SOLVING LINEAR SYSTEMS OF EQUATIONS

- Background on linear systems
- Gaussian elimination and the Gauss-Jordan algorithms
- The LU factorization
- Gaussian Elimination with pivoting – permutation matrices.
- Case of banded systems

Background: Linear systems

The Problem: A is an $n \times n$ matrix, and b a vector of $\mathbb{R}^n$. Find x such that:

$$Ax = b$$

➤ x is the unknown vector, b the right-hand side, and A is the coefficient matrix

Example:

$$\begin{cases} 2x_1 + 4x_2 + 4x_3 = 6 \\ x_1 + 5x_2 + 6x_3 = 4 \\ x_1 + 3x_2 + x_3 = 8 \end{cases} \quad \text{or} \quad \begin{pmatrix} 2 & 4 & 4 \\ 1 & 5 & 6 \\ 1 & 3 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 6 \\ 4 \\ 8 \end{pmatrix}$$

✍1 Solution of above system ?

- Standard mathematical solution by Cramer's rule:

$$\boxed{x_i = \det(A_i) / \det(A)}$$

A_i = matrix obtained by replacing i-th column by b.

- Note: This formula is useless in practice beyond $n = 3$ or $n = 4$.

Three situations:

1. The matrix A is nonsingular. There is a unique solution given by $x = A^{-1}b$.
2. The matrix A is singular and $b \in \mathrm{Ran}(A)$. There are infinitely many solutions.
3. The matrix A is singular and $b \notin \mathrm{Ran}(A)$. There are no solutions.

Example: (1) Let $A = \begin{pmatrix} 2 & 0 \\ 0 & 4 \end{pmatrix}$ $b = \begin{pmatrix} 1 \\ 8 \end{pmatrix}$. A is nonsingular ➤ a unique solution $x = \begin{pmatrix} 0.5 \\ 2 \end{pmatrix}$.

Example: (2) Case where A is singular & $b \in \text{Ran}(A)$:

$$A = \begin{pmatrix} 2 & 0 \\ 0 & 0 \end{pmatrix}, \quad b = \begin{pmatrix} 1 \\ 0 \end{pmatrix}.$$

➤ infinitely many solutions: $x(\alpha) = \begin{pmatrix} 0.5 \\ \alpha \end{pmatrix} \quad \forall \alpha.$

Example: (3) Let A same as above, but $b = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$.

➤ No solutions since 2nd equation cannot be satisfied

Triangular linear systems

Example:

$$\begin{pmatrix} 2 & 4 & 4 \\ 0 & 5 & -2 \\ 0 & 0 & 2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 2 \\ 1 \\ 4 \end{pmatrix}$$

- One equation can be trivially solved: the last one. $x_3 = 2$
- x_3 is known we can now solve the 2nd equation:

$$5x_2 - 2x_3 = 1 \;\rightarrow\; 5x_2 - 2 \times 2 = 1 \;\rightarrow\; x_2 = 1$$

- Finally x_1 can be determined similarly:

$$2x_1 + 4x_2 + 4x_3 = 2 \rightarrow \;\ldots\; \rightarrow \; x_1 = -5$$

ALGORITHM : 1. *Back-Substitution algorithm*

```
For $i = n : -1 : 1$ do:
    $t := b_i$
    For $j = i + 1 : n$ do
        $t := t - a_{ij} x_j$
    End
    $x_i = t / a_{ii}$
End
```

$$\left.\begin{array}{l}\text{For } j = i+1:n \text{ do} \\ \quad t := t - a_{ij}x_j \\ \text{End}\end{array}\right\} \begin{array}{l} t := b_i - (a_{i,i+1:n}, x_{i+1:n}) \\ \quad = b_i - \textit{an inner product}\end{array}$$

- We must require that each $a_{ii} \neq 0$
- Operation count?

Column version of back-substitution

Back-Substitution algorithm. Column version

```
For $j = n : -1 : 1$ do:
    $x_j = b_j / a_{jj}$
    For $i = 1 : j - 1$ do
        $b_i := b_i - x_j * a_{ij}$
    End
End
```

✍2 Justify the above algorithm [Show that it does indeed compute the solution]

➤ Analogous algorithms for *lower* triangular systems.

Linear Systems of Equations: Gaussian Elimination

- Back to arbitrary linear systems.

Principle of the method: Since triangular systems are easy to solve, we will transform a linear system into one that is triangular. Main operation: combine rows so that zeros appear in the required locations to make the system triangular.

Notation: use a Tableau:

$$\begin{cases} 2x_1 + 4x_2 + 4x_3 = 2 \\ x_1 + 3x_2 + 1x_3 = 1 \\ x_1 + 5x_2 + 6x_3 = -6 \end{cases} \quad \text{tableau:} \quad \begin{array}{|ccc|c|} \hline 2 & 4 & 4 & 2 \\ 1 & 3 & 1 & 1 \\ 1 & 5 & 6 & -6 \\ \hline \end{array}$$

- Main operation used: scaling and adding rows.

Example: Replace row2 by: row2 - $\frac{1}{2}$*row1:

$$\left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 1 & 3 & 1 & 1 \\ 1 & 5 & 6 & -6 \end{array}\right] \rightarrow \left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 1 & 5 & 6 & -6 \end{array}\right]$$

- This is equivalent to:

$$\left[\begin{array}{ccc} 1 & 0 & 0 \\ -\frac{1}{2} & 1 & 0 \\ 0 & 0 & 1 \end{array}\right] \times \left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 1 & 3 & 1 & 1 \\ 1 & 5 & 6 & -6 \end{array}\right] = \left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 1 & 5 & 6 & -6 \end{array}\right]$$

- The left-hand matrix is of the form $M = I - ve_1^T$ with $v = \begin{pmatrix} 0 \\ \frac{1}{2} \\ 0 \end{pmatrix}$

Linear Systems of Equations: Gaussian Elimination

Go back to original system. Step 1 must transform:

$$\begin{array}{|rrr|r|}\hline 2 & 4 & 4 & 2 \\ 1 & 3 & 1 & 1 \\ 1 & 5 & 6 & -6 \\ \hline \end{array} \text{ into: } \begin{array}{|ccc|c|}\hline x & x & x & x \\ 0 & x & x & x \\ 0 & x & x & x \\ \hline \end{array}$$

$$row_2 := row_2 - \tfrac{1}{2} \times row_1: \qquad row_3 := row_3 - \tfrac{1}{2} \times row_1:$$

$$\begin{array}{|rrr|r|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 1 & 5 & 6 & -6 \\ \hline \end{array} \qquad \begin{array}{|rrr|r|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \\ \hline \end{array}$$

- Equivalent to

$$\begin{array}{|ccc|}\hline 1 & 0 & 0 \\ -\frac{1}{2} & 1 & 0 \\ -\frac{1}{2} & 0 & 1 \\ \hline \end{array} \times \begin{array}{|ccc|c|}\hline 2 & 4 & 4 & 2 \\ 1 & 3 & 1 & 1 \\ 1 & 5 & 6 & -6 \\ \hline \end{array} = \begin{array}{|ccc|c|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \\ \hline \end{array}$$

$$[\boldsymbol{A}, \boldsymbol{b}] \rightarrow [\boldsymbol{M}_1\boldsymbol{A}, \boldsymbol{M}_1\boldsymbol{b}]; \quad \boldsymbol{M}_1 = \boldsymbol{I} - \boldsymbol{v}^{(1)}\boldsymbol{e}_1^T; \quad \boldsymbol{v}^{(1)} = \begin{pmatrix} 0 \\ \frac{1}{2} \\ \frac{1}{2} \end{pmatrix}$$

- New system $\boldsymbol{A}_1\boldsymbol{x} = \boldsymbol{b}_1$. Step 2 must now transform:

$$\begin{array}{|ccc|c|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \\ \hline \end{array} \text{ into: } \begin{array}{|ccc|c|}\hline x & x & x & x \\ 0 & x & x & x \\ 0 & 0 & x & x \\ \hline \end{array}$$

$$row_3 := row_3 - 3 \times row_2 : \rightarrow \left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 0 & 7 & -7 \end{array}\right]$$

- Equivalent to

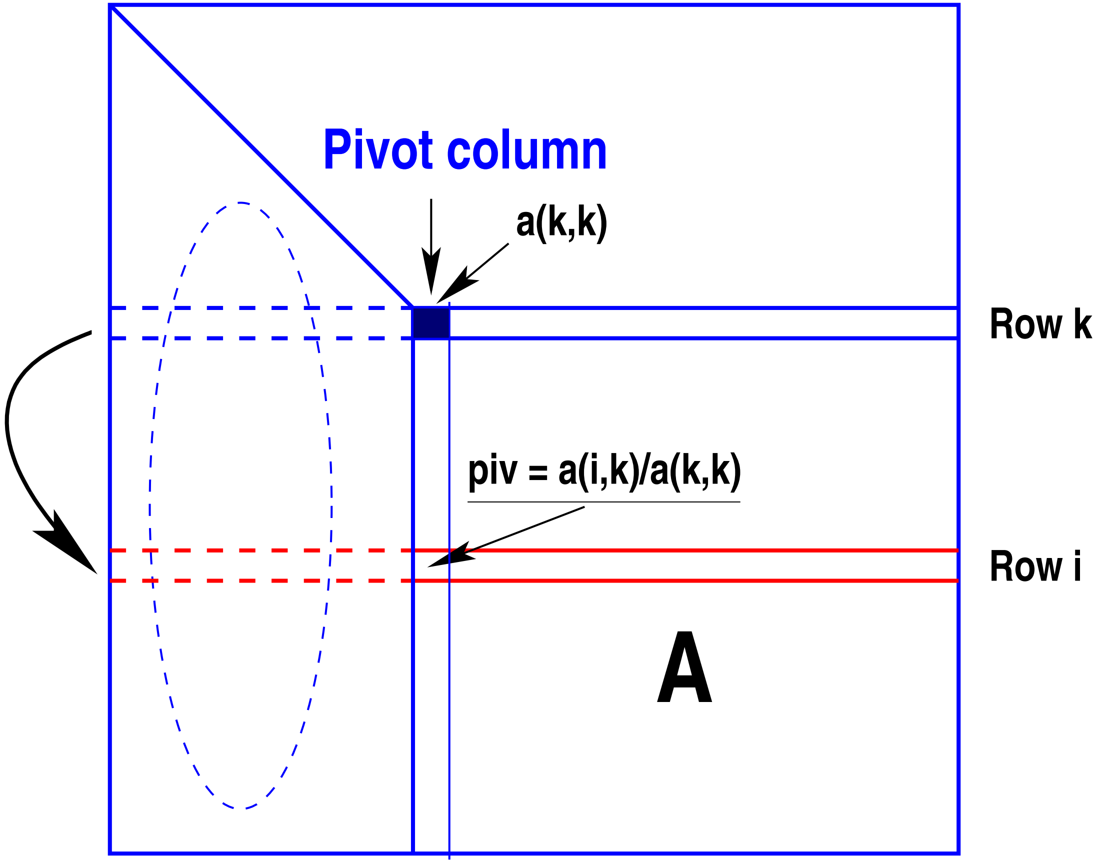
$$\left[\begin{array}{ccc} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -3 & 1 \end{array}\right] \times \left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \end{array}\right] = \left[\begin{array}{ccc|c} 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 0 & 7 & -7 \end{array}\right]$$

- Second transformation is as follows:

$$[A_1, b_1] \rightarrow [M_2 A_1, M_2 b_1]; \; M_2 = I - v^{(2)} e_2^T; \; v^{(2)} = \begin{pmatrix} 0 \\ 0 \\ 3 \end{pmatrix}$$

- Triangular system ➤ Solve.

Gaussian Elimination in a picture



```
For i=k+1:n Do:
    piv = a(i,k)/a(k,k)
    row(i):=row(i) – piv*row(k)
```

ALGORITHM : 2 ▪ *Gaussian Elimination*

```
For $k = 1 : n-1$ Do:
    For $i = k+1 : n$ Do:
        $piv := a_{ik}/a_{kk}$
        For $j := k+1 : n+1$ Do :
            $a_{ij} := a_{ij} - piv * a_{kj}$
        End
    End
End
```

- Operation count:

$$T = \sum_{k=1}^{n-1} \sum_{i=k+1}^{n} [1 + \sum_{j=k+1}^{n+1} 2] = \sum_{k=1}^{n-1} \sum_{i=k+1}^{n} (2(n-k)+3) = \ldots$$

✍3 Complete the above calculation. Order of the cost?

The LU factorization

- Now ignore the right-hand side from the transformations.

Observation: Gaussian elimination is equivalent to $n - 1$ successive Gaussian transformations, i.e., multiplications with matrices of the form $M_k = I - v^{(k)} e_k^T$, where the first k components of $v^{(k)}$ equal zero.

- Set $A_0 \equiv A$

$$\begin{aligned} A \rightarrow M_1 A_0 = A_1 &\rightarrow M_2 A_1 = A_2 \rightarrow M_3 A_2 = A_3 \cdots \\ &\rightarrow M_{n-1} A_{n-2} = A_{n-1} \equiv U \end{aligned}$$

- Last $A_k \equiv U$ is an upper triangular matrix.

- At each step we have: $\boxed{A_k = M_{k+1}^{-1} A_{k+1}}$. Therefore:

$$\begin{aligned} A_0 &= M_1^{-1} A_1 \\ &= M_1^{-1} M_2^{-1} A_2 \\ &= M_1^{-1} M_2^{-1} M_3^{-1} A_3 \\ &= \ldots \\ &= M_1^{-1} M_2^{-1} M_3^{-1} \cdots M_{n-1}^{-1} A_{n-1} \end{aligned}$$

- $$\boxed{L = M_1^{-1} M_2^{-1} M_3^{-1} \cdots M_{n-1}^{-1}}$$

- Note: L is Lower triangular, A_{n-1} is upper triangular

- LU decomposition : $A = LU$

How to get L?

$$L = M_1^{-1} M_2^{-1} M_3^{-1} \cdots M_{n-1}^{-1}$$

- Consider only the first 2 matrices in this product.
- Note $M_k^{-1} = (I - v^{(k)} e_k^T)^{-1} = (I + v^{(k)} e_k^T)$. So:

$$M_1^{-1} M_2^{-1} = (I + v^{(1)} e_1^T)(I + v^{(2)} e_2^T) = I + v^{(1)} e_1^T + v^{(2)} e_2^T.$$

- Generally,

$$M_1^{-1} M_2^{-1} \cdots M_k^{-1} = I + v^{(1)} e_1^T + v^{(2)} e_2^T + \cdots v^{(k)} e_k^T$$

The L factor is a lower triangular matrix with ones on the diagonal. Column k of L, contains the multipliers l_{ik} used in the k-th step of Gaussian elimination.

A matrix A has an LU decomposition if

$$\det(A(1:k,1:k)) \neq 0 \quad \text{for} \quad k = 1, \cdots, n-1.$$

In this case, the determinant of A satisfies:

$$\det A = \det(U) = \prod_{i=1}^{n} u_{ii}$$

If, in addition, A is nonsingular, then the LU factorization is unique.

✍4 Practical use: Show how to use the LU factorization to solve linear systems with the same matrix $\boldsymbol{A}$ and different $\boldsymbol{b}$'s.

✍5 LU factorization of the matrix $\boldsymbol{A} = \begin{pmatrix} \mathbf{2} & \mathbf{4} & \mathbf{4} \\ \mathbf{1} & \mathbf{5} & \mathbf{6} \\ \mathbf{1} & \mathbf{3} & \mathbf{1} \end{pmatrix}$?

✍6 Determinant of $\boldsymbol{A}$?

✍7 True or false: "Computing the LU factorization of matrix $\boldsymbol{A}$ involves more arithmetic operations than solving a linear system $\boldsymbol{Ax} = \boldsymbol{b}$ by Gaussian elimination".

Gauss-Jordan Elimination

Principle of the method: We will now transform the system into one that is even easier to solve than triangular systems, namely a diagonal system. The method is very similar to Gaussian Elimination. It is just a bit more expensive.

Back to original system. Step 1 must transform:

$$\begin{array}{|ccc|c|}\hline \mathbf{2} & \mathbf{4} & \mathbf{4} & \mathbf{2} \\ \mathbf{1} & \mathbf{3} & \mathbf{1} & \mathbf{1} \\ \mathbf{1} & \mathbf{5} & \mathbf{6} & \mathbf{-6} \\ \hline \end{array} \text{ into: } \begin{array}{|ccc|c|}\hline \boldsymbol{x} & \boldsymbol{x} & \boldsymbol{x} & \boldsymbol{x} \\ \mathbf{0} & \boldsymbol{x} & \boldsymbol{x} & \boldsymbol{x} \\ \mathbf{0} & \boldsymbol{x} & \boldsymbol{x} & \boldsymbol{x} \\ \hline \end{array}$$

$row_2 := row_2 - 0.5 \times row_1$: $row_3 := row_3 - 0.5 \times row_1$:

$$\begin{array}{|ccc|c|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 1 & 5 & 6 & -6 \\ \hline\end{array} \qquad \begin{array}{|ccc|c|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \\ \hline\end{array}$$

Step 2: $\begin{array}{|ccc|c|}\hline 2 & 4 & 4 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \\ \hline\end{array}$ into: $\begin{array}{|ccc|c|}\hline x & 0 & x & x \\ 0 & x & x & x \\ 0 & 0 & x & x \\ \hline\end{array}$

$row_1 := row_1 - 4 \times row_2$: $row_3 := row_3 - 3 \times row_2$:

$$\begin{array}{|ccc|c|}\hline 2 & 0 & 8 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 3 & 4 & -7 \\ \hline\end{array} \qquad \begin{array}{|ccc|c|}\hline 2 & 0 & 8 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 0 & 7 & -7 \\ \hline\end{array}$$

There is now a third step:

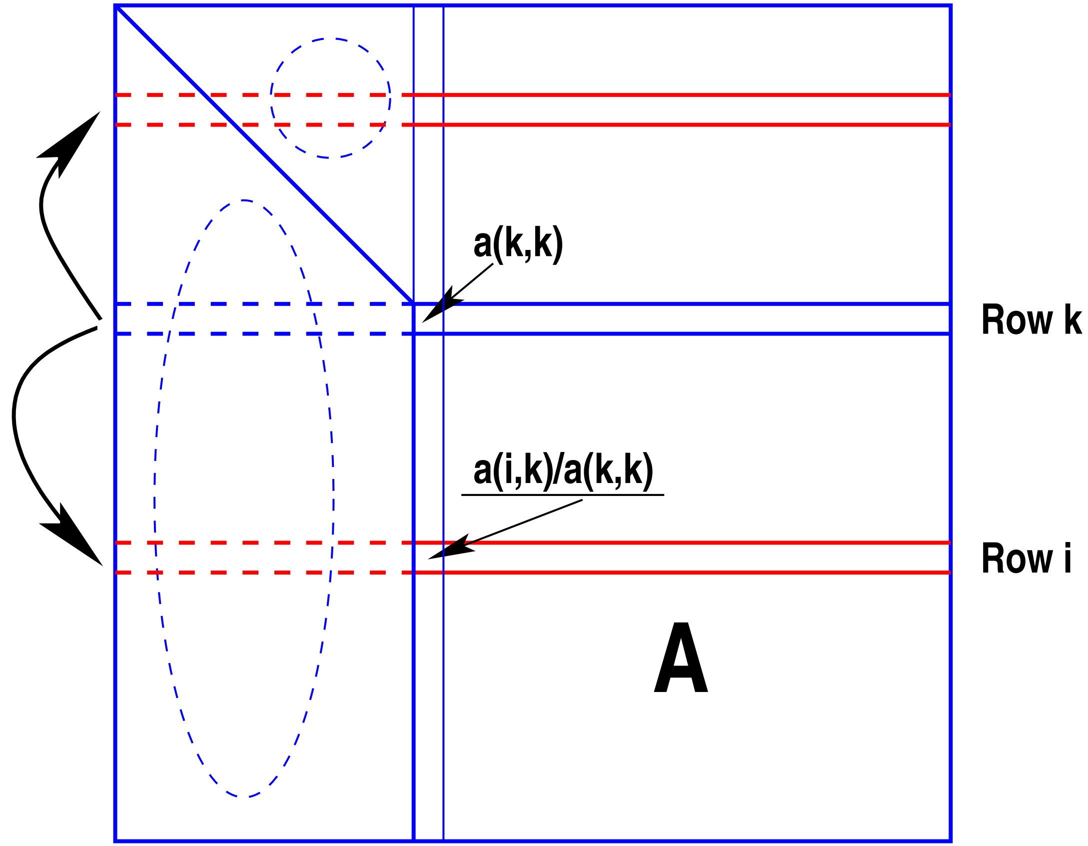
To transform: $\left[\begin{array}{ccc|c} 2 & 0 & 8 & 2 \\ 0 & 1 & -1 & 0 \\ 0 & 0 & 7 & -7 \end{array}\right]$ into: $\left[\begin{array}{ccc|c} x & 0 & 0 & x \\ 0 & x & 0 & x \\ 0 & 0 & x & x \end{array}\right]$

$row_1 := row_1 - \frac{8}{7} \times row_3$: $row_2 := row_2 - \frac{-1}{7} \times row_3$:

$$\left[\begin{array}{ccc|c} 2 & 0 & 0 & 10 \\ 0 & 1 & -1 & 0 \\ 0 & 0 & 7 & -7 \end{array}\right] \qquad \left[\begin{array}{ccc|c} 2 & 0 & 0 & 10 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 7 & -7 \end{array}\right]$$

Solution: $x_3 = -1$; $x_2 = -1$; $x_1 = 5$

Gauss-Jordan Elimination in a picture



ALGORITHM : 3 ■ *Gauss-Jordan elimination*

```
For $k = 1 : n$ Do:
    For $i = 1 : n$ and if $i! = k$ Do :
        $piv := a_{ik}/a_{kk}$
        For $j := k + 1 : n + 1$ Do :
            $a_{ij} := a_{ij} - piv * a_{kj}$
        End
    End
End
```

- Operation count:

$$T = \sum_{k=1}^{n} \sum_{i=1}^{n-1} [1 + \sum_{j=k+1}^{n+1} 2] = \sum_{k=1}^{n} \sum_{i=1}^{n-1} (2(n-k) + 3) = \cdots$$

✍8 Complete the above calculation. Order of the cost? How does it compare with Gaussian Elimination?

```
 function x = gaussj (A, b)
%---------------------------------------------------
%  function x = gaussj (A, b)
%  solves A x  = b by Gauss-Jordan elimination
%---------------------------------------------------
 n = size(A,1) ;
 A = [A,b];
 for k=1:n
   for i=1:n
     if (i ~= k)
         piv = A(i,k) / A(k,k) ;
         A(i,k+1:n+1) = A(i,k+1:n+1) - piv*A(k,k+1:n+1);
      end
   end
 end
 x = A(:,n+1) ./ diag(A) ;
```

Gaussian Elimination: Partial Pivoting

Consider again GE for the system:

$$\begin{cases} 2x_1 + 2x_2 + 4x_3 = 2 \\ x_1 + x_2 + x_3 = 1 \\ x_1 + 4x_2 + 6x_3 = -5 \end{cases}$$

Or:

$$\left[\begin{array}{ccc|c} 2 & 2 & 4 & 2 \\ 1 & 1 & 1 & 1 \\ 1 & 4 & 6 & -5 \end{array}\right]$$

- $row_2 := row_2 - \frac{1}{2} \times row_1$:

$$\left[\begin{array}{ccc|c} 2 & 2 & 4 & 2 \\ 0 & 0 & -1 & 0 \\ 1 & 4 & 6 & -5 \end{array}\right]$$

- $row_3 := row_3 - \frac{1}{2} \times row_1$:

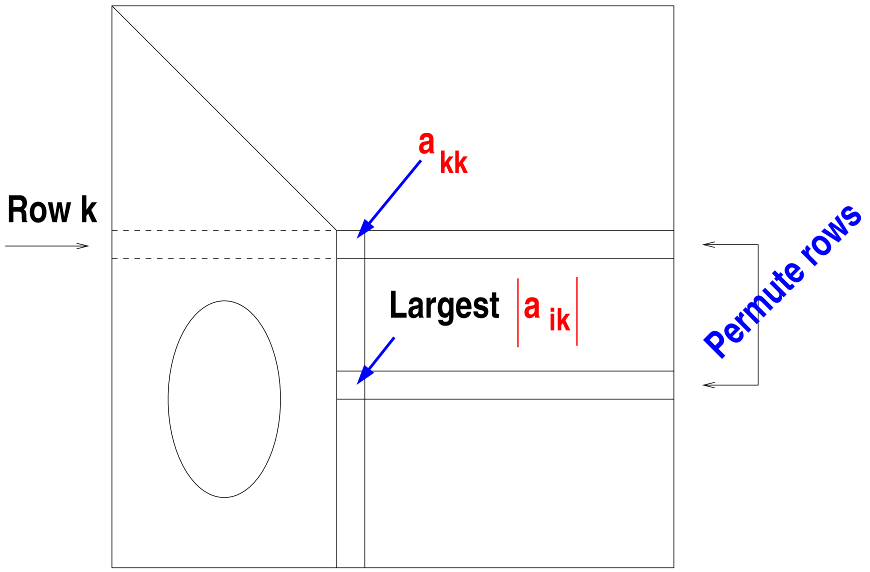
$$\left[\begin{array}{ccc|c} 2 & 2 & 4 & 2 \\ 0 & 0 & -1 & 0 \\ 0 & 3 & 4 & -6 \end{array}\right]$$

- Pivot a_{22} is zero. Solution : permute rows 2 and 3:

$$\left[\begin{array}{ccc|c} 2 & 2 & 4 & 2 \\ 0 & 3 & 4 & -6 \\ 0 & 0 & -1 & 0 \end{array}\right]$$

Gaussian Elimination with Partial Pivoting

Partial Pivoting



- General situation:

Always permute row k with row l such that

$$|a_{lk}| = \max_{i=k,\ldots,n} |a_{ik}|$$

- More 'stable' algorithm.

✍9 The matlab script *gaussp* will be provided. Explore it from the angle of an actual implementation in a language like C. Is it necessary to 'physically' move the rows? (moving data around is not free).

Pivoting and permutation matrices

- A permutation matrix is a matrix obtained from the identity matrix by permuting its rows

- For example for the permutation $\pi = \{3, 1, 4, 2\}$ we obtain

$$P = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{pmatrix}$$

- Important observation: the matrix PA is obtained from A by permuting its rows with the permutation π

$$\boxed{(PA)_{i,:} = A_{\pi(i),:}}$$

✍10 What is the matrix PA when

$$P = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{pmatrix} \quad A = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 5 & 6 & 7 & 8 \\ 9 & 0 & -1 & 2 \\ -3 & 4 & -5 & 6 \end{pmatrix} ?$$

➤ Any permutation matrix is the product of interchange permutations, which only swap two rows of I.

➤ Notation: E_{ij} = Identity with rows i and j swapped

Example: To obtain $\pi = \{3,1,4,2\}$ from $\pi = \{1,2,3,4\}$ – we need to swap $\pi(2) \leftrightarrow \pi(3)$ then $\pi(3) \leftrightarrow \pi(4)$ and finally $\pi(1) \leftrightarrow \pi(2)$. Hence:

$$P = \begin{pmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{pmatrix} = E_{1,2} \times E_{3,4} \times E_{2,3}$$

✍11 In the previous example where

```
>> A = [ 1 2 3 4; 5 6 7 8; 9  0 -1 2 ; -3 4 -5 6]
```

Matlab gives $det(A) = -896$. What is $det(PA)$?

- At each step of G.E. with partial pivoting:

$$\boxed{M_{k+1} E_{k+1} A_k = A_{k+1}}$$

where E_{k+1} encodes a swap of row $k+1$ with row $l > k+1$.

- Notes: (1) $E_i^{-1} = E_i$ and (2) $M_j^{-1} \times E_{k+1} = E_{k+1} \times \tilde{M}_j^{-1}$ for $k \ge j$, where $\tilde{M}_j$ has a permuted Gauss vector:

$$\begin{aligned}(I + v^{(j)} e_j^T) E_{k+1} &= E_{k+1} (I + E_{k+1} v^{(j)} e_j^T) \\ &\equiv E_{k+1} (I + \tilde{v}^{(j)} e_j^T) \\ &\equiv E_{k+1} \tilde{M}_j\end{aligned}$$

- Here we have used the fact that above row $k+1$, the permutation matrix E_{k+1} looks just like an identity matrix.

Result:

$$\begin{aligned} A_0 &= E_1 M_1^{-1} A_1 \\ &= E_1 M_1^{-1} E_2 M_2^{-1} A_2 = E_1 E_2 \tilde{M}_1^{-1} M_2^{-1} A_2 \\ &= E_1 E_2 \tilde{M}_1^{-1} M_2^{-1} E_3 M_3^{-1} A_3 \\ &= E_1 E_2 E_3 \tilde{M}_1^{-1} \tilde{M}_2^{-1} M_3^{-1} A_3 \\ &= \dots \\ &= E_1 \cdots E_{n-1} \ \times \ \tilde{M}_1^{-1} \tilde{M}_2^{-1} \tilde{M}_3^{-1} \cdots \tilde{M}_{n-1}^{-1} \ \times \ A_{n-1} \end{aligned}$$

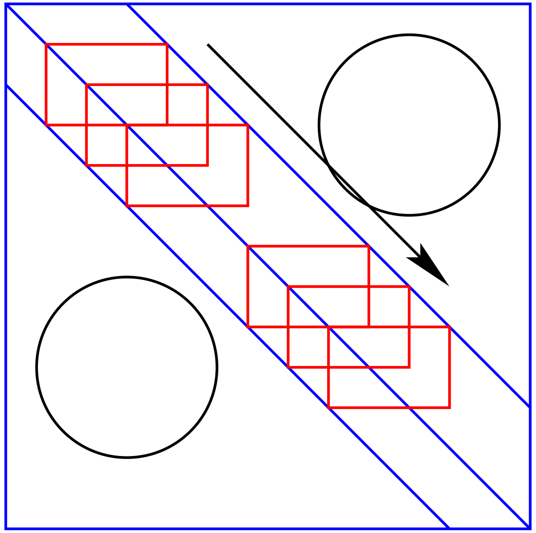
- In the end

$$PA = LU \text{ with } P = E_{n-1} \cdots E_1$$

Special case of banded matrices

- Banded matrices arise in many applications
- A has upper bandwidth q if $a_{ij} = 0 \textit{ for } j - i > q$
- A has lower bandwidth p if $a_{ij} = 0 \textit{ for } i - j > p$

✍12 Explain how GE would work on a banded system (you want to avoid operations involving zeros) – Hint: see picture



- Simplest case: tridiagonal ➤ $p = q = 1$.

➤ First observation: Gaussian elimination (no pivoting) preserves the initial banded form. Consider first step of Gaussian elimination:

```
2.  For $i = 2 : n$ Do:
3.      $a_{i1} := a_{i1}/a_{11}$ (pivots)
4.      For $j := 2 : n$ Do :
5.          $a_{ij} := a_{ij} - a_{i1} * a_{1j}$
6.      End
7.  End
```

➤ If A has upper bandwidth q and lower bandwidth p then so is the resulting $[L/U]$ matrix. ➤ Band form is preserved (induction)

✍13 Operation count?

What happens when partial pivoting is used?

If A has lower bandwidth p, upper bandwidth q, and if Gaussian elimination with partial pivoting is used, then the resulting U has upper bandwidth $p+q$. L has at most $p+1$ nonzero elements per column (bandedness is lost).

- Simplest case: tridiagonal ➤ $p = q = 1$.

Example:

$$A = \begin{pmatrix} 1 & 1 & 0 & 0 & 0 \\ 2 & 1 & 1 & 0 & 0 \\ 0 & 2 & 1 & 1 & 0 \\ 0 & 0 & 2 & 1 & 1 \\ 0 & 0 & 0 & 2 & 1 \end{pmatrix}$$