MPI Basics

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CS 499: Spring 2016 GMU

Logistics

Reading: Grama Ch 6 + 4

- Ch 6: MPI basics
- Ch 4: Communication patterns

Assignment 1

- Posted, due Thu 2/4
- Groups of 2 permitted
- A few tips on broadcast Problem 5
- Questions?

Today

Begin discussion of MPI programming

Generic Send and Receive

Minimum required functionality to do distributed memory parallel computing:

```
send(void *sendbuf, int nelems, int dest)
receive(void *recvbuf, int nelems, int source)
```

Sample Use

```
1  P0                      P1
2
3  a = 100;                receive(&a, 1, 0)
4  send(&a, 1, 1);         printf("%d\n", a);
5  a=0;
```

- Proc 0 sends a single integer to Proc 1
- Proc 0 then 0s that integer
- Proc 1 receives and prints the integer

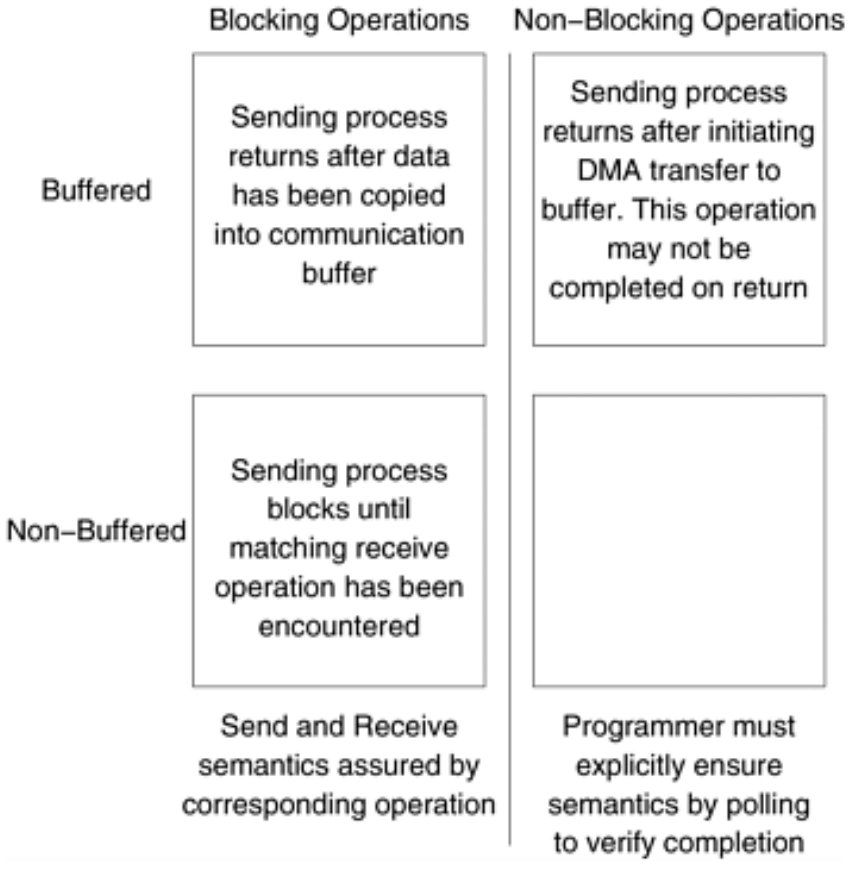
More typical appearance

Will typically write this as a single program which every processor runs.

```
void exchange(){
  int a = 100;
  int my_proc = get_processor_number();
  if(my_proc == 0){
    send(&a, 1, 1);
    a=0;
  }
  else if(my_proc == 1){
    receive(&a, 1, 0);
    printf("%d\n", a);
  }
}
```

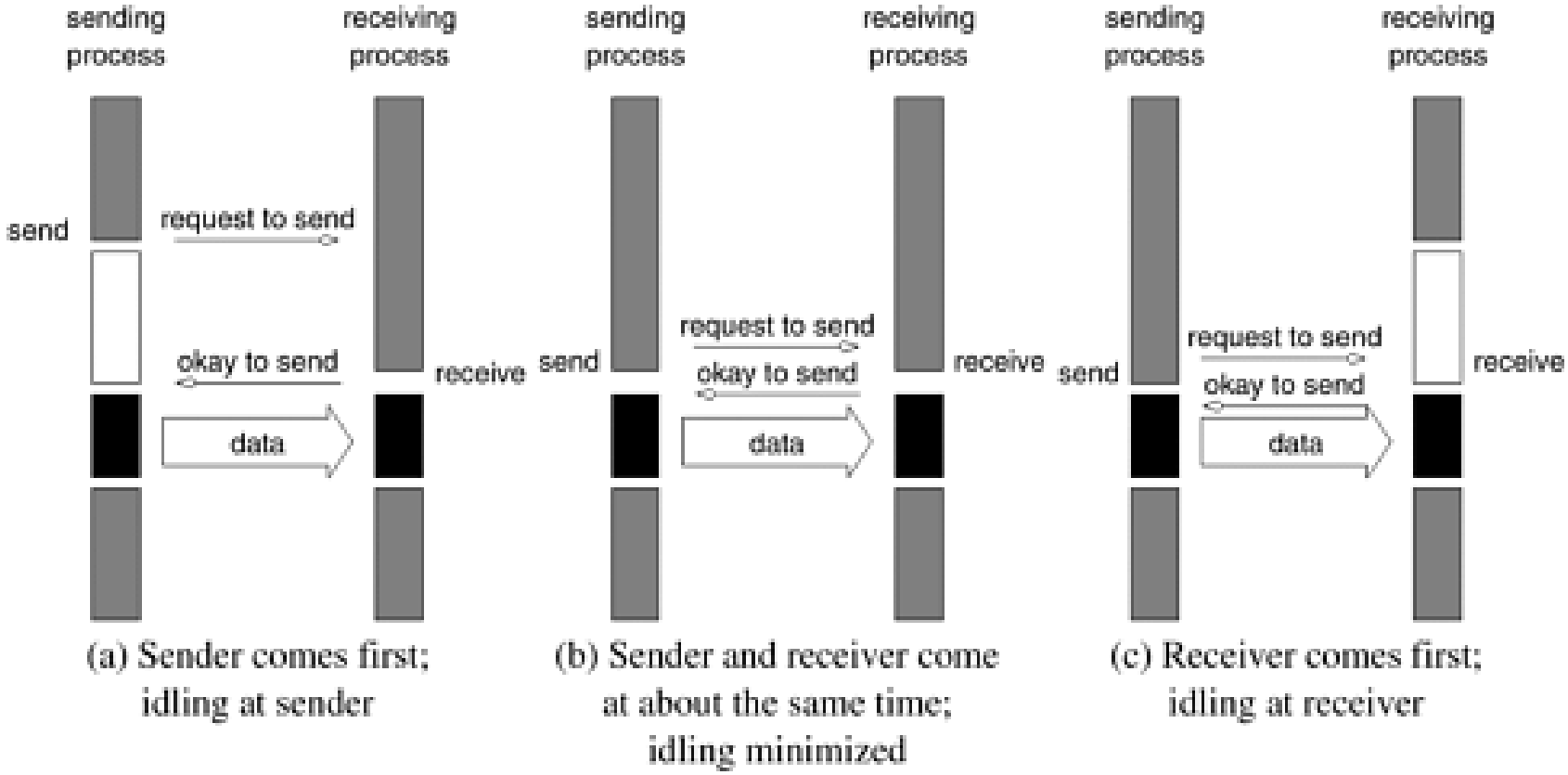
- Function to identify proc number
- Branching on proc number to take different actions

Flavors Send/Receive



- Hardware/OS support for buffered communication tends to make things run faster
- Usually have function calls available to do `send()` (blocking) and `send_nonblocking()` but must have some hardware support for it

Blocked and Unbuffered



(a) Sender comes first; idling at sender

(b) Sender and receiver come at about the same time; idling minimized

(c) Receiver comes first; idling at receiver

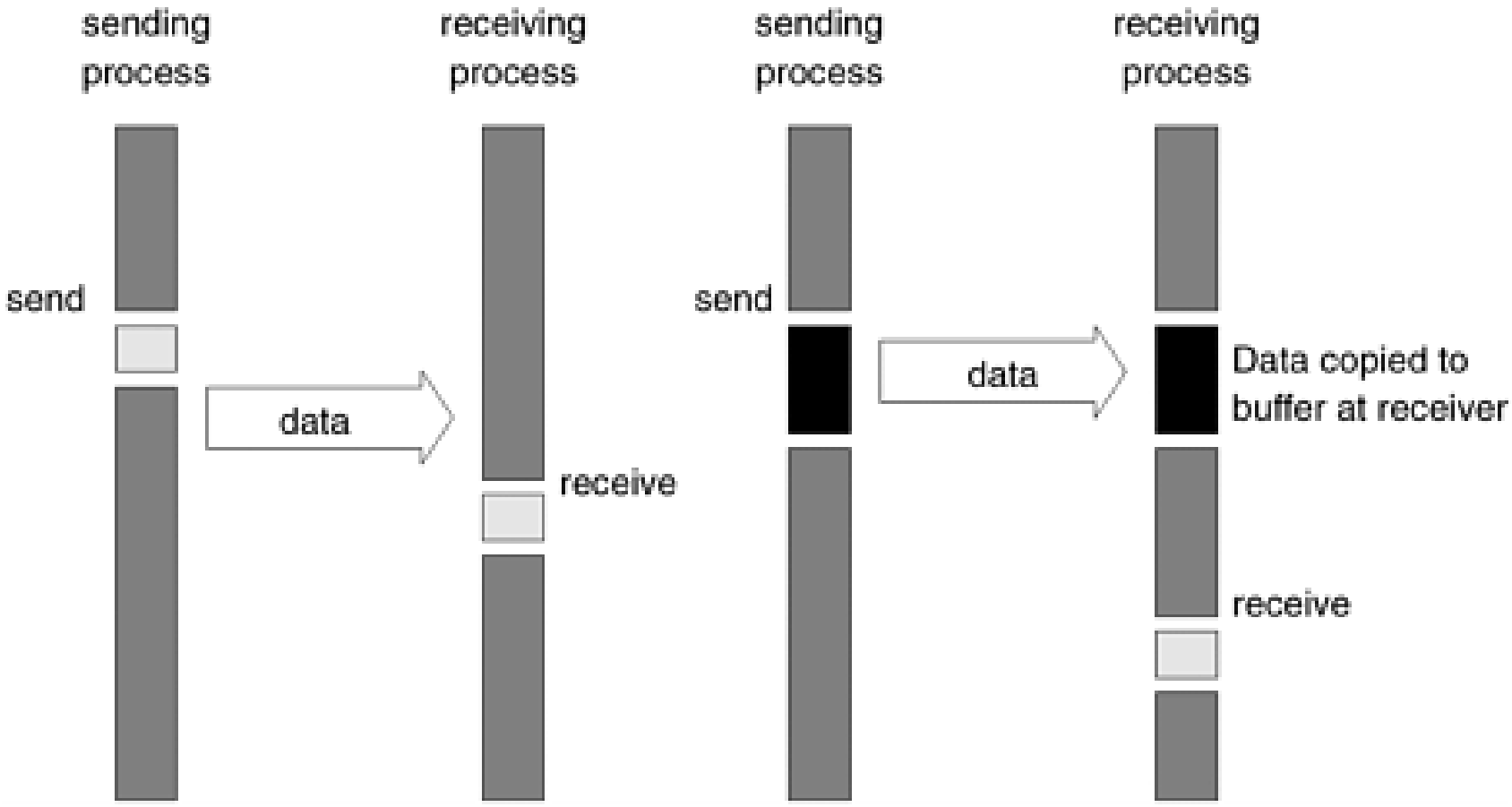
Blocking/Unbuffered: no extra buffer available to hold pending sends/receives so must wait, wait until message is sent to proceed
Blocked processors are idle, do no work, which cuts into speedup

Danger

```
1     P0                          P1
2
3     send(&a, 1, 1);             send(&a, 1, 0);
4     receive(&b, 1, 1);          receive(&b, 1, 0);
```

Assuming `send`/`receive` blocked/unbuffered, what's wrong with the above code?

Blocking with Buffers



Hardware buffer support, sender and receiver have a memory minion

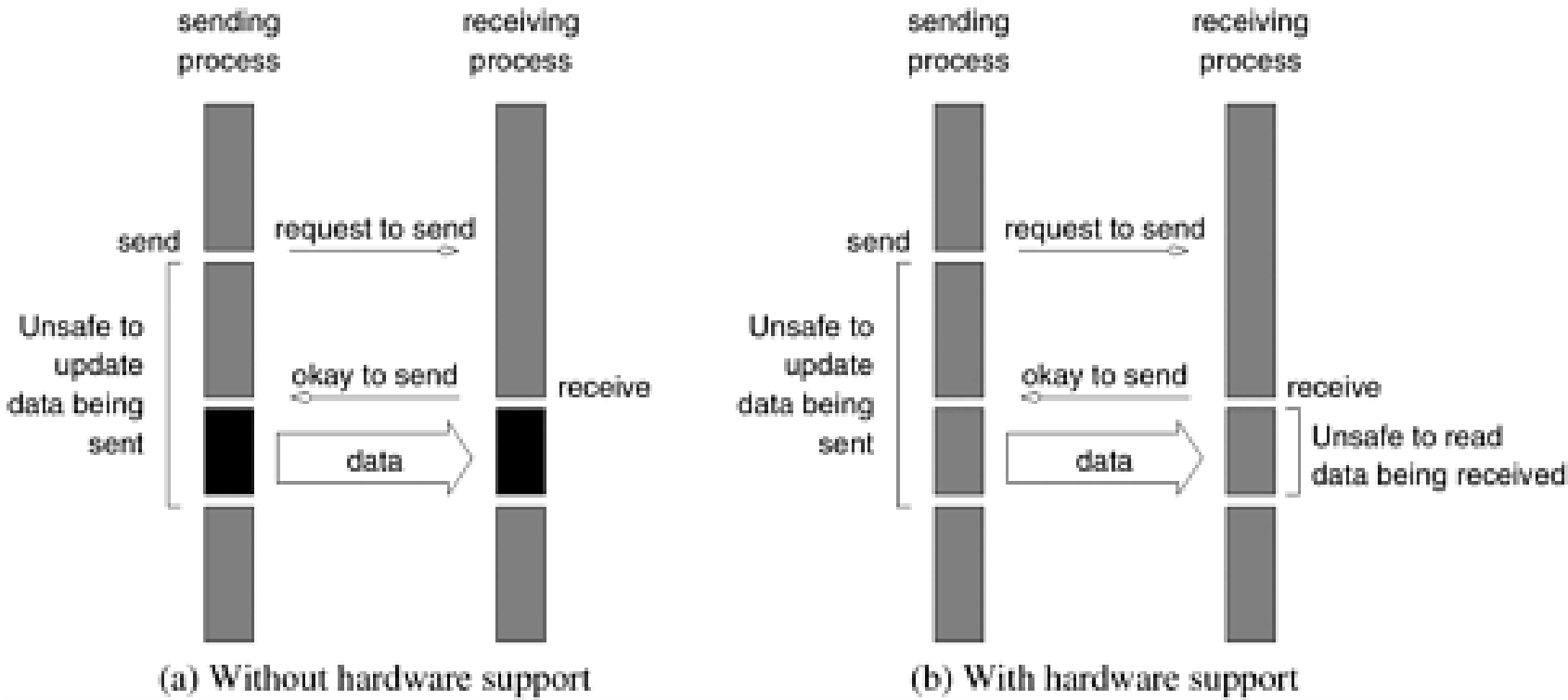
No buffer support: sender interrupts receiver

The Danger Continues

```
1  P0                          P1
2
3  receive(&a, 1, 1);          receive(&a, 1, 0);
4  send(&b, 1, 1);             send(&b, 1, 0);
```

- `receive()` always blocks until message is obtained
- Does the above code work even in the buffered setting?

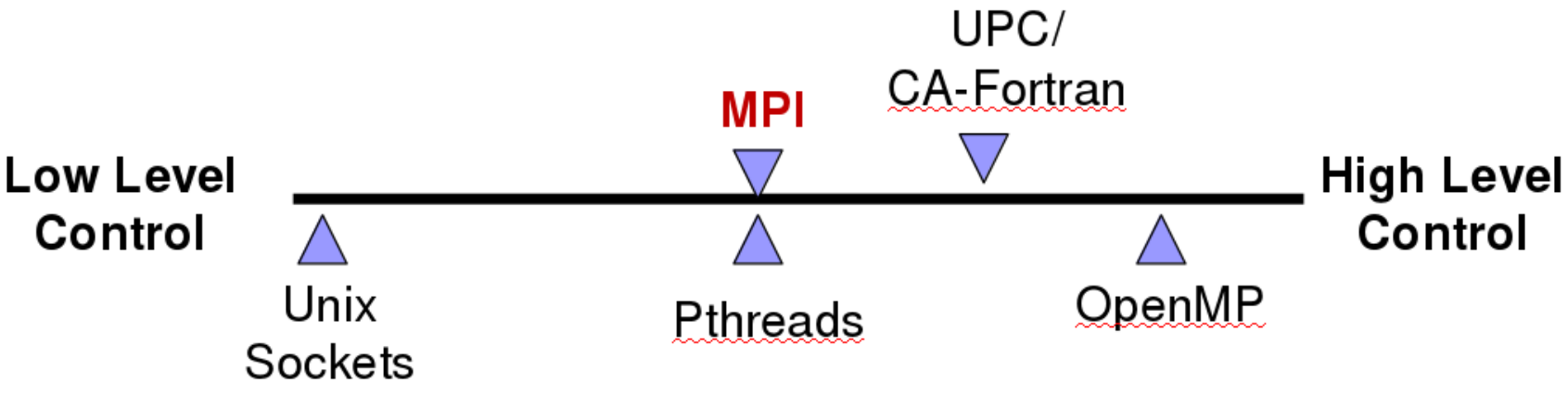
Non-blocking Communication



- Takes a bit more work on the programming side
- Must explicitly ensure that transaction completes with function calls
- `isend(data,dest,status)`: send w/o waiting
- `ireceive(data,dest,status)`: receive w/o waiting
- `wait(status)`: wait until a message has been sent or received before moving one

MPI: Message Passing Interface

- Standardized library of functions for C/C++/Fortran
- Communicate between processors in a distributed memory machine
- Open source implementations: MPICH, Open MPI
- Proprietary: Intel, Platform, IBM, Platform, Cray
- Typically geared for particular architecture
- May exploit specifics of a particular machine



MPI In a Nutshell: 6 Essential Functions

```
// Initializes MPI.
int MPI_Init(int *argc, char ***argv) ;

// Terminates MPI.
int MPI_Finalize() ;

// Determines the number of processes.
int MPI_Comm_size(MPI_Comm comm, int *size);

// Determines the label of the calling process.
int MPI_Comm_rank(MPI_Comm comm, int *rank);

// Sends a message.
int MPI_Send(void *buf, int count, MPI_Datatype datatype,
             int dest, int tag, MPI_Comm comm);

// Receives a message.
int MPI_Recv(void *buf, int count, MPI_Datatype datatype,
             int source, int tag, MPI_Comm comm,
             MPI_Status *status);
```

MPI Setup: Hello World

```
#include <stdio.h>
#include <mpi.h>

int main (int argc, char *argv[]){
  int rank, size;
  MPI_Init (&argc, &argv);                /* starts MPI */
  MPI_Comm_rank (MPI_COMM_WORLD, &rank); /* get current process id */
  MPI_Comm_size (MPI_COMM_WORLD, &size); /* get number of processes */
  int i;
  for(i=0; i<1; i++){
    printf( "Hello world from process %d of %d\n", rank, size );
  }
  MPI_Finalize();
  return 0;
}
```

- Note the use of `MPI_COMM_WORLD` which is a predefined constant corresponding to all processors.
- Can also set up other communicators that correspond to subsets of processors

Compilation and Running

- Demo using openmpi implementation
- `mpirun` for interactive running
- `mpirun -np 4 progr` sets number of "processors" to 4

```
lila [test-code]% mpicc -o hello hello.c

lila [test-code]% ./hello
Hello world from process 0 of 1

lila [test-code]% mpirun hello
Hello world from process 0 of 4
Hello world from process 1 of 4
Hello world from process 2 of 4
Hello world from process 3 of 4

lila [test-code]% mpirun -np 2 hello
Hello world from process 0 of 2
Hello world from process 1 of 2

lila [test-code]% mpirun -np 8 hello
Hello world from process 7 of 8
Hello world from process 0 of 8
Hello world from process 2 of 8
Hello world from process 3 of 8
Hello world from process 4 of 8
```

MPI Send and Recieve

```
int a[10], b[10];
int partner = 1;
...

// Send contents of a to partner proc with tag=1
MPI_Send(a, 10, MPI_INT, partner, 1, MPI_COMM_WORLD);

// Receive message into b from partner proc with tag=1,
// ignore status of receipt
MPI_Recv(b, 10, MPI_INT, partner, 1, MPI_COMM_WORLD, MPI_STATUS_IGNORE)
```

Analyze the program `send-receive-test.c`

Tag Trouble

```
int a[10], b[10], myrank;
MPI_Status status;
...
MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
if (myrank == 0) {
  MPI_Send(a, 10, MPI_INT, 1, 1, MPI_COMM_WORLD);
  MPI_Send(b, 10, MPI_INT, 1, 2, MPI_COMM_WORLD);
}
else if (myrank == 1) {
  MPI_Recv(b, 10, MPI_INT, 0, 2, MPI_COMM_WORLD);
  MPI_Recv(a, 10, MPI_INT, 0, 1, MPI_COMM_WORLD);
}
```

- Tags must be honored on receive
- Above code may deadlock if not buffered due to the misordering of tags
- Mostly we will not deal with tags (`tag=1`)

Issues with Untyped Data in MPI

```
// Sends a message.
int MPI_Send(void *buf, int count, MPI_Datatype datatype,
             int dest, int tag, MPI_Comm comm);

// Receives a message.
int MPI_Recv(void *buf, int count, MPI_Datatype datatype,
             int source, int tag, MPI_Comm comm,
             MPI_Status *status);
```

- Type of buffer is always untyped (`void* buf`)
- To try to get at slightly better safety, MPI has standard datatypes

`MPI_CHAR`	signed char
`MPI_INT`	signed int
`MPI_LONG`	signed long int
`MPI_FLOAT`	float
`MPI_DOUBLE`	double
`MPI_BYTE`	Last two used for sending
`MPI_PACKED`	structure arrays

Unsigned types also available

Exercise: Heat Transfer

- Discuss conversion of the following HW1 code to an MPI version
- How is data in `H` divided up?
- Is communication required?
- How would one arrange `MPI_Send / MPI_Recv` calls?

```
// Simulate the temperature changes for internal cells
for(t=0; t<max_time-1; t++){
  for(p=1; p<width-1; p++){
    double left_diff  = H[t][p] - H[t][p-1];
    double right_diff = H[t][p] - H[t][p+1];
    double delta = -k*( left_diff + right_diff );
    H[t+1][p] = H[t][p] + delta;
  }
}
```

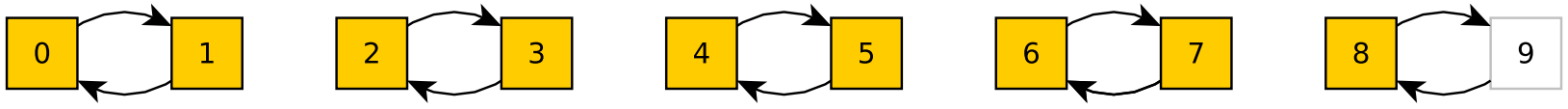
Some Patterns that occur in the problem

- Pair exchange of items: made easier with `MPI_sendrecv`
- Collecting final output for display: `MPI_Gather`

Exchange: Sendrecv for exchanging data between pairs

```
{
  double send[10], recv[10]; int partner;
  if(procid % 2 == 1 ){ // odd procs send left, receive left
    partner = procid-1;
    MPI_Send(send, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD);
    MPI_Recv(recv, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD,
             MPI_STATUS_IGNORE);
  }
  else{                    // even procs receive right, send right
    partner = procid+1;
    MPI_Recv(recv, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD,
             MPI_STATUS_IGNORE);
    MPI_Send(send, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD);
  }
}
{ // Sendrecv simplifies this pattern
  double send[10], recv[10]; int partner;
  partner = (procid % 2 == 1) ? procid-1 : procid+1;
  MPI_Sendrecv(send, 10, MPI_DOUBLE, partner, 1,
               recv, 10, MPI_DOUBLE, partner, 1,
               MPI_COMM_WORLD, MPI_STATUS_IGNORE);
}
```

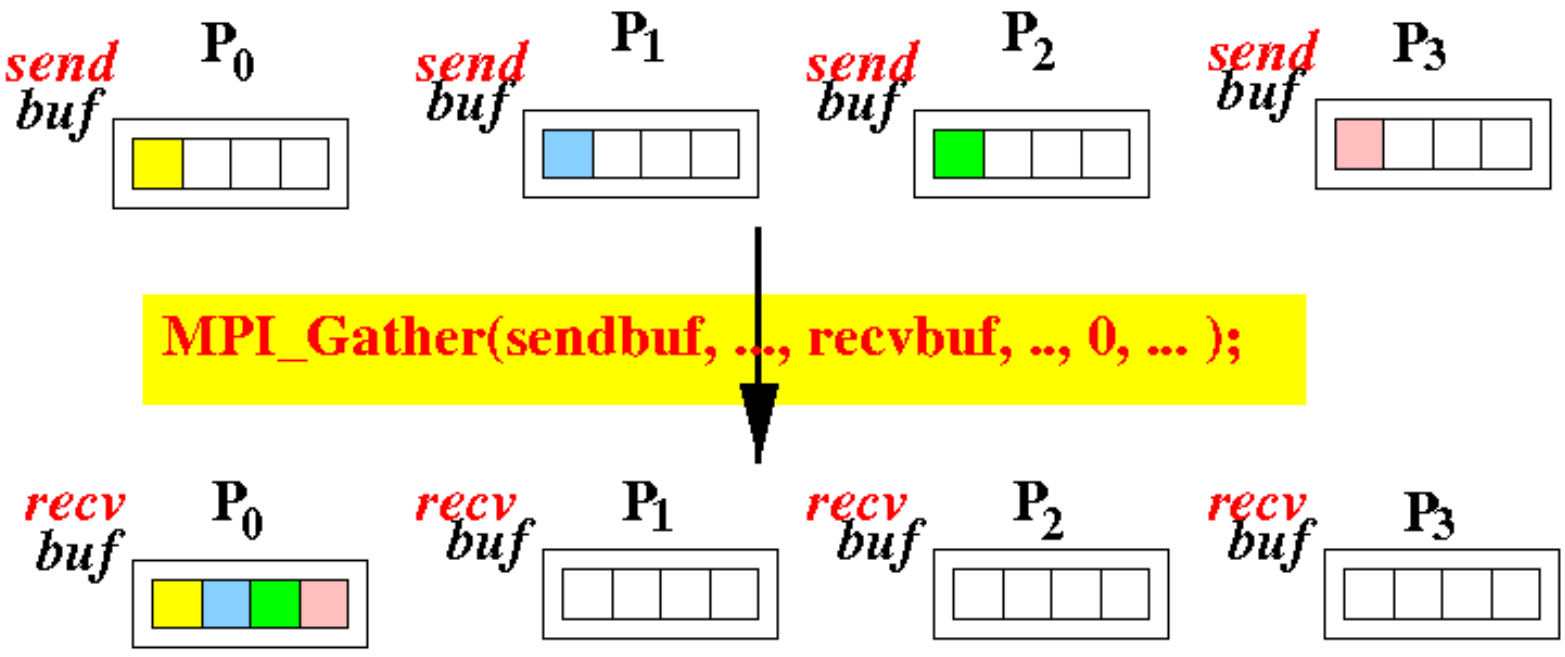
Take Care: Pair exchange can hang




```
{
  double send[10], recv[10]; int partner;
  partner = (procid % 2 == 1) ? procid-1 : procid+1;
  MPI_Sendrecv(send, 10, MPI_DOUBLE, partner, 1,
               recv, 10, MPI_DOUBLE, partner, 1,
               MPI_COMM_WORLD, MPI_STATUS_IGNORE);
}
```

- With 9 processors, logic is broken
- Proc 8 will wait to communicate with a partner that doesn't exist
- Program never terminates

Gather



Source: Shun Yan Cheung Notes on MPI

- Every processor has computed columns
- One processor (usually `procid 0`) needs to gather all of the data
- Everyone calls `MPI_Gather()`

MPI_Gather Sample

Use of Gather

```
// Preamble for any code
MPI_Comm comm = MPI_COMM_WORLD;
int sendarray[100];
int procid, total_procs, *rbuf;
...;
// Only proc 0 needs space for all
// data
if(procid == 0) {
  rbuf = malloc(total_procs*100*
                sizeof(int));
}

// Everyone calls gather
// proc 0 gets all data eventually
MPI_Gather(sendarray, 100, MPI_INT,
           rbuf, 100, MPI_INT,
           0, comm);
```

Equivalent Non-Gather Code

```
if(rank == 0){
  for(i=0; i<100; i++){
    rbuf[i] = sendarray[i];
  }
  for(i=1; i<total_procs; i++){
    int *rloc = &rbuf[i*100];
    MPI_Recv(rloc, 100,
             MPI_INT, i,
             tag, MPI_COMM_WORLD,
             MPI_STATUS_IGNORE);
  }
}
else{
  MPI_Send(sendarray, 100,
           MPI_INT, 0,
           tag, MPI_COMM_WORLD);
}
```