backward simultaneously. Not satisfied with the number of deaths and permanent maimings from that invention he invents C and Unix.

– A Brief, Incomplete, and Mostly Wrong History of Programming Languages
Beneath the C

C is a “high-level” as it abstracts away from a real machine. It must be translated to lower levels to be executed.

Assembly Language

➤ Specific to each CPU architecture (Intel, etc)
➤ Still “human readable” but fairly directly translated to binary using Assemblers

<table>
<thead>
<tr>
<th>INTEL x86-64 ASSEMBLY</th>
<th>HEXADECIMAL/BINARY OPCODES</th>
</tr>
</thead>
<tbody>
<tr>
<td>cmpl $1, %ecx</td>
<td>1124: 83 f9 01</td>
</tr>
<tr>
<td>jle .END</td>
<td>1127: 7e 1e = 0111 1110 0001 1110</td>
</tr>
<tr>
<td>movl $2, %esi</td>
<td>1129: be 02 00 00 00</td>
</tr>
<tr>
<td>movl %ecx,%eax</td>
<td>112e: 89 c8</td>
</tr>
<tr>
<td>cqto</td>
<td>1130: 48 99</td>
</tr>
<tr>
<td>idivl %esi</td>
<td>1132: f7 fe</td>
</tr>
<tr>
<td>cmpl $1,%edx</td>
<td>1134: 83 fa 01</td>
</tr>
<tr>
<td>jne .EVEN</td>
<td>1137: 75 07</td>
</tr>
</tbody>
</table>

Binary Opcodes

➤ 1’s and 0’s, represent the digital signal of the machine
➤ Codes corresponds to instructions directly understood by processor

Looks like **fun**, right? You bet it is! Assembly coding is 1 month away…
CSCI 2021: Course Goals

▶ Basic proficiency at C programming
▶ Knowledge of running programs in physical memory including the stack, heap, global, and text areas of memory
▶ Understanding of the essential elements of assembly languages
▶ Knowledge of the correspondence between high-level program constructs.
▶ Ability to use a symbolic debugger
▶ Basic understanding of how data is encoded in binary
▶ Knowledge of computer memory systems
▶ Basic knowledge of computer architecture
A Word on Safety

Please wear your mask during lecture

► For your safety, my safety, the safety of the class, and the safety of all the old, young, and immuno-compromised loved ones that we see but do not want to hurt, mask up.

► Refrain from eating/drinking during lecture

► Keep your mask on the whole time

► If you feel sick, stay home, watch the videos, notify me if the illness is prolonged and we will make arrangements