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$1 + g(z)$  (with  $g(z)$  given by (2.2)) produce improved values for the Golay merit factor, which measures how far  $|f(z)|$  is away from  $(n + 1)^{1/2}$  on average as  $z$  runs over  $|z| = 1$ . (We note that cyclic shifts of coefficients of  $f(z)$  do not affect the value of  $R_p(f)$ .) A nonexhaustive search of cyclic shifts of the sequences constructed in Section 2 with  $n = 138$ ,  $|S| \leq 2$ , found a sequence with  $R_a(f) = 110.2457$ , which is better than the Ruprecht et al. sequence of [19], since the length is less. Thus modifications of our construction yield good values even for  $R_a(f)$ , although there is no proof that they will work for large lengths. It is also possible to try other modifications, which might yield even better results.

#### 4. Final remarks

How large are the  $R_p(f)$  produced by the above construction for moderate lengths  $n$ ? For  $n = 82$ , the largest  $R_p(f)$  that is known is 72.02 [16]. The construction of this section produces a sequence with  $R_p(f) = 69.90$ . Surprisingly, this result is achieved with  $|S| = 1$ . As was noted at the beginning of this section, if  $|S| = 0$ , then  $R_p(f) \sim q/2$  as  $q \rightarrow \infty$  (for  $n = q-1$ ,  $q$  a prime). However, if we choose  $|S| = 1$ ,  $a = -b = 1$ , then we obtain  $R_p(f) \sim 9q/10$  as  $q \rightarrow \infty$ . Choices of  $S$  with  $|S| \geq 2$  give better  $R_p(f)$  only for  $q \gtrsim 130$ , and the improvement is slight initially. (We note also that while  $R_p(f)$  is the same for all choices of  $S$  with  $|S| = 1$ ,  $a = -b$ , the precise selection of  $S$  does make a slight difference for  $|S| \geq 2$ .) The resulting sequences for  $p < 180$  are not as good (say, when judged by the value of the ratio  $R_p(f)/(n+1)$ ) as the sequence obtained from the 13-term Barker sequence (see the discussion preceding Eq. (1.8)), but they are better than some other sequences that have been proposed. For example, Ruprecht, Neeser, and Hufschmid [19] list a sequence with  $n = 143$  and  $R_p(f) = 120.69862$ . Our construction with  $n = q-1 = 138$  and  $|S| = 2$  yields a value of  $R_p(f) = 121.32578$ , so that  $R_p(f)$  is higher even though  $n$  is lower. (It should be mentioned, though, that the Ruprecht et al. sequence was chosen to have a high  $R_a(f)$ , not a high  $R_p(f)$ .)

The construction of Theorem 1 produces a sequence with high  $R_p(f)$  because the polynomials associated to the Legendre sequences already have the desired behavior at the points  $z = \exp(2\pi ik/q)$  for  $1 \leq k \leq q-1$ , and it is only at  $z = 1$  that they need to be modified. Unfortunately the behavior of these polynomials at other points on the unit circle is not as well controlled. Montgomery [14] showed that

$$\max_{|z|=1} |f(z)| > 2\pi^{-1} q^{1/2} \log \log q \quad (4.1)$$

for all sufficiently large  $q$ , and he conjectured that this bound is of the right order of magnitude. If Montgomery's conjecture is right, these polynomials will be smaller than random ones, which reach  $q^{1/2}(\log q)^{1/2}$  in size (cf. [16]). However, these polynomials do have small minimal absolute values. B. Conrey and A. Granville have observed (unpublished) that the polynomial  $g(z)$  of Eq. (2.2) has  $> p/2$  zeros with  $|z| = 1$ . Therefore it is not straightforward to modify those polynomials to obtain large  $R_a(f)$ . The highest value of  $R_a(f)$  that our construction obtains for  $n = 138$ ,  $|S| \leq 2$  is 28.764, while the Ruprecht et al. sequence of [19] has  $R_a(f) = 110.57658$ . However, there are ways of modifying our construction to obtain higher values of  $R_a(f)$ . For example, it is known (see [16] for references) that cyclic shifts of the coefficients of

where the  $a_k$  are real constants,  $|a_k| \leq 1$  for all  $k$ , and the  $\tau_k$  are independent random variables with

$$Pr(\tau_k = -\gamma_k) = 1 - \gamma_k, \quad Pr(\tau_k = 1 - \gamma_k) = \gamma_k, \quad (3.8)$$

for some constants  $\gamma_k$ ,  $0 \leq \gamma_k \leq 1$ , then

$$Pr\left(|W| > C \left(\sum_{k=1}^n \gamma_k\right)^{1/2} (\log n)^{1/2}\right) < n^{-10}. \quad (3.9)$$

**Proof.** We have, for any  $\lambda > 0$ ,

$$Pr(W > x)e^{\lambda x} \leq \mathcal{E}(e^{\lambda W}). \quad (3.10)$$

Now the  $\tau_k$  are independent, so

$$\mathcal{E}(e^{\lambda W}) = \prod_{k=1}^n \mathcal{E}(e^{\lambda \tau_k a_k}). \quad (3.11)$$

We next note that

$$\mathcal{E}(e^{\lambda \tau_k a_k}) = e^{-\lambda \gamma_k a_k} (1 - \gamma_k) + e^{\lambda (1 - \gamma_k) a_k} \gamma_k \leq e^{C' \lambda^2 \gamma_k} \quad (3.12)$$

if  $C'$  is sufficiently large. Therefore

$$Pr(W > x) \leq \exp\left(C' \lambda^2 \sum_{k=1}^n \gamma_k - \lambda x\right). \quad (3.13)$$

This bound holds for all  $\lambda > 0$ , so for  $x > 0$  we select  $\lambda = x(2C' \sum \gamma_k)^{-1}$  and obtain

$$Pr(W > x) \leq \exp\left(-x^2 \left(4C' \sum_{k=1}^n \gamma_k\right)^{-1}\right). \quad (3.14)$$

Since the same bound for  $Pr(W < -x)$  follows by applying (3.14) to the problem with  $a_k$  replaced by  $-a_k$ , we easily obtain the claim of the lemma.  $\blacksquare$

To conclude the proof of Theorem 2, we apply Lemma 1 to the real and imaginary parts of

$$h(\zeta^j) - \mathcal{E}(h(\zeta^j)), \quad 0 \leq j \leq q - 1.$$

We find that with probability  $\geq 1 - n^{-8}$ ,

$$\left|h(\zeta^j) - \mathcal{E}(h(\zeta^j))\right| < 10Cq^{1/4}(\log q)^{1/2} \quad (3.15)$$

holds for all  $j$ ,  $0 \leq j \leq q - 1$ . Therefore

$$|f(\zeta^j)| = q^{1/2} + O(q^{1/4}(\log q)^{1/2}) \quad (3.16)$$

for all  $j$ , which yields Theorem 2.

There is a method of Kolountzakis [8] that often manages to remove factors such as the  $(\log q)^{1/2}$  in the estimate (3.16). However, it does not seem to apply in this case.

### 3. Proof of Theorem 2

Theorem 2 follows from a modification of the proof of Theorem 1, using methods similar to those of [15]. As in the preceding section, we define  $f(z)$  by (2.6) and (2.7) with  $a = 1$ . However, this time we will take  $S$  to be of size about  $q^{1/2}$ , and it will contain only nonresidues. The set  $S$  will be chosen at random, with each  $k$ ,  $1 \leq k \leq q - 1$ ,  $\left(\frac{k}{q}\right) = -1$ , selected independently to be in  $S$  with probability

$$Pr(k \in S) = q^{-1/2}/2 . \quad (3.1)$$

Thus we have

$$h(z) = 1 - 2 \sum_{k=1}^{q-1} \eta_k \left(\frac{k}{q}\right) z^k , \quad (3.2)$$

where  $\eta_k = 0$  or  $1$  is a random variable with  $\eta_k$  identically  $0$  if  $\left(\frac{k}{q}\right) = 1$ , and  $\mathcal{E}(\eta_k) = q^{-1/2}/2$  if  $\left(\frac{k}{q}\right) = -1$ .

We need to determine the behavior of  $h(\zeta^j)$  for  $0 \leq j \leq q - 1$ , where  $\zeta$  is defined by (2.1). We first consider the expected value  $\mathcal{E}(h(\zeta^j))$  for a fixed  $j$ . For  $j = 0$  we have

$$\mathcal{E}(h(1)) = 1 + 2 \sum_{\substack{k=1 \\ \left(\frac{k}{q}\right)=-1}}^{q-1} \mathcal{E}(\eta_k) = 1 + (q-1)q^{-1/2} = q^{1/2} + 1 - q^{-1/2} . \quad (3.3)$$

For  $1 \leq j \leq q - 1$ , on the other hand,

$$\mathcal{E}(h(\zeta^j)) = 1 + q^{-1/2} \sum_{\substack{k=1 \\ \left(\frac{k}{q}\right)=-1}}^{q-1} \zeta^{kj} . \quad (3.4)$$

Since  $\sum_{k=0}^{q-1} \zeta^{kj} = 0$  for  $1 \leq j \leq q - 1$ , the sum in (3.4) is  $-\left(\left(\frac{j}{q}\right)q(\zeta) + 1\right)/2$ . Hence

$$\mathcal{E}(h(\zeta^j)) = 1 - q^{-1/2}/2 - \left(\frac{j}{q}\right)q^{-1/2}g(\zeta)/2 , \quad (3.5)$$

and so

$$\mathcal{E}(h(\zeta^j)) = O(1) . \quad (3.6)$$

We conclude that  $\mathcal{E}(h(\zeta^j))$  has the desired behavior uniformly for all  $j$ ,  $0 \leq j \leq q - 1$ .

It remains to prove that for some choice of coefficients,  $h(\zeta^j)$  will be close to  $\mathcal{E}(h(\zeta^j))$  for all  $j$ . This will follow from the following result, which is similar to those in [6, 15].

**Lemma 1.** *There exists a constant  $C > 0$  such that if*

$$W = \sum_{k=1}^n \tau_k a_k , \quad (3.7)$$

where  $b = \pm 1$ . The precise selection of  $a$  and  $S$  will be discussed later. We now observe that all coefficients of  $f(z)$  are  $\pm 1$ . Further, we have

$$f(1) = g(1) = a - 2b|S|. \quad (2.9)$$

For  $1 \leq j \leq q-1$ ,

$$|f(\zeta^j)|^2 = |g(\zeta^j) + h(\zeta^j)|^2 = q + |h(\zeta^j)|^2 + 2 \operatorname{Re} \left( \overline{g(\zeta^j)} h(\zeta^j) \right). \quad (2.10)$$

Since  $|S| < q^{1/2}/100$ , we find that for large  $q$ ,

$$|h(\zeta^j)| < q^{1/2}/10 = |g(\zeta^j)|/10. \quad (2.11)$$

Therefore we can write, for  $1 \leq j \leq q-1$ ,

$$|f(\zeta^j)|^{-2} = q^{-1} \left( 1 - 2q^{-1} \operatorname{Re} \left( \overline{g(\zeta^j)} h(\zeta^j) \right) + O(q^{-1}|h(\zeta^j)|^2) \right). \quad (2.12)$$

This implies, by (2.3), that

$$|f(\zeta^j)|^{-2} = q^{-1} \left( 1 - 2q^{-1} \binom{j}{q} \operatorname{Re} \left( \overline{g(\zeta)} h(\zeta^j) \right) + O(q^{-1}|h(\zeta^j)|^2) \right), \quad (2.13)$$

and therefore

$$\sum_{j=1}^{q-1} |f(\zeta^j)|^{-2} = \frac{q-1}{q} - 2q^{-2} \operatorname{Re} \overline{g(\zeta)} \sum_{j=1}^{q-1} \binom{j}{q} h(\zeta^j) + O \left( q^{-2} \sum_{j=1}^{q-1} |h(\zeta^j)|^2 \right). \quad (2.14)$$

Now

$$\sum_{j=1}^{q-1} |h(\zeta^j)|^2 \leq \sum_{j=0}^{q-1} |h(\zeta^j)|^2 = 4q|S|. \quad (2.15)$$

On the other hand, by (2.8),

$$\begin{aligned} \sum_{j=1}^{q-1} \binom{j}{q} h(\zeta^j) &= a \sum_{j=1}^{q-1} \binom{j}{q} - 2b \sum_{k \in S} \sum_{j=1}^{q-1} \binom{j}{q} \zeta^{kj} \\ &= -2b \sum_{k \in S} \left( \frac{k}{q} \right) g(\zeta) = -2g(\zeta)|S|. \end{aligned} \quad (2.16)$$

If we now combine (2.9), (2.14), (2.15), and (2.16), we find that

$$\sum_{j=0}^{q-1} |f(\zeta^j)|^{-2} = (2b|S| - a)^{-2} + \frac{q-1}{q} + O(q^{-1}|S|). \quad (2.17)$$

If we select  $|S| \sim q^{1/3}$  as  $q \rightarrow \infty$ , say, then we obtain

$$\sum_{j=0}^{q-1} |f(\zeta^j)|^{-2} = 1 + O(q^{-2/3}), \quad (2.18)$$

which yields the claim of Theorem 1.

$$g(z) = \sum_{k=1}^{q-1} \left(\frac{k}{q}\right) z^k, \quad (2.2)$$

where  $\left(\frac{k}{q}\right)$  is the Legendre symbol. (Thus  $\left(\frac{k}{q}\right)$  is 0 for  $k = 0$ , 1 if  $k$  is a nonzero quadratic residue modulo  $q$ , and  $-1$  if  $k$  is a nonresidue modulo  $q$ .) The  $g(\zeta^j)$  are Gauss sums, and have an extensive literature. It is known (and easy to derive [1]) that

$$g(1) = 0, \quad g(\zeta^j) = \left(\frac{j}{q}\right) g(\zeta) \quad \text{for } 1 \leq j \leq q-1. \quad (2.3)$$

It is also easy to see (cf. [1]) that

$$g(\zeta)^2 = (-1)^{(q-1)/2} q. \quad (2.4)$$

It is further known that

$$g(\zeta) = \begin{cases} q^{1/2}, & q \equiv 1 \pmod{4}, \\ iq^{1/2}, & q \equiv 3 \pmod{4}, \end{cases} \quad (2.5)$$

but this is much harder to prove, and we will not use it. It is also known that  $g(z)$  is large for some  $z$  with  $|z| = 1$  [14].

We cannot use the sequence of coefficients of  $g(z)$ , because (i)  $a_0 = 0$  and (ii)  $g(1) = 0$ . The main idea behind the construction below is to modify  $g(z)$  slightly. We note that if we take  $f(z) = 1 + g(z)$ , then the coefficient sequence does consist of  $\pm 1$ 's, and  $f(1) = 1$ ,  $|f(\zeta^k)| \geq q^{1/2} - 1$  for  $1 \leq k \leq q-1$ . Therefore  $R(f) \sim q/2$  as  $q \rightarrow \infty$ , and this already gives a merit factor far superior to that of almost all  $\pm 1$  sequences.

We set

$$f(z) = g(z) + h(z), \quad (2.6)$$

where

$$h(z) = a - 2 \sum_{k \in S} \left(\frac{k}{q}\right) z^k, \quad (2.7)$$

$a = \pm 1$ , and  $S \subseteq \{1, \dots, q-1\}$ ,  $|S| < q^{1/2}/100$ . It is easy to see, using the results on maximal values of random trigonometric polynomials, that random choices of  $S$  give  $R(f) \sim n$  as  $n \rightarrow \infty$ . What we show, however, is that a nonrandom choice produces much better answers due to the special number theoretic properties of the Legendre sequence. We will select  $S$  to consist entirely of residues or else entirely of nonresidues, so that

$$\left(\frac{k}{q}\right) = b \quad \text{for all } k \in S, \quad (2.8)$$



The construction of Theorem 1 produces sequences for which  $n^{-1/2}|f(\exp(2\pi ik/(n+1)))| = 1 + o(1)$  as  $n \rightarrow \infty$  uniformly in  $k$  satisfying  $1 \leq k \leq n$ . For  $k = 0$ , though,  $|f(1)|$  is of order  $n^{1/3}$ . However, we prove the following result.

**Theorem 2.** *If  $n = q - 1$  for  $q$  a prime, then there exists a sequence  $a_0, \dots, a_n$  with  $a_j = \pm 1$  for all  $j$  such that*

$$n^{-1/2}|f(\exp(2\pi ik/(n+1)))| = 1 + O(n^{-1/4}(\log n)^{1/2}) \quad \text{as } n \rightarrow \infty \quad (1.11)$$

uniformly in  $k$ ,  $0 \leq k \leq n$ .

If we use only the bound (1.11) for the sequences of Theorem 2, we find that these sequences have  $R_p(f) \geq n - c'n^{3/4}(\log n)^{1/2}$  for some constant  $c' > 0$ . With more care, one can show that these sequences have larger  $R_p(f)$ , but the bound for  $n - R_p(f)$  that one can prove for these sequences appears to be considerably weaker than that given by Theorem 1 for its sequences.

We note that if

$$n^{-1/2}|f(e^{2\pi ik/(n+1)})| = 1, \quad 0 \leq k \leq n, \quad (1.12)$$

which is equivalent to  $R_p(f) = n + 1$ , then  $a_0, \dots, a_n$  is a Barker sequence and also the first row of a circulant Hadamard matrix, and so is thought not to exist for  $n > 3$  [3, 21]. However, there is still no proof of this conjecture.

We leave several problems open. For example, can Theorems 1 and 2 be generalized so that  $n$  does not have to be of the form  $n = q - 1$  for  $q$  a prime? Also, can one prove analogs of Theorems 1 and 2 for the aperiodic merit factor  $R_a(f)$ ? Numerical evidence (cf. [16]) suggests that there do exist  $\pm 1$  sequences  $a_0, \dots, a_n$  for  $n \geq 10$  such that the associated polynomials  $f(z)$  have

$$\min_{|z|=1} n^{-1/2}|f(z)| \geq 1/2. \quad (1.13)$$

A sequence satisfying (1.13) is guaranteed to have  $R_a(f) \geq n/4$ . However, since  $R_a(f)$  is an average result, we might expect that some of these sequences might have  $R_a(f) \sim n$  as  $n \rightarrow \infty$ . That is what seems to happen for the sequences listed in [16].

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## 2. Proof of Theorem 1

Let  $q$  be an odd prime, and define

$$\zeta = \exp(2\pi i/q), \quad (2.1)$$

associated to a  $\pm 1$  sequence of 169 terms, and

$$R_a(f(z^{13})f(z)) = 153.1014\dots, \quad R_p(f(z^{13})f(z)) = 154.6331\dots \quad (1.8)$$

However, even this construction does not produce good asymptotic results.

The main result of this note is to show that high periodic Ruprecht merit factors can be achieved for a dense sequence of values of  $n$ .

**Theorem 1.** *There is a constant  $c > 0$  such that if  $n = q - 1$  for  $q$  a prime, then there exists a sequence  $a_0, \dots, a_n$  with  $a_j = \pm 1$  for all  $j$  such that*

$$n - cn^{1/3} \leq R_p(f) \leq n + 1. \quad (1.9)$$

The proof of Theorem 1, given in Section 2, shows how to construct these sequences. The sequences of Theorem 1 do have higher  $R_a(f)$  than random sequences, but not very high ones. There is a discussion of this disappointing behavior in Section 4.

The search for  $\pm 1$  sequences with large Ruprecht merit factors is just one part of the huge subject of extremal and statistical properties of  $\pm 1$  sequences. For results, surveys, and applications, see [11, 16, 20]. In particular, there are connections to the search for sequences with large Golay merit factor [5, 12, 16].

For  $R_p(f)$  to be large,  $|f(\exp(2\pi ik/(n+1)))|$  has to be large for most  $k$ . Erdős [4] and Littlewood [9, 10] have raised the question of whether there exist  $\pm 1$  sequences  $a_0, \dots, a_n$  such that the associated polynomials  $f(z)$  satisfy

$$\min_{|z|=1} n^{-1/2}|f(z)| = 1 + o(1) \quad \text{as } n \rightarrow \infty. \quad (1.10)$$

If such sequences existed, then we would have  $R_a(f) \sim n$  and  $R_p(f) \sim n$  as  $n \rightarrow \infty$  for their polynomials. The current evidence is that such sequences don't exist (cf. [16]). However, to obtain large  $R_p(f)$  we do not require (1.10) to hold. We even do not require  $n^{-1/2}|f(\exp(2\pi ik/(n+1)))| = 1 + o(1)$  as  $n \rightarrow \infty$  to hold uniformly for all  $k$ ,  $0 \leq k \leq n$ . Instead, we prove Theorem 1 by modifying the Legendre sequence  $a_j = \left(\frac{j}{q}\right)$ . It is easy to see that modifications of that sequence achieve  $R_p(f) \sim n$  as  $n = q - 1 \rightarrow \infty$ , but the difference  $R_p(f) - n$  usually turns out to be much larger than  $n^{1/3}$  when one uses some of the obvious methods. By a careful analysis of what happens to  $R_p(f)$  as the Legendre sequence is changed, we can obtain the bound of Theorem 1.

defined as

$$R_a(f) = \left( \int_0^1 |f(e^{2\pi it})|^{-2} dt \right)^{-1}. \quad (1.5)$$

(For  $R_a(f)$  to exist, we require that  $f(z)$  be an invertible sequence.) Sequences with large  $R_a(f)$  are more desirable than those with large  $R_p(f)$ , since they can be used for transmission [19], not just for multipath estimation. Unfortunately while we will provide constructions of sequences with large  $R_p(f)$ , the problem of obtaining large  $R_a(f)$  remains open.

Since

$$\sum_{k=0}^n \left| f(e^{2\pi ik/(n+1)}) \right|^2 = (n+1)^2 \quad (1.6)$$

and

$$\int_0^1 |f(e^{2\pi it})|^2 dt = n+1 \quad (1.7)$$

by a familiar calculation, the Cauchy-Schwarz inequality shows that  $R_p(f) \leq n+1$ ,  $R_a(f) \leq n+1$  for any sequence  $a_0, \dots, a_n$ . How close can  $R_a(f)$  and  $R_p(f)$  come to  $n+1$ ? Ruprecht [18] lists in Table B.6 the sequences  $a_0, \dots, a_n$  with the highest values of  $R_p(f)$  for  $n \leq 29$ , as well as some sequences with high values of  $R_p(f)$  for  $30 \leq n \leq 32$ . The maximal value of  $R_p(f)$  for  $n = 29$  is 26.6583, for example. Ruprecht also gives, in Table B.8, the best sequences drawn from a restricted class, that of the *skew-symmetric*  $a_j$  (i.e., those with even  $n$  and  $a_{n/2-r} = (-1)^r a_{n/2+r}$ ) for  $n \leq 44$ . (The value for  $n = 44$  is incorrect, though. See [16].) The maximal value of  $R_p(f)$  for  $n = 42$  is 37.4244. In Tables B.9 and B.10 of [18] Ruprecht lists sequences with large  $R_a(f)$ , for  $n \leq 23$  in the general case and  $n \leq 44$  for the skew-symmetric case. For example, for  $n = 44$  he gives a skew-symmetric sequence with  $R_a(f) = 39.7753$ . Most of the values, especially for large  $n$ , are not known to be maximal. Skew-symmetric sequences with large  $R_a(f)$  and  $R_p(f)$  for  $n \leq 90$  (obtained from a search for other types of extremal  $\pm 1$  sequences) are given in [16]. The nonexhaustive search for high  $R_a(f)$  and  $R_p(f)$  that is documented in that paper has produced a value of  $R_p(f) = 77.5820$  for  $n = 90$ , for example.

What can one do for larger lengths  $n$ ? Random choices of the  $a_j$  almost always give small values of  $R(f)$  (cf. [16]). This is because random trigonometric polynomials have small minimal absolute values [7, 17], as was conjectured by Littlewood [9, 10]. Thus the situation is completely different than it is in coding theory, where random codes are good.

Sometimes one can construct a sequence with a large Ruprecht merit factor from shorter sequences. For example, if  $n = 12$  and  $(a_0, \dots, a_n) = (1, 1, 1, 1, 1, -1, -1, 1, 1, -1, 1, -1, 1)$  is the 13-term Barker sequence, with associated polynomial  $f(z)$ , then  $f(z^{13})f(z)$  is a polynomial

# Construction of invertible sequences for multipath estimation

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## 1. Introduction

In the Ph.D. thesis [18], written under the supervision of Jim Massey, Jürg Ruprecht has proposed coding schemes designed for effective multipath estimation. Such schemes might be useful in indoor wireless systems [19, 21] or other communication settings. These schemes use *invertible sequences*, which are sequences  $a_0, \dots, a_n$ , with  $a_j = \pm 1$  for each  $j$ , such that the associated polynomial

$$f(z) = \sum_{j=0}^n a_j z^j \quad (1.1)$$

satisfies

$$f(e^{2\pi i t}) \neq 0 \quad \text{for all real } t. \quad (1.2)$$

In some situations these schemes use *invertible periodic sequences*, for which the polynomial  $f(z)$  only has to satisfy

$$f(e^{2\pi i k/(n+1)}) \neq 0, \quad 0 \leq k \leq n. \quad (1.3)$$

(These invertible periodic sequences possess inverses under periodic convolution, which is required for Ruprecht's maximum likelihood estimation methods [18].) For best performance in estimating multipath interference, it is desirable to find invertible periodic sequences that maximize

$$R_p(f) = \frac{n+1}{\sum_{k=0}^n \left| f(e^{2\pi i k/(n+1)}) \right|^{-2}}. \quad (1.4)$$

In [18], this figure of merit is referred to as even processing gain  $G_e^{(vs)}$  of a sequence  $s$  and its periodic inverse  $v$ , and is defined in a much more complicated form. However, a short calculation based on the formulas on p. 27 and in Appendix A of [18] shows that it equals our  $R_p(f)$ . We will call  $R_p(f)$  the periodic Ruprecht merit factor, to distinguish it from other merit factors, such as that of Golay [5, 12, 16], as well as the aperiodic Ruprecht merit factor,

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*Dedicated to Jim Massey on the occasion of his 60th birthday*

## ABSTRACT

J. Ruprecht has proposed coding schemes that allow for multipath estimation. They use sequences  $a_0, \dots, a_n$  with  $a_j = \pm 1$  for each  $j$  such that the associated polynomial  $f(z) = \sum a_j z^j$  has a large

$$R_p(f) = \frac{n+1}{\sum_{k=0}^n |f(e^{2\pi i k/(n+1)})|^{-2}}.$$

Most sequences have a small  $R_p(f)$ , and those with maximal  $R_p(f)$  are hard to find. This note shows for  $n$  of the form  $n = q - 1$ ,  $q$  a prime, one can construct sequences with  $R_p(f) \geq n - O(n^{1/3})$ . Since  $R_p(f) \leq n + 1$  for any sequence, this construction is asymptotically close to optimal. It also produces large values of  $R_p(f)$  for small  $n$ .

It is also shown that for  $n = q - 1$ ,  $q$  a prime, there exist sequences  $a_0, \dots, a_n$  such that the associated polynomial  $f(z)$  satisfies

$$|f(e^{2\pi i k/(n+1)})| = (1 + o(1))n^{1/2} \quad \text{as } n \rightarrow \infty$$

uniformly for  $0 \leq k \leq n$ .