CSCI 2021: Virtual Memory

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Last Updated:
Mon Apr 25 01:06:15 PM CDT 2022
Logistics

Reading Bryant/O’Hallaron

- Ch 9: Virtual Memory
- Ch 7: Linking (next)

P5

- 1 Problem
- Implement a small version of malloc() / free()
- Post later today with video, due last day of class

Goals

- Address Spaces, Translation, Paged Memory
- mmap(), Sharing Pages

<table>
<thead>
<tr>
<th>Date</th>
<th>Event</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fri 4/22</td>
<td>Virtual Mem 1/2</td>
</tr>
<tr>
<td>Mon 4/25</td>
<td>Virtual Mem 2/2</td>
</tr>
<tr>
<td>Wed 4/27</td>
<td>ELF Files/Linking 1/2</td>
</tr>
<tr>
<td></td>
<td>Lab 14 mmap()</td>
</tr>
<tr>
<td></td>
<td>HW 14 Linking</td>
</tr>
<tr>
<td>Fri 4/29</td>
<td>Obj Code/Linking 2/2</td>
</tr>
<tr>
<td>Mon 5/2</td>
<td>Last Lecture, Review P5 Due</td>
</tr>
</tbody>
</table>
Exercise: The View of Memory Addresses so Far

- Every **process** (running program) has some memory, divided into roughly 4 areas (which are...?)
- Reference different data/variables through their addresses
- If only a single program could run at time, no trouble: load program into memory and go
- Running multiple programs gets interesting particularly if they both reference the *same memory location*, e.g. address 1024

```
PROGRAM 1                        PROGRAM 2
...                              ...
## load global from #1024        ## add to global at #1024
movq 1024, %rax                 addl %esi, 1024
...                              ...
```

- What **conflict** exists between these programs?
- What are possible **solutions** to this conflict?
Answers: The View of Memory Addresses so Far

- 4 areas of memory are roughly: (1) Stack (2) Heap (3) Globals (4) Text/Instructions
- Both programs use physical address #1024, behavior depends on order that instructions are interleaved between them

ORDER A: Program 1 loads first

<table>
<thead>
<tr>
<th>PROGRAM 1</th>
<th>PROGRAM 2</th>
</tr>
</thead>
<tbody>
<tr>
<td>movq 1024, %rax</td>
<td>...</td>
</tr>
<tr>
<td>...</td>
<td>addl %esi, 1024</td>
</tr>
</tbody>
</table>

ORDER B: Program 2 adds first

<table>
<thead>
<tr>
<th>PROGRAM 1</th>
<th>PROGRAM 2</th>
</tr>
</thead>
<tbody>
<tr>
<td>...</td>
<td>addl %esi, 1024</td>
</tr>
<tr>
<td>movq 1024, %rax</td>
<td>...</td>
</tr>
</tbody>
</table>

- **Solution 1:** Never let Programs 1 and 2 run together (bleck!)
- **Solution 2:** Translate every memory address in every program on **loading** it, run with physical addresses
  - Tough/impossible as not all addresses are known at compile/load time...
- **Solution 3:** Translate every memory address/access in every program while it runs (!!!)
Paged Memory

- Physical memory is divided into hunks called **pages**
- Common page size supported by many OS’s (Linux) and hardware is 4KB = 4096 bytes, can be larger with OS config
- CPU models use some # of bits for **Virtual Addresses**
  ```
  > cat /proc/cpuinfo
  vendor_id : GenuineIntel
  cpu family : 6
  model : 79
  model name : Intel(R) Xeon(R) CPU E5-1620 v4 @ 3.50GHz
  ...
  address sizes : 46 bits physical, 48 bits virtual
  ```
- Example of address with page number and offset labelled
  ```
  xxxxPagenumberOff : 48 bits used
  0x00007ffa0997a428 : 64 bit address
  |   |   |   |
  |   |   --> Offset 0x428 within page, 12 bits
  |   | +--> Page number 0x7ffa0997a, 36 bits
  | +--> Constant bits, not used by processor
  ```
Translation happens at the Page Level

- Within a page, addresses are sequential
- Between pages, may be non-sequential

### Page Table:

<table>
<thead>
<tr>
<th>Virtual Page</th>
<th>Size</th>
<th>Physical Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>00007ffa0997a000</td>
<td>4K</td>
<td>RAM: 0000564955aa1000</td>
</tr>
<tr>
<td>00007ffa0997b000</td>
<td>4K</td>
<td>RAM: 0000321e46937000</td>
</tr>
<tr>
<td>...</td>
<td></td>
<td>...</td>
</tr>
</tbody>
</table>

### Address Space From Page Table:

<table>
<thead>
<tr>
<th>Virtual Address</th>
<th>Page Offset</th>
<th>Physical Address</th>
</tr>
</thead>
<tbody>
<tr>
<td>00007ffa0997a000</td>
<td>0</td>
<td>0000564955aa1000</td>
</tr>
<tr>
<td>00007ffa0997a001</td>
<td>1</td>
<td>0000564955aa1001</td>
</tr>
<tr>
<td>00007ffa0997a002</td>
<td>2</td>
<td>0000564955aa1002</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td>00007ffa0997afff</td>
<td>4095</td>
<td>0000564955aa1fff</td>
</tr>
</tbody>
</table>

| 00007ffa0997b000 | 0           | 0000321e46937000 |
| 00007ffa0997b001 | 1           | 0000321e46937001 |
| ...             | ...         | ...              |
Addresses Translation Hardware

- Translation must be **FAST** so usually involves hardware
- **MMU (Memory Manager Unit)** is a hardware element specifically designed for address translation
- Usually contains a special cache, **TLB** (Translation Lookaside Buffer), which stores recently translated addresses
- **OS Kernel interacts with MMU**
- Provides location of the **Page Table**, data structure relating Virtual/Physical Addresses
- **Page Fault**: MMU couldn’t map Virtual to Physical page, runs a Kernel routine to handle the fault
Exercise: Translating Virtual Addresses

Nearby diagram illustrates relation of Virtual Pages to Physical Pages

1. How many page tables are there?
2. Where can a page table entry refer to?
3. Count the number of Virtual pages, compare to the number of physical pages - which is larger?
4. What happens if PID #123 accesses its Virtual Page #2
5. What happens if PID #456 accesses its Virtual Page #2
On using a Virtual Memory address, MMU will search TLB for physical DRAM address.

If found in TLB, Hit, use physical DRAM address.

If not found, MMU will search Page Table, if found and in DRAM, cache in TLB.

Else Miss = Page fault, OS decides..

1. Page is swapped to Disk, move to DRAM, potentially evicting another page
2. Page not in page table = Segmentation Fault
Each process has its own page table, OS maintains mapping of Virtual to Physical addresses.

Processes “compete” for RAM.

OS gives each process impression it owns all of RAM.

OS may not have enough memory to back up all or even 1 process.

Disk used to supplement ram as Swap Space.

Thrashing may occur when too many processes want too much RAM, “constantly swapping”
Trade-offs of Address Translation

Wins of Virtual Memory

1. Avoids processes each referencing the same address, conflicting
2. Allows each Process (running program) to believe it has entire memory to itself
3. Gives OS tons of flexibility and control over memory layout
   - Present a continuous Virtual chunk which is spread out in Physical memory
   - Use Disk Space as memory
   - Check for out of bounds memory references

Losses of Virtual Memory

1. Address translation is not constant O(1), has an impact on performance of real algorithms*
2. Requires special hardware to make translation fast enough: MMU/TLB
3. Not needed if only a single program is running on a machine

Wins outweigh Losses in most systems so Virtual Memory is used widely, a great idea in CS

*See On a Model of Virtual Address Translation (2015)
The Many Other Advantages of Virtual Memory

1. Caching: Seen that VirtMem can treat main memory as a cache for larger memory
2. Security: Translation allows OS to check memory addresses for validity, segfault on out-of-bounds access
3. Debugging: Valgrind checks addresses for validity
4. Sharing Data: Processes can share data with one another; request OS to map virtual addresses to same physical addresses
5. **Sharing Libraries**: Can share same program text between programs by mapping address space to same shared library
6. **Convenient I/O**: Map internal OS data structures for files to virtual addresses to make working with files free of `read()`/`write()`
Virtual Memory and `mmap()`

- Normally programs interact indirectly with Virtual Memory system
  - Stack/Heap/Globals/Text are mapped automatically to regions in Virtual Memory System
  - Maps are adjusted as Stack/Heap Grow/Shrink
- `mmap()` / `munmap()` directly manipulate page tables
  - `mmap()` creates new entries in page table
  - `munmap()` deletes entries in the page table
  - Can map arbitrary or specific addresses into memory
- `mmap()` is used to initially set up Stack / Heap / Globals / Text when a program is loaded by the program loader
- While a program is running can also use `mmap()` to interact with virtual memory
- A convenient way to do File I/O via Memory Mapped Files
Exercise: Printing Contents of file

Examine the two programs below which print the contents of a file

▶ Identify differences between them
▶ Which has a higher memory requirement?

1 // print_file.c
2 int main(int argc, char *argv[]){
3     FILE *fin = fopen(argv[1], "r");
4     char inbuf[256];
5     while(1){
6         int nread =
7             fread(inbuf, sizeof(char),
8                 256, fin);
9         if(nread == 0){
10             break;
11         }
12         fwrite(inbuf,sizeof(char),
13                 nread,stdout);
14     }
15     fclose(fin);
16     return 0;
17 }

1 // mmap_print_file.c
1 int main(int argc, char *argv[]){
2     int fd = open(argv[1], O_RDONLY);
3         struct stat stat_buf;
4         fstat(fd, &stat_buf);
5         int size = stat_buf.st_size;
6         char *file_chars =
7             mmap(NULL, size,
8                 PROT_READ, MAP_SHARED,
9                 fd, 0);
10     for(int i=0; i<size; i++){
11         printf("%c",file_chars[i]);
12     }
13     printf("\n");
14     munmap(file_chars, size);
15     close(fd);
16     return 0;
17 }
Answers: Printing Contents of file

1. Write a simple program to print all characters in a file. What are key features of this program?
   - Open file
   - Read up to 256 characters into memory using `fread()`/`fscanf()`
   - Print those characters with `printf()`
   - Read more characters and print
   - Stop when end of file is reached
   - Close file

2. Examine ` mmap_print_file.c`: does it contain all of these key features? Which ones are missing?
   - Missing the `fread()`/`fscanf()` portion
   - Uses `mmap()` to get **direct access** to the bytes of the file
   - Treat bytes as an array of characters and print them directly
mmap(): Mapping Addresses is Amazing

- ptr = mmap(NULL, size,...,fd,0) arranges backing entity of fd to be mapped to be mapped to ptr
- fd often a file opened with open() system call

```c
int fd = open("gettysburg.txt", O_RDONLY);
// open file to get file descriptor

char *file_chars = mmap(NULL, size, PROT_READ, MAP_SHARED,
                         fd, 0);
// call mmap to get a direct pointer to the bytes in file associated // with fd; NULL indicates don't care what address is returned; // specify file size, read only, allow sharing, offset 0

printf("%c",file_chars[0]); // print 0th file char
printf("%c",file_chars[5]); // print 5th file char
```
For efficiency, part of files are stored in RAM by the OS.

OS manages internal data structures to track which parts of a file are in RAM, whether they need to be written to disk.

`mmap()` alters a process Page Table to translate addresses to the cached file page.

OS tracks whether page is changed, either by file write or `mmap()` manipulation.

Automatically writes back to disk when needed.

Changes by one process to cached file page will be seen by other processes.

See diagram on next slide.
Diagram of Kernel Structures for `mmap()`

**USER SPACE**

Process Memory #1234

<table>
<thead>
<tr>
<th>read_fd</th>
<th>3</th>
</tr>
</thead>
<tbody>
<tr>
<td>file_bytes</td>
<td>#4096</td>
</tr>
</tbody>
</table>

```c
int read_fd = open(...);
char *file_bytes = mmap(..., read_fd);
```

Shows some internal structures of the kernel that allow `mmap()` to be used to map a user pointer `file_bytes` to the same page of memory where a file is cached by the OS.

**KERNEL SPACE**

Process #1234 Page Table / File Descriptors

<table>
<thead>
<tr>
<th>#0</th>
<th>#4096</th>
<th>#12288</th>
</tr>
</thead>
<tbody>
<tr>
<td>/</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**file table entry**

- **fd**
  - `RD_ONLY | pos=0`
  - size=16K
  - cached pages
  - disk blocks

**inode for file**

<table>
<thead>
<tr>
<th>RAM PAGES</th>
<th>PHYSICAL STORAGE</th>
</tr>
</thead>
<tbody>
<tr>
<td>ABCD</td>
<td>ABCD</td>
</tr>
<tr>
<td>EFGH</td>
<td>EFGH</td>
</tr>
<tr>
<td></td>
<td>HIJK</td>
</tr>
<tr>
<td></td>
<td>LMNO</td>
</tr>
</tbody>
</table>
mmap() exposes several capabilities from the OS

```
char *file_chars =
    mmap(NULL, size,
        PROT_READ | PROT_WRITE, // map allowing read + write
        MAP_SHARED, // share changes with original file
        fd, 0); // file to map + offset from start
```

Assign new value to memory, OS writes changes into the file

**Example**: `mmap_tr.c` to transform one character to another
Mapping things that aren’t characters

\texttt{mmap()} just gives a pointer: can assert type of what it points at

- Example \texttt{int *}: treat file as array of binary ints
- Notice changing array will write to file

\begin{verbatim}
// mmap_increment.c: demonstrate working with mmap()'d binary data

int fd = open("binary_nums.dat", O_RDWR);
// open file descriptor, like a FILE *

int *file_ints = mmap(NULL, size, PROT_READ | PROT_WRITE, MAP_SHARED, fd, 0);
// get pointer to file bytes through mmap,
// treat as array of binary ints

int len = size / sizeof(int);
// how many ints in file

for(int i=0; i<len; i++){
    printf("%d\n",file_ints[i]); // print all ints
}

for(int i=0; i<len; i++){
    file_ints[i] += 1; // increment each file int, writes back to disk
}
\end{verbatim}
mmap() Compared to Traditional fread()/fwrite() I/O

Advantages of mmap()

▶ Avoid following cycle
  ▶ fread()/fscanf() file contents into memory
  ▶ Analyze/Change data
  ▶ fwrite()/fscanf() write memory back into file
▶ Saves memory and time
▶ Many Linux mechanisms backed by mmap() like processes sharing memory

Drawbacks of mmap()

▶ Always maps pages of memory: multiple of 4096b (4K)
▶ For small maps, lots of wasted space
▶ Cannot change size of files with mmap(): must used fwrite() to extend or other calls to shrink
▶ No bounds checking, just like everything else in C
Page Table Size

- Page tables map a virtual page to physical location
- Page tables maintained by operating system in Kernel Memory
- A direct page table has one entry per virtual page
- Each page is $4K = 2^{12}$ bytes, so 12 bits for offset of address into a page
- Virtual Address Space is $2^{48}$ bytes
- So, $2^{36}$ virtual pages mapped in the page table...
  - 68,719,476,736 pages
  - At 8 bytes per page entry...
  - 1 Terabyte for a page table

How big does the page table mapping virtual to physical pages need to be?
“Tables” are Multi-Level Sparse Trees

Fix this absurdity with **multi-level page tables**: a sparse tree

- Virtual address divided into sections which indicate which PTE to access at different table levels
- 3-4 level page table is common in modern architectures
- Programs typically use only small amounts of virtual memory: most entries in different levels are NULL (not mapped) leading to much smaller page tables than a direct (array) map
Direct Page Table vs Sparse Tree Page Table

**Direct Page Table: Array-Like, 5-bit addresses**

<table>
<thead>
<tr>
<th>VP#</th>
<th>Valid</th>
<th>PP#</th>
<th>Contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>0000</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>00001</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>00010</td>
<td>1</td>
<td>01001</td>
<td>654</td>
</tr>
<tr>
<td>00011</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>00100</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>00101</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>00110</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>00111</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>01000</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>01001</td>
<td>0</td>
<td>/</td>
<td>987</td>
</tr>
<tr>
<td>...</td>
<td>0</td>
<td>/</td>
<td>...</td>
</tr>
<tr>
<td>11011</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>11000</td>
<td>1</td>
<td>00001</td>
<td>321</td>
</tr>
<tr>
<td>11010</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>11100</td>
<td>1</td>
<td>11100</td>
<td>-</td>
</tr>
<tr>
<td>11101</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>11110</td>
<td>1</td>
<td>11100</td>
<td>-</td>
</tr>
<tr>
<td>11111</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
</tbody>
</table>

**Physical Memory**

<table>
<thead>
<tr>
<th>PP#</th>
<th>Contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>0000</td>
<td>-</td>
</tr>
<tr>
<td>00001</td>
<td>654</td>
</tr>
<tr>
<td>00010</td>
<td>-</td>
</tr>
<tr>
<td>00011</td>
<td>-</td>
</tr>
<tr>
<td>00100</td>
<td>-</td>
</tr>
<tr>
<td>00101</td>
<td>-</td>
</tr>
<tr>
<td>00110</td>
<td>-</td>
</tr>
<tr>
<td>00111</td>
<td>-</td>
</tr>
<tr>
<td>01000</td>
<td>-</td>
</tr>
<tr>
<td>01001</td>
<td>987</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td>11011</td>
<td>-</td>
</tr>
<tr>
<td>11000</td>
<td>321</td>
</tr>
<tr>
<td>11010</td>
<td>-</td>
</tr>
<tr>
<td>11100</td>
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</tr>
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</tr>
<tr>
<td>11110</td>
<td>-</td>
</tr>
<tr>
<td>11111</td>
<td>-</td>
</tr>
</tbody>
</table>

**Two-level Page Table: Sparse Tree, 5-bit addresses**

<table>
<thead>
<tr>
<th>VP High Bits</th>
<th>Valid Node</th>
<th>VP Low Bits</th>
<th>Valid</th>
<th>PP#</th>
<th>Contents</th>
</tr>
</thead>
<tbody>
<tr>
<td>000</td>
<td>1</td>
<td>00</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>001</td>
<td>0</td>
<td>/</td>
<td>0</td>
<td>/</td>
<td>-</td>
</tr>
<tr>
<td>010</td>
<td>0</td>
<td>/</td>
<td>10</td>
<td>01001</td>
<td>-</td>
</tr>
<tr>
<td>011</td>
<td>0</td>
<td>/</td>
<td>11</td>
<td>0</td>
<td>-</td>
</tr>
<tr>
<td>100</td>
<td>0</td>
<td>/</td>
<td>00</td>
<td>1</td>
<td>00001</td>
</tr>
<tr>
<td>101</td>
<td>0</td>
<td>/</td>
<td>01</td>
<td>0</td>
<td>-</td>
</tr>
<tr>
<td>110</td>
<td>1</td>
<td>00001</td>
<td>10</td>
<td>1</td>
<td>11000</td>
</tr>
<tr>
<td>111</td>
<td>1</td>
<td>11100</td>
<td>11</td>
<td>0</td>
<td>-</td>
</tr>
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**Physical Memory**

<table>
<thead>
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<td>654</td>
</tr>
<tr>
<td>00010</td>
<td>-</td>
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<tr>
<td>00011</td>
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<td>-</td>
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<tr>
<td>01000</td>
<td>-</td>
</tr>
<tr>
<td>01001</td>
<td>987</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
</tr>
<tr>
<td>11100</td>
<td>-</td>
</tr>
<tr>
<td>11101</td>
<td>-</td>
</tr>
<tr>
<td>11110</td>
<td>-</td>
</tr>
<tr>
<td>11111</td>
<td>-</td>
</tr>
</tbody>
</table>

Both data structures map 3 virtual pages to 3 physical pages as indicated in the map to the left but use different amounts of space to do so.

**Direct Table**
- 3 pages mapped
- 32 entries required

**Mult-level Table**
- 3 pages mapped
- 16 entries required
- 50% space saved

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Textbook Example: Two-level Page Table

Space savings gained via NULL portions of the page table/tree

32 bit addresses, 4KB pages, 4-byte PTEs

Source: Bryant/O’Hallaron, CSAPP 3rd Ed
Pages and Mapping

- Memory is segmented into hunks called **pages**, 4Kb is common (use page-size.c to see your system’s page size)
- OS maintains tables of which pages of memory exist in RAM, which are on disk
- OS maintains tables per process that translate process virtual addresses to physical pages
- **Shared Memory** can be arranged by mapping virtual addresses for two processes to the same memory page

![Diagram showing virtual addresses and physical memory]

**Virtual Address Space for Process 1:**

- VP 1
- VP 2
- ⋮
- N-1

**Virtual Address Space for Process 2:**

- VP 1
- VP 2
- ⋮
- N-1

**Address translation**

**Physical Address Space (DRAM):**

- PP 2
- PP 6
- PP 8
- ⋮
- M-1

(e.g., read-only library code)
Shared Libraries: *.so Files

- Code for libraries can be shared
- libc.so: shared library with malloc(), printf() etc in it
- OS puts into one page, maps all linked procs to it

Source: John T. Bell Operating Systems Course Notes
pmap: show virtual address space of running process

> ./memory_parts
0x5575555a71e9 : main()
0x5575555aa0c0 : global_arr
0x5575555b482a0 : heap_arr
0x600000000000 : mmap'd block1
0x600000001000 : mmap'd block2
0x7f2244dc4000 : mmap'd file
0x7ffff0133b70 : stack_arr
my pid is 496605
press any key to continue

▶ Determine **process id**
   of running program

▶ **pmap** reports its virtual
   address space

▶ More details of **pmap**
   output in this article
   from Andreas Fester

▶ His diagram is awesome

    ▶
    ▶ pmap 496605
    ▶

    > pmap 496605
    496605: ./memory_parts
    00005575555a6000  4K r----- memory_parts
    00005575555a7000  4K r-x-- memory_parts TEXT
    00005575555a8000  4K r----- memory_parts
    00005575555a9000  4K r----- memory_parts
    00005575555aa000  4K rw--- memory_parts GLOBALS
    00005575555ab000  4K rw--- [ anon ]
    00005575555ac000  132K rw--- [ anon ]  HEAP
    0000600000000000  8K rw--- [ anon ]
    00007f2244bca000  8K rw--- [ anon ]
    00007f2244bcc000  152K r----- libc-2.32.so
    00007f2244bf2000  1332K r-x-- libc-2.32.so
    00007f2244d3f000  304K r----- libc-2.32.so
    00007f2244d8e000  12K rw--- libc-2.32.so
    00007f2244d91000  24K rw--- [ anon ]
    00007f2244dc4000  4K r----- gettysburg.txt
    00007f2244dc5000  8K r----- ld-2.32.so
    00007f2244dc7000  132K r-x-- ld-2.32.so
    00007f2244de8000  36K r----- ld-2.32.so
    00007f2244df2000  8K rw--- ld-2.32.so
    00007ffff0114000  132K rw--- [ stack ] STACK
    00007ffff014d000  12K r----- [ anon ]
    total            2352K

fanatical
Memory Protection

- Output of `pmap` indicates another feature of virtual memory: **protection**
- OS marks pages of memory with Read/Write/Execute/Share permissions like files
- Attempt to violate these and get segmentation violations (segfault)
- Ex: Executable page (instructions) usually marked as `r-x`: no write permission.
- Ensures program don’t accidentally write over their instructions and change them
- Ex: By default, pages are not shared (no 's' permission) but can make it so with the right calls
Exercise: Quick Review

1. While running a program, memory address #1024 always refers to a physical location in DRAM (True/False: why?)
2. Two programs which both use the address #1024 cannot be simultaneously run (True/False: why?)
3. What do MMU and TLB stand for and what do they do?
4. What is a memory page? How big is it usually?
5. What is a Page Table and what is it good for?
Answers: Quick Review

1. While running a program, memory address #1024 always refers to a physical location in DRAM (True/False: why?)
   ▶ False: #1024 is usually a **virtual address** which is translated by the OS/Hardware to a physical location which *may* be in DRAM but may instead be paged out to disk

2. Two programs which both use the address #1024 cannot be simultaneously run (True/False: why?)
   ▶ False: The OS/Hardware will likely translate these identical virtual addresses to **different physical locations** so that the programs do not clobber each other’s data

3. What do MMU and TLB stand for and what do they do?
   ▶ Memory Management Unit: a piece of hardware involved in translating Virtual Addresses to Physical Addresses/Locations
   ▶ Translation Lookaside Buffer: a special cache used by the MMU to make address translation **fast**

4. What is a memory page? How big is it usually?
   ▶ A discrete hunk of memory usually 4Kb (4096 bytes) big

5. What is a Page Table and what is it good for?
   ▶ A table maintained by the operating system that is used to map Virtual Addresses to Physical addresses for each page
Additional Review Questions

- What OS data structure facilitates the Virtual Memory system? What kind of data structure is it?
- What does pmap do?
- What does the mmap() system call do that enables easier I/O? How does this look in a C program?
- Describe at least 3 benefits a Virtual Memory system provides to a computing system