## Math 5251 Error-correcting codes and finite fields Spring 2022, Vic Reiner Midterm exam 2

## Due Wednesday Apr. 6 by 11:59pm, on Canvas

**Instructions:** This is an open book, open library, open notes, open web, take-home exam, but you are *not* allowed to collaborate. The instructor is the only human source you are allowed to consult.

1. (30 points total; 5 points each) True or False. Your answers must be justified either by counterexamples or proofs to receive full credit.

(c) There exists an integer  $m \ge 1$  for which  $\mathbb{Z}/(5^m - 1)$  is a field.

(d) When *n* is odd, an  $\mathbb{F}_2$ -linear code  $\mathcal{C}$  and its dual code  $\mathcal{C}^{\perp}$  inside  $(\mathbb{F}_2)^n$  will always intersect only in the zero vector  $\underline{0}$ , that is,  $\mathcal{C} \cap \mathcal{C}^{\perp} = \{\underline{0}\}$ .

(e) Let  $\mathcal{C}$  be the  $\mathbb{F}_{11}$ -linear code in  $(\mathbb{F}_{11})^6$  whose dual code  $\mathcal{C}^{\perp}$  has as its generator matrix the  $1 \times 6$  matrix

$$H = \begin{bmatrix} \overline{2} & \overline{3} & \overline{4} & \overline{5} & \overline{6} & \overline{7} \end{bmatrix}.$$

Then  $m := |\mathcal{C}| = 161051.$ 

(f) For C the same code as in (e), C has minimum distance d(C) = 3.

2. (a) (10 points) The integer 7919 is prime, and so we know  $\alpha = \overline{100}$  in  $\mathbb{F}_{7919}$  has a multiplicative inverse  $\alpha^{-1}$ . Find  $\alpha^{-1}$  explicitly, using the extended Euclid algorithm.

(b) (10 points) The polynomials

$$f(x) = x^2$$
$$g(x) = x^5 + x + 1$$

in  $\mathbb{F}_2[x]$  have no common factors. Hence there will exist polynomials a(x), b(x) in  $\mathbb{F}_2[x]$  satisfying a(x)f(x) + b(x)g(x) = 1. Find such polynomials a(x), b(x) explicitly, using the extended Euclid algorithm.

3. Let  $\mathcal{C} \subseteq (\mathbb{F}_5)^{12}$  be the following  $\mathbb{F}_5$ -linear code:

$$\mathcal{C} = \{ \mathbf{x} = [x_1, \dots, x_{12}] \in (\mathbb{F}_5)^{12} : x_1 = x_2 = x_3 = x_4, \\ x_5 = x_6 = x_7 = x_8, \\ x_9 = x_{10} = x_{11} = x_{12} \}$$

So  $\mathcal{C}$  is the third extension of the 4-fold repetition code over  $\mathbb{F}_5$ .

(a) (5 points) What is the dimension  $k = \dim_{\mathbb{F}_5}(\mathcal{C})$ ?

(b) 5 points) What is the 5-ary rate of C?

(c) (5 points) What is the minimum distance  $d(\mathcal{C})$ ?

- (d) (5 points) Write down a generator matrix G for  $\mathcal{C}$ .
- (e) (5 points) Write down a generator matrix H for its dual code  $\mathcal{C}^{\perp}$ .
- (f) (5 points) What is the 5-ary rate of its dual code  $\mathcal{C}^{\perp}$ ?

4. (a) (5 points) Find a representative for  $\overline{1000}$  in  $\mathbb{Z}/37$  that lies within the set of residues  $\{\overline{0}, \overline{1}, \overline{2}, \ldots, \overline{36}\}$ .

(b) (5 points) Do the same for  $\overline{1,000,000}$  in  $\mathbb{Z}/37$ .

(c) (10 points) Prove that if a number N is written in decimal notation with digits  $a_{\ell}a_{\ell-1}\cdots a_2a_1a_0$  (so that  $a_0$  is the ones digit,  $a_1$  is the tens digit,  $a_2$  the hundreds digit, etc), then in  $\mathbb{Z}/37$  one has

$$\overline{N} = \dots + \overline{a_5 a_4 a_3} + \overline{a_2 a_1 a_0}.$$

For example, in  $\mathbb{Z}/37$  one has  $\overline{41, 246, 789, 963} = \overline{41} + \overline{246} + \overline{789} + \overline{963}$ .

2