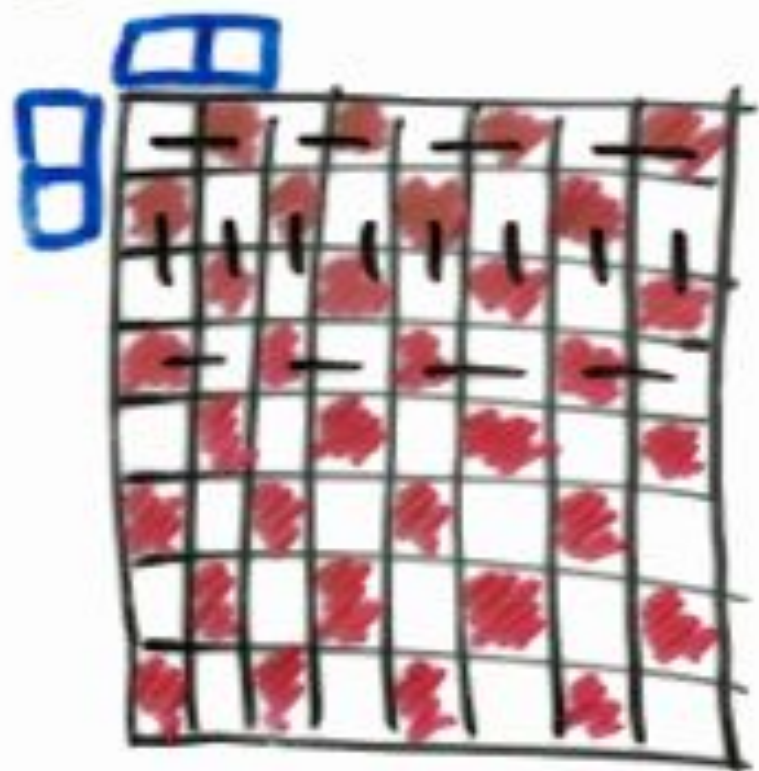


§§ 3-4 Techniques of Proof

"Proof" and "Theory" (or "Theorem") have very different meanings in mathematics, compared to other fields.

Ex Chessboard Tiling.

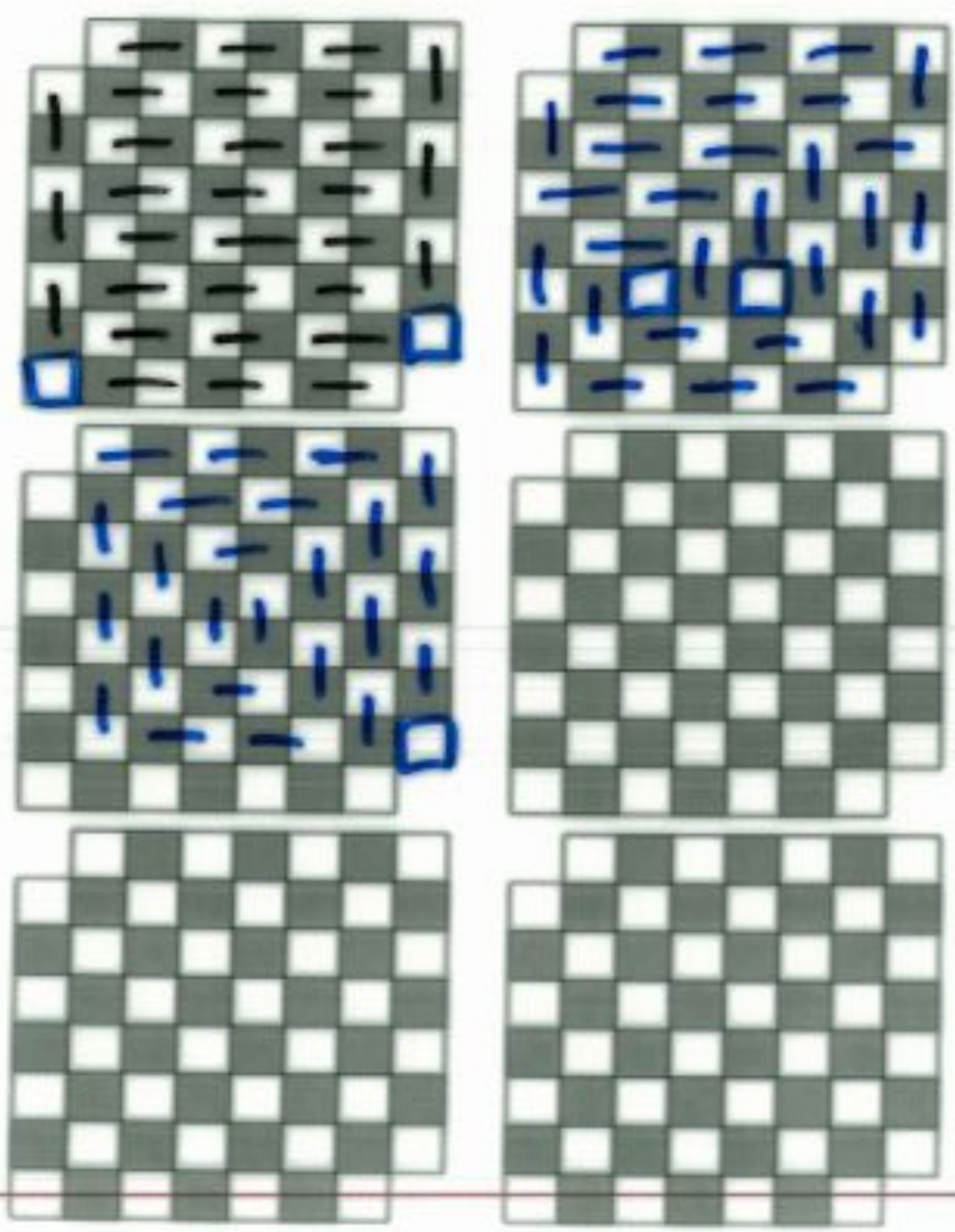


Suppose dominoes cover exactly 2 squares, either vertically or horizontally.

Can you cover all 64 squares using dominoes?

What if we remove two (opposite) corners of the board?

Mutilated Chessboard Problem



A mathematical proof is an airtight logical argument. A correct pf is true for all eternity. (!!)

In symbols, to prove

$$p \Rightarrow q$$

If p , then q

etc...

we want to create a series of implications w/ statements:

$$p \Rightarrow s_1 \Rightarrow s_2 \Rightarrow s_3 \Rightarrow \dots \Rightarrow s_n \Rightarrow q$$

KEY: if p is true and each implication is true, then q is true as well!

⚠ Basic issue for students in this section: what can you assume is true?

- basic arithmetic
- basic algebra
- also, these def^s:

An integer n is even if
 $n = 2k$ for some integer k .

n is odd if it has form
 $n = 2k + 1$ for some integer k .

Note these are iff, by convention
 so: if I know $n = 2(3) + 1$,
 then I can say n is odd.

Also if I know/assume n is odd,
 \exists integer $k \ni n = 2k + 1$.

Ex: Prove: If n is an odd integer, then n^2 is odd as well. 4

p : [$p(n)$] : n is odd

f : [$f(n)$] : n^2 is odd.

Pf: $p(n)$: n is an odd integer

$$\Rightarrow S_1: \exists k \text{ (integer)} \text{ s.t. } n = 2k+1.$$

$$\Rightarrow S_2: n^2 = (2k+1)^2$$

$$\Rightarrow S_3: n^2 = 4k^2 + 4k + 1$$

$$\Rightarrow S_4: n^2 = 2(\underbrace{2k^2 + 2k}_{\text{even}}) + 1$$

$$\Rightarrow f: n^2 \text{ is odd.}$$

ICK!

In words:

Let n be an odd integer, so \exists integer k s.t. $n = 2k+1$. Then

$$\begin{aligned} n^2 &= (2k+1)^2 \\ &= 4k^2 + 4k + 1 \\ &= 2(2k^2 + 2k) + 1 \end{aligned}$$

which has the form of an odd number. Hence n^2 is odd. ■ ● GET

etc.

Proving $p \Rightarrow q$ via $p \Rightarrow s_1 \Rightarrow \dots \Rightarrow s_n \Rightarrow q$ is called a direct proof. \exists other methods

Today / Friday: Pf by contrapositive, contradiction, (cases)

[induction to come later...]

Def The contrapositive of $p \Rightarrow q$ is.
 $\sim q \Rightarrow \sim p$.

Ex Write contrapositive of:

If $x > 1$, then $x^2 > 1$

$x^2 \leq 1 \Rightarrow x \leq 1$.

If it's raining, the sidewalk is wet.
sidewalk dry \Rightarrow not raining.

Contrapositive is useful b/c of this tautology: $(p \Rightarrow q) \Leftrightarrow (\sim q \Rightarrow \sim p)$

Proof by Contrapositive: prove $p \Rightarrow q$ indirectly, via a direct proof of $\neg q \Rightarrow \neg p$

Ex: Prove: for an integer n ,
 n^2 even $\Rightarrow n$ even.

Direct Proof Let n^2 be even,
so $n^2 = 2k$, some k .
Then $n = \sqrt{2k} = \sqrt{2} \cdot \sqrt{k}$???

Alternatively, we prove the
contrapositive, n odd $\Rightarrow n^2$ odd.
We did this
in ≈ 3 lines
last time.

Proof by Contradiction Let c be a 7
"contradiction" - a stmt which is
always false, e.g.: 2 is odd, $0=1$, etc.

Then: $(\sim p \Rightarrow c) \Leftrightarrow p$

if assuming $\sim p$ leads to
total nonsense (a contradiction),
then our assumption was
wrong, and p is true.

$[(p \wedge \sim q) \Rightarrow c] \Leftrightarrow (p \Rightarrow q)$

assume $p \Rightarrow q$ is false, show
that it leads to a contradiction
Thus our assumption was
wrong, and $p \Rightarrow q$ true.



} other tautologies used in
pls - see book - but they're
based on same ideas.

Proof: There are infinitely many primes.

Pf: Assume there are finitely many primes, exactly n of them:

$$p_1, p_2, \dots, p_n$$

Let $a = (p_1 p_2 \dots p_n) + 1$.

Clearly $a \neq p_i$ for any i , so a is not prime.

Thus a can be factored into primes. But $a \div p_i$ has a remainder of 1, so a cannot be factored.

Hence a is not prime and not composite - a contradiction!

Thus our assumption is wrong, and there are infinitely many primes!

Prove $\sqrt{2}$ is irrational

Rephrase If $x^2 = 2$, then x not rational

Pf:


Summary of terms

given $p \Rightarrow q$ implication
 $\sim q \Rightarrow \sim p$ contrapositive

} log. equiv.

$q \Rightarrow p$ converse
 $\sim p \Rightarrow \sim q$ inverse.

} log. equiv.
contrapos. of converse.

 In general \exists no connection between truth values of $p \Rightarrow q$ and its converse.

Ex. n^2 even \Leftrightarrow n even (T)
 n even \Rightarrow n^2 even (T)

$f(x)$ diff'ble \Rightarrow f continuous. (T)
cont \Rightarrow diff'ble. (F)

Deductive Reasoning: Showing a conclusion follows from certain premises. $(p \Rightarrow q)$

Inductive Reasoning: Pattern recognition

We often use INDuctive reasoning to find what to prove, and DEductive reasoning to prove it!

Last Method: Proof by Cases.

Ex Prove: $\left| \frac{a}{b} \right| = \frac{|a|}{|b|}$

Pf: Case 1 $a \geq 0, b > 0$. $\frac{|a|}{|b|} = \frac{a}{b} \geq 0$

\vdots

So $\frac{|a|}{|b|} = \frac{a}{b} = \left| \frac{a}{b} \right|$

Case 4 $a \leq 0, b < 0 \dots$



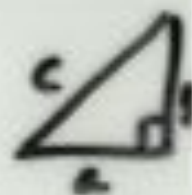
Last warning: to prove a statement is false, it suffices to give one counter-example:

Ex Give ctr-ex to "For all triangles w/ side lengths a, b, c , $a^2 + b^2 = c^2$."



But you can't prove a universal statement by checking 1 (or 1,000,000) examples:

Ex For a right Δ w/ hyp c , we have $a^2 + b^2 = c^2$.



Have to prove Pyth. Tho
can't just check
3-4-5 Δ .