

 1 Non associativity in the presence of round-off.

**Solution:** This is done in a class demo and the diary should be posted.

Here are the commands.

```
n = 10000;  
a = randn(n,1);  b = randn(n,1);  c = randn(n,1);  
t = ((a+b)+c == a+(b+c));  
sum(t)
```

Right-hand side in 3rd line returns 1 for each instance when the two numbers are the same.

 2 Find machine epsilon in matlab.

**Solution:**

```
u = 1;  
for i=0:999  
    fprintf(1,' i = %d , u = %e \n',i,u)  
    if (1.0 +u == 1.0) break, end
```

$$u = u/2;$$

end

$$u = u*2$$



**4** Proof of Lemma: If  $|\delta_i| \leq \underline{u}$  and  $n\underline{u} < 1$  then

$$\prod_{i=1}^n (1 + \delta_i) = 1 + \theta_n \quad \text{where} \quad |\theta_n| \leq \frac{n\underline{u}}{1 - n\underline{u}}$$

### **Solution:**

The proof is by induction on  $n$ .

1) Basis of induction. When  $n = 1$  then the product reduces to  $1 + \delta_i$  and so we can take  $\theta_n = \delta_n$  and we know that  $|\delta_n| \leq \underline{u}$  from the assumptions and so

$$|\theta_n| \leq \underline{u} \leq \frac{\underline{u}}{1 - \underline{u}},$$

as desired.

2) Induction step. Assume now that the result as stated is true for  $n$  and consider a product with  $n + 1$  terms:  $\prod_{i=1}^{n+1} (1 + \delta_i)$ . We can write this as  $(1 + \delta_{n+1})\prod_{i=1}^n (1 + \delta_i)$  and from the induction hypothesis


we get:

$$\prod_{i=1}^{n+1} (1 + \delta_i) = (1 + \theta_n)(1 + \delta_{n+1}) = 1 + \theta_n + \delta_{n+1} + \theta_n \delta_{n+1}$$

with  $\theta_n$  satisfying the inequality  $\theta_n \leq (n\underline{u})/(1 - n\underline{u})$ . We call  $\theta_{n+1}$  the quantity  $\theta_{n+1} = \theta_n + \delta_{n+1} + \theta_n \delta_{n+1}$ , and we have

$$\begin{aligned} |\theta_{n+1}| &= |\theta_n + \delta_{n+1} + \theta_n \delta_{n+1}| \\ &\leq \frac{n\underline{u}}{1 - n\underline{u}} + \underline{u} + \frac{n\underline{u}}{1 - n\underline{u}} \times \underline{u} \\ &= \frac{n\underline{u} + \underline{u}(1 - n\underline{u}) + n\underline{u}^2}{1 - n\underline{u}} = \frac{(n + 1)\underline{u}}{1 - n\underline{u}} \\ &\leq \frac{(n + 1)\underline{u}}{1 - (n + 1)\underline{u}} \end{aligned}$$

This establishes the result with  $n$  replaced by  $n + 1$  as wanted and completes the proof.  $\square$

 **5** Assume you use single precision for which you have  $\underline{u} = 2. \times 10^{-6}$ . What is the largest  $n$  for which  $n\underline{u} \leq 0.01$  holds? Any conclusions for the use of single precision arithmetic?

**Solution:** We need  $n \leq 0.01/(2.0 \times 10^{-4})$  which gives  $n \leq 5,000$ . Hence, single precision is inadequate for computations involving long inner products.

 **6** What does the main result on inner products imply for the case

when  $\mathbf{y} = \mathbf{x}$ ? [Contrast the relative accuracy you get in this case vs. the general case when  $\mathbf{y} \neq \mathbf{x}$ ]  $\square$

**Solution:** In this case we have

$$|fl(\mathbf{x}^T \mathbf{x}) - (\mathbf{x}^T \mathbf{x})| \leq \gamma_n \mathbf{x}^T \mathbf{x}$$

which implies that we will always have a small relative error. Not true for the general case because the final result (forward form)

$$|fl(\mathbf{y}^T \mathbf{x}) - (\mathbf{y}^T \mathbf{x})| \leq \gamma_n |\mathbf{y}|^T |\mathbf{x}|$$

does not imply a small relative error which would mean  $|fl(\mathbf{y}^T \mathbf{x}) - (\mathbf{y}^T \mathbf{x})| \leq \epsilon |\mathbf{y}^T \mathbf{x}|$  where  $\epsilon$  is small.  $\square$

**7** Show for any  $\mathbf{x}, \mathbf{y}$ , there exist  $\Delta \mathbf{x}, \Delta \mathbf{y}$  such that


$$fl(\mathbf{x}^T \mathbf{y}) = (\mathbf{x} + \Delta \mathbf{x})^T \mathbf{y}, \quad \text{with } |\Delta \mathbf{x}| \leq \gamma_n |\mathbf{x}|$$

$$fl(\mathbf{x}^T \mathbf{y}) = \mathbf{x}^T (\mathbf{y} + \Delta \mathbf{y}), \quad \text{with } |\Delta \mathbf{y}| \leq \gamma_n |\mathbf{y}|$$

**Solution:** The main result we proved is that

$$fl(\mathbf{y}^T \mathbf{x}) = \sum_{i=1}^n x_i y_i (1 + \theta_i) \quad \text{where } |\theta_i| \leq \gamma_n$$

The first relation comes from just attaching each  $(1 + \theta_i)$  to  $x_i$  so  $x_i$  is replaced by  $x_i + \theta_i x_i$  ... Similarly for the second relation.  $\square$


 8 (Continuation) Let  $A$  an  $m \times n$  matrix,  $x$  an  $n$ -vector, and  $y = Ax$ . Show that there exist a matrix  $\Delta A$  such

$$fl(y) = (A + \Delta A)x, \quad \text{with} \quad |\Delta A| \leq \gamma_n |A|$$

**Solution:** The result comes from applying the result on inner products to each entry  $y_i$  of  $y$  – which is the inner product of row  $i$  with  $y$ . We use the first of the two results above:

$$fl(y_i) = (a_{i,:} + \Delta a_{i,:})^T y \quad \text{with} \quad |\Delta a_{i,:}| \leq \gamma_n |a_{i,:}|$$

the result follows from expressing this in matrix form.  $\square$

 9 (Continuation) From the above derive a result about a column of the product of two matrices  $A$  and  $B$ . Does a similar result hold for the product  $AB$  as a whole?

**Solution:** We can have a result for each column since this is just a matrix-vector product. However this does not translate into a result for  $AB$  because the  $\Delta A$  we get for each column will depend on the column. Specifically, for the  $j$ -th column of  $B$  you will have a certain matrix  $(\Delta A)_j$  such that  $fl(AB(:, j)) = (A + (\Delta A)_j)B(:, j)$  with certain conditions as set in previous exercise. However this  $(\Delta A)_j$  is \*NOT\* the same matrix for each column. So you cannot

say  $fl(A) = (A + \Delta A)B, \dots \square$

### Supplemental notes

The importance of floating point analysis cannot be overstated. There were many instances where poor implementation of algorithms failed and led to - on occasion - disastrous results. One of the best examples is the failed launch of the European Ariane rocket in 1996 [Ariane flight V88]. See the story in this wikipedia page

[https://en.wikipedia.org/wiki/Ariane\\_flight\\_V88](https://en.wikipedia.org/wiki/Ariane_flight_V88)