MPI Basics

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Last Updated: Tue Feb 7 01:45:04 PM CST 2023

Logistics

Reading: Grama Ch 6 + 4

- Ch 6: MPI basics
- Ch 4: Communication patterns

Assignments

- A1 Due Soon
 - On-time by Thu 02-Feb
 - Late through Sat 04-Feb

Questions?

 A2 up next week: MPI Programming

Today

- Primitives for Distributed Memory Computing
- MPI Programming

Next Week

- Comm Patterns
- Thu 09-Feb: Mini-Exam 1

Generic Send and Receive

Distributed memory machines require explicit sharing of data

Minimum required functionality is:

send(void *sendbuf, int nelems, int dest) receive(void *recvbuf, int nelems, int source)

Referred to as a "point-to-point" communication

Sample Use

- // PO runs // P1 runs 1 $(0 \rightarrow 0)$ a = 100;
- send(&a, 1, 1); 3
- a=50; 4

- Proc 0 sends a single integer to Proc 1
- Proc 0 then changes that integer
- Proc 1 receives and prints the integer

More typical appearance

- Typically write this as a single program which every processor runs: Single Program, Multiple Data (SPMD)
- Assume availability of a function giving logical Proc Number
- Branching on proc number to take different actions

```
1 void exchange(){
2
    int a = 100:
    int my proc = get processor number();
3
    if(my_proc == 0){
4
      send(&a, 1, 1); // send data 100 to Proc 1
5
      a=50:
6
7
    }
    else if(my proc == 1){
8
      receive(&a, 1, 0); // receive data from Proc 0
9
      printf("%d\n", a);
10
    }
11
12 }
```

Flavors Send/Receive

- Hardware+OS may support copying message into a "buffer" space to make allowing program to proceed faster
- Functions usually available to do both blocking send() and send_nonblocking() with hardware support BUT without OS/hardware support they are the same

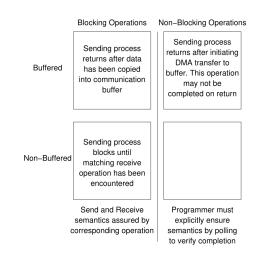


Figure 6.3 Space of possible protocols for send and receive operations.

Blocked and Unbuffered

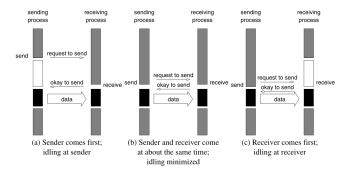


Figure 6.1 Handshake for a blocking non-buffered send/receive operation. It is easy to see that in cases where sender and receiver do not reach communication point at similar times, there can be considerable idling overheads.

Blocking/Unbuffered: no extra buffer available to hold pending sends/receives so must wait until message is sent to proceed Blocked processors are idle, do no work, which cuts into speedup

Ordering of Send Receive

```
1 P0 P1
2 3 send(&a, 1, 1); send(&a, 1, 0);
4 receive(&b, 1, 1); receive(&b, 1, 0);
```

Assuming send/receive blocked/unbuffered, what's wrong with the above code?

Blocking with Buffers

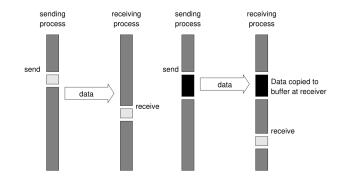


Figure 6.2 Blocking buffered transfer protocols: (a) in the presence of communication hardware with buffers at send and receive ends; and (b) in the absence of communication hardware, sender interrupts receiver and deposits data in buffer at receiver end.

Hardware buffer support, sender and receiver have a memory minion No buffer support: sender interrupts receiver

The Danger Continues

```
1 P0 P1
2 3 receive(&a, 1, 1); receive(&a, 1, 0);
4 send(&b, 1, 1); send(&b, 1, 0);
```

receive() always blocks until message is obtained

Does the above code work even in the buffered setting?

Non-blocking Communication

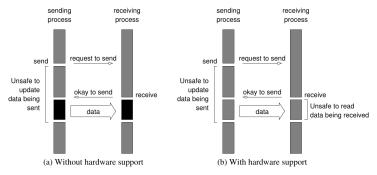
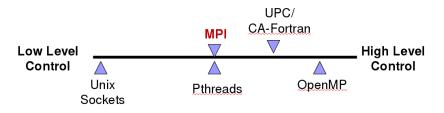


Figure 6.4 Non-blocking non-buffered send and receive operations (a) in absence of communication hardware; (b) in presence of communication hardware.

- Takes a bit more work on the programming side
- Must explicitly ensure that transaction completes with function calls
- isend(data,dest,status): send w/o waiting
- ireceive(data,dest,status): receive w/o waiting
- wait(status): wait until a message has been sent or

MPI: Message Passing Interface

- ► Standardized library of functions for C/C++/Fortran
- Communicate between processors in a distributed memory machine first appearing around 1992
- ▶ MPI Version 1.x universally deployed, Version 2.x less so
- Open source implementations: MPICH, Open MPI
- Proprietary: Intel, Platform, IBM, Platform, Cray
- Typically vendor configures MPI for particular architecture / network of a large-scale machine



MPI In a Nutshell: 6 Essential Functions

```
// Initialize and Terminate MPI
int MPI_Init(int *argc, char ***argv) ;
int MPI_Finalize() ;
```

// Get total number of processors
int MPI_Comm_size(MPI_Comm comm, int *size);

// Get logical proc number of calling process
int MPI_Comm_rank(MPI_Comm comm, int *rank);

MPI Hello World

```
1 // mpi_hello.c: C Example of hello world with MPI. Compile and run as
2 // > mpicc -o mpi_hello mpi_hello.c
3 // > mpirun ./mpi_hello  # use number of processors equal to total machine
4 // > mpirun -np 2 mpi_hello # use 2 processors
5 // > mpirun -np 8 mpi hello # use 8 processors
6
7 #include <stdio.h>
8 #include <mpi.h>
9
10 int main (int argc, char *argv[]){
  int rank;
                                  // the id of this processor
11
12 int size;
                                  // the number of processors being used
13
   MPI Init (&argc, &argv); // starts MPI
14
   MPI_Comm_rank (MPI_COMM_WORLD, &rank); // get current process id
15
    MPI_Comm_size (MPI_COMM_WORLD, &size); // get number of processes
16
17
18
    // Say hello from this proc
    printf( "Proc %d of %d says 'Hello world'\n", rank, size );
19
20
    MPI Finalize():
21
    return 0:
22
23 }
```

Compilation and Running

- Demo using openmpi implementation
- mpirun for interactive running
- mpirun -np 4 progr sets number of "processors" to 4

```
>> cd 04-mpi-code/
>> mpicc -o mpi_hello mpi_hello.c
>> ./mpi hello
. . .
Proc 0 of 1 says 'Hello world'
>> mpirun -np 2 mpi_hello
. . .
Proc 0 of 2 says 'Hello world'
Proc 1 of 2 says 'Hello world'
>> mpirun mpi_hello
. . .
Proc 2 of 4 says 'Hello world'
Proc 0 of 4 says 'Hello world'
Proc 1 of 4 says 'Hello world'
Proc 3 of 4 says 'Hello world'
```

MPI Implementations and OpenMPI Warnings

- Several Implementations of MPI:
 - ► OpenMPI and MPICH are free, open source, widely available
 - HPC Vendors like IBM and Cray provide their own tailored MPI versions
- Recent Versions of OpenMPI can complain a LOT about various items missing
- The many machines with MPI are not configured perfectly leading to additional errors
 - Example: --mca btl_base_warn_component_unused 0 to warn about missing HPC network components during mpirun
 - Example: --mca opal_warn_on_missing_libcuda 0 if not intending to use GPU libraries
- Exact nature of warnings/errors varies a lot, look at messages which often dictate how to disable them
- Provided mpiopts.sh script can be sourced to set suppress common errors

Warning Suppression in OpenMPI

>> mpicc mpi_hello_plus.c

>> mpirun -np 2 a.out

The library attempted to open the following supporting CUDA libraries, but each of them failed. CUDA-aware support is disabled. libcuda.so.1: cannot open shared object file: No such file or directory libcuda.dylib: cannot open shared object file: No such file or directory /usr/lib64/libcuda.so.1: cannot open shared object file: No such file or directory /usr/lib64/libcuda.dylib: cannot open shared object file: No such file or directory If you are not interested in CUDA-aware support, then run with --mca opal_warn_on_missing_libcuda 0 to suppress this message. If you are interested in CUDA-aware support, then try setting LD_LIBRARY_PATH to the location of libcuda.so.1 to get passed this issue. _____ P0000 [val]: Hello world from process 0 of 2 P0001 [val]: Hello world from process 1 of 2 [val:558294] 1 more process has sent help message help-mpi-common-cuda.txt / dlopen failed [val:558294] Set MCA parameter "orte base_help aggregate" to 0 to see all help / error messages >> source mpiopts.sh >> mpirun \$MPIOPTS -np 2 a.out

 P0001 [val]: Hello world from process
 1 of 2

 P0000 [val]: Hello world from process
 0 of 2

MPI Oversubscribing

Default OpenMPI config uses all processors on a single machine, fails for larger requests unless --oversubscribe

```
>> mpirun -np 2 a.out
P0001 [val]: Hello world from process 1 of 2
P0000 [val]: Hello world from process 0 of 2
>> mpirun -np 8 a.out
      _____
There are not enough slots available in the system to satisfy the 8
slots that were requested by the application:
 a.out
Either request fewer slots for your application, or make more slots
available for use.
Alternatively, you can use the --oversubscribe option to ignore the
number of available slots when deciding the number of processes to
launch.
_____
>> source mpiopts.sh
>> echo $MPTOPTS
--mca opal_warn_on_missing_libcuda 0 --oversubscribe
>> mpirun $MPIOPTS -np 16 a.out
P0009 [val]: Hello world from process 9 of 16
P0014 [val]: Hello world from process 14 of 16
P0012 [val]: Hello world from process 12 of 16
```

Hostfiles

- For simple cluster configurations, can pass a hostfile to mpirun to indicate host names of other machines in cluster
- Simplest form of hostfile is a list of symbolic or IP addresses for machines to recruit for the run
- CSE Labs CUDA cluster¹ has the following machines which can be used for experimentation

cuda01.cselabs.umn.edu cuda02.cselabs.umn.edu cuda03.cselabs.umn.edu cuda04.cselabs.umn.edu cuda05.cselabs.umn.edu

40 physical cores per machine

¹CUDA cluster is present to support this class, thus MPI is set up for it. We will also use it later for GPU programming. https://cse.umn.edu/cseit/classrooms-labs#cs

Extended Example on CUDA cluster 1/2

log in to lab machines

```
>> ssh cuda01.cselabs.umn.edu
...
csel-cuda-01>> cat hostfile-cuda-full.txt
cuda01.cselabs.umn.edu
cuda02.cselabs.umn.edu
cuda03.cselabs.umn.edu
cuda04.cselabs.umn.edu
cuda05.cselabs.umn.edu
```

compile + run mpi program

```
csel-cuda-01>> mpicc mpi_hello_plus.c
csel-cuda-01>> mpirun -hostfile hostfile-cuda-full.txt -np 32 ./a.out
No protocol specified
P0023 [csel-cuda-01]: Hello world from process 23 of 32
P0024 [csel-cuda-01]: Hello world from process 24 of 32
...
P0002 [csel-cuda-01]: Hello world from process 2 of 32
P0013 [csel-cuda-01]: Hello world from process 13 of 32
```

40 processors per machine

```
csel-cuda-01>> mpirun -hostfile hostfile-cuda-full.txt -np 64 ./a.out
...
P0011 [csel-cuda-01]: Hello world from process 11 of 64
P0020 [csel-cuda-01]: Hello world from process 20 of 64
...
P0039 [csel-cuda-01]: Hello world from process 39 of 64
P0041 [csel-cuda-02]: Hello world from process 41 of 64
P0045 [csel-cuda-02]: Hello world from process 45 of 64
P0046 [csel-cuda-02]: Hello world from process 46 of 64
```

. . .

Extended Example on CUDA cluster 2/2

```
### utilize whole cluster
csel-cuda-01>> mpirun -hostfile hostfile-cuda-full.txt -np 200 ./a.out
...
P0003 [csel-cuda-01]: Hello world from process 3 of 200
...
P0089 [csel-cuda-03]: Hello world from process 89 of 200
...
P0190 [csel-cuda-05]: Hello world from process 190 of 200
P0123 [csel-cuda-04]: Hello world from process 123 of 200
...
P0077 [csel-cuda-02]: Hello world from process 77 of 200
```

There are not enough slots available in the system to satisfy the 201 slots that were requested by the application:

./a.out

Either request fewer slots for your application, or make more slots available for use.

A "slot" is the Open MPI term for an allocatable unit where we can launch a process. The number of slots available are defined by the environment in which Open MPI processes are run:

- Hostfile, via "slots=N" clauses (N defaults to number of processor cores if not provided)
- The --host command line parameter, via a ":N" suffix on the hostname (N defaults to 1 if not provided)
- 3. Resource manager (e.g., SLURM, PBS/Torque, LSF, etc.)
- 4. If none of a hostfile, the --host command line parameter, or an RM is present, Open MPI defaults to the number of processor cores

Distributed Memory (MPI) Systems at UMN CUDA Cluster

- csel-cuda01.cselabs.umn.edu to csel-cuda05.cselabs.umn.edu
- Have OpenMPI installed, honor hostfile
- Hostfile in 04-mpi-code.zip as hostfile-cuda-full.txt
- Good for experimentation but not a true HPC system
- Requires setting up SSH keys / Known Hosts²

MSI Systems

- Will use MSI to evaluate scalability of program performance for A2
- Usually no need to use a hostfile as MPI jobs are run in batch and number of nodes is requested as part of job
- Requires use of job scheduling system SLURM, discuss later

²See Accessing Unix/Linux Programming Environments Section 3 for instructions on setting up keys for password/Duo free login to CSE Labs

MPI Send and Recieve

```
Most basic functionality is point-to-point message transfer via
MPI_Send() / MPI_Recv()
1 int count = 5;
2 int a[count]={10,20,30,40,50};
3 int b[count];
4 int partner = 1;
5 int tag = 1;
6
7 // Send contents of a to partner proc with tag=1
8 MPI_Send(a, count, MPI_INT, partner, tag, MPI_COMM_WORLD);
9
10 // Receive message into b from partner proc
11 MPI_Recv(b, count, MPI_INT, partner, tag, MPI_COMM_WORLD,
12 MPI_STATUS_IGNORE); // ignore status of receipt
```

Analyze the program send_receive_test.c

Compare with send_bugs.c which demos stall problems

Note MPI_ANY_SOURCE may be used for recv's source

Tags Make Messages Unique

```
int a[10], b[10], myrank;
MPI_Status status;
...
MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
if (myrank == 0) {
    MPI_Send(a, 10, MPI_INT, 1, 1, MPI_COMM_WORLD);
    MPI_Send(b, 10, MPI_INT, 1, 2, MPI_COMM_WORLD);
}
else if (myrank == 1) {
    MPI_Recv(b, 10, MPI_INT, 0, 2, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
    MPI_Recv(a, 10, MPI_INT, 0, 1, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
}
```

- Tags must be honored on receive
- Above code may deadlock if not buffered due to the misordering of tags
- Mostly we will use tag=1 for simplicity
- Alternatively MPI_ANY_TAG, possible to query what tag was received later on (though we won't have cause to do this)

MPI Data Types Supported

- Buffer is always untyped (void* buf)
- To strive for slightly better safety, MPI has standard datatypes

MPI_CHAR	signed char
MPI_INT	signed int
MPI_LONG	signed long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_BYTE	Last two used for sending
MPI_PACKED	structure arrays

Unsigned types also available

Exercise: Heat Transfer in MPI

- Discuss conversion of the following A1 code to an MPI version
- How is data in H divided up?
- Is communication required?
- How would one arrange MPI_Send / MPI_Recv calls?
- How much data needs to be transferred and between who?
- When the computation is finished, how can all data be displayed?

```
// Simulate the temperature changes for internal cells
for(t=0; t<max_time-1; t++){
  for(p=1; p<width-1; p++){
    double left_diff = H[t][p] - H[t][p-1];
    double right_diff = H[t][p] - H[t][p+1];
    double delta = -k*( left_diff + right_diff );
    H[t+1][p] = H[t][p] + delta;
  }
}</pre>
```

Some Patterns that occur in the problem

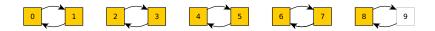
- Pair exchange of items: made easier with MPI_sendrecv
- Collecting final output for display: MPI_Gather
 - Previewed here
 - Discussed in following lectures

Exchange: Sendrecv for exchanging data between pairs

```
{
```

```
double send[10], recv[10]; int partner;
  if(procid % 2 == 1 ){ // odd procs send left, receive left
    partner = procid-1;
    MPI_Send(send, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD);
    MPI_Recv(recv, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD,
             MPI STATUS IGNORE);
  }
  else{
                        // even procs receive right, send right
    partner = procid+1;
    MPI_Recv(recv, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD,
             MPI STATUS IGNORE);
    MPI_Send(send, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD);
  }
}
{ // Sendrecv simplifies this pattern
  double send[10], recv[10]; int partner;
  partner = (procid \% 2 == 1) ? procid-1 : procid+1;
 MPI_Sendrecv(send, 10, MPI_DOUBLE, partner, 1,
               recv, 10, MPI_DOUBLE, partner, 1,
               MPI_COMM_WORLD, MPI_STATUS_IGNORE);
```

Take Care: Pair exchange can hang

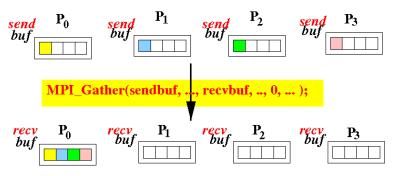


```
{
```

}

- With 9 processors, logic is broken
- Proc 8 will wait to communicate with a partner that doesn't exist
- Program never terminates

Gather Preview



Source: Shun Yan Cheung Notes on MPI

- Every processor has computed columns
- One processor (usually procid 0) needs to gather all of the data for printing / saving
- Everyone calls MPI_Gather()

MPI_Gather Sample

Use of Gather

Equivalent Non-Gather Code

```
if(rank == 0){
  for(i=0; i<100; i++){</pre>
    rbuf[i] = sendarray[i];
  }
  for(i=1; i<total_procs; i++){</pre>
    int *rloc = &rbuf[i*100];
    MPI_Recv(rloc, 100,
              MPI INT, i,
              tag, MPI_COMM_WORLD,
              MPI_STATUS_IGNORE);
 }
}
else{
  MPI_Send(sendarray, 100,
           MPI INT, 0,
           tag, MPI_COMM_WORLD);
}
```

Collective Communication Patterns Next

- gather is an example of a class of Collective Communication Patterns
- Will study more of these in subsequent lectures
- Using built-in collective comm. patterns simplifies programs
- MPI implementation typically uses the most efficient underlying communications for a particular network

Sending Structs

}

- Simple and effective if all compute nodes use the same binary layout
- MPI also provides a (complex) method for situations where struct layout differs between nodes
- Must Dictate # of struct fields, types, and ordering into a MPI_Datatype and use MPI_Type_create_struct()
- Likely hurts performance if struct layout differs so will not discuss in detail

Non-blocking Send / Receive

BLOCKING	NON-BLOCKING
MPI_Send()	<pre>MPI_Isend(, &request)</pre>
<pre>MPI_Recv()</pre>	<pre>MPI_Irecv(, &request)</pre>
SYCNRONIZE	<pre>MPI_Wait(&request,)</pre>

- Non-blocking calls trigger send/receive to be initiated but do not block process(or) to completion
- MPI_Request struct tracks whether operation has completed
- Block via MPI_wait() until send/recv completes
- Prior to blocking, unsafe to alter/use data in buffers
- Can pair MPI_ISend() / MPI_Recv() and vice versa

Faux Example of MPI_Isend() / MPI_IRecv()

```
1 int data_a[100] = {...};
 2 int data_b[100] = {...};
3 int partner = ...;
 4 int tag = ...;
5
6 ...; // compute data_a[]
7
8 MPI Request request:
9 MPI_Isend(data_a, 100, MPI_INT, partner, tag, MPI_COMM_WORLD, &request);
10
11 ...; // unsafe to alter data_a[] so compute data_b[]
12
13 MPI_Wait(&request, MPI_STATUS_IGNORE); // block until data_a[] has been sent
                                          // now safe to alter data a[]
14
15 for(int i=0; i<100; i++){
                                    // more computations on data_a[]
    data a[i]++;
16
17 }
18 ...;
19 MPI Irecv(data a, 100, MPI INT, partner, tag, MPI COMM WORLD, &request);
20
21 ...; // unsafe to do anything with data_a[], compute on data_b[]
22
23 MPI_Wait(&request, MPI_STATUS_IGNORE); // block until data_a[] has been sent
                                      // now safe to alter data_a[]
24
25 for(int i=0; i<100; i++){
                                       // more computations on data_a[]
    data a[i]++:
26
27 }
```